expkv

an expandable \(\key \rangle = \langle value \rangle \) implementation

Jonathan P. Spratte*

2020-07-04 V1.4

Abstract

expkv provides a small interface for $\langle key \rangle = \langle value \rangle$ parsing. The parsing macro is fully expandable, the $\langle code \rangle$ of your keys might be not. expkv is pretty fast, but not the fastest available $\langle key \rangle = \langle value \rangle$ solution (keyval is one and a half times as fast, but not expandable and it might strip braces it shouldn't have stripped).

Contents

1	Documentation	2
	1.1 Setting up Keys	2
	1.2 Parsing Keys	3
	1.3 Miscellaneous	4
	1.3.1 Other Macros	4
	1.3.2 Bugs	5
	1.3.3 Comparisons	6
	1.4 Examples	7
	1.4.1 Standard Use-Case	7
	1.4.2 An Expandable (key)=(value) Macro Using \ekvsneak	9
	1.5 Error Messages	11
	1.5.1 Load Time	11
	1.5.2 Defining Keys	11
	1.5.3 Using Keys	12
	1.6 License	12
2	Implementation	13
	2.1 The LATEX Package	13
	2.2 The Generic Code	13
Ind	lex	28

^{*}jspratte@yahoo.de

1 Documentation

expkv provides an expandable $\langle key \rangle = \langle value \rangle$ parser. The $\langle key \rangle = \langle value \rangle$ pairs should be given as a comma separated list and the separator between a $\langle key \rangle$ and the associated $\langle value \rangle$ should be an equal sign. Both, the commas and the equal signs, might be of category 12 (other) or 13 (active). To support this is necessary as for example babel turns characters active for some languages, for instance the equal sign is turned active for Turkish.

expkv is usable as generic code or as a LATEX package. To use it, just use one of:

```
\usepackage{expkv} % LaTeX
\input expkv % plainTeX
```

The LATEX package doesn't do more than expkv.tex, except calling \ProvidesPackage and setting things up such that expkv.tex will use \ProvidesFile.

In the expkv family are other packages contained which provide additional functionality. Those packages currently are:

```
expkvider a key-defining frontend for expkv using a \langle key \rangle = \langle value \rangle syntax expkvides define expandable \langle key \rangle = \langle value \rangle macros using expkv expkvider parse package and class options with expkv
```

Note that while the package names are stylised with a vertical rule, their names are all lower case with a hyphen (*e.g.*, expkv-def).

1.1 Setting up Keys

expkv provides a rather simple approach to setting up keys, similar to keyval. However there is an auxiliary package named expkvider which provides a more sophisticated interface, similar to well established packages like pgfkeys or l3keys.

Keys in $\exp_{\mathbf{k}\mathbf{v}}$ (as in almost all other $\langle key \rangle = \langle value \rangle$ implementations) belong to a *set* such that different sets can contain keys of the same name. Unlike many other implementations $\exp_{\mathbf{k}\mathbf{v}}$ doesn't provide means to set a default value, instead we have keys that take values and keys that don't (the latter are called NoVal keys by $\exp_{\mathbf{k}\mathbf{v}}$), but both can have the same name (on the user level).

The following macros are available to define new keys. Those macros containing "def" in their name can be prefixed by anything allowed to prefix \def (but don't use \outer, keys defined with it won't ever be usable), prefixes allowed for \let can prefix those with "let" in their name, accordingly. Neither \(\set \) nor \(\let key \) are allowed to be empty for new keys. \(\set \) will be used as is inside of \csname \(\cdot\).\\end{endcsname} and \(\let key \) will get \\detokenized.

Defines a $\langle key \rangle$ taking a value in a $\langle set \rangle$ to expand to $\langle code \rangle$. In $\langle code \rangle$ you can use #1 to refer to the given value.

 $\verb|\ekvdefNoVal| {$\langle set \rangle} {$\langle key \rangle} {$\langle code \rangle}$

Defines a no value taking (key) in a (set) to expand to (code).

\ekvlet \ekv

Let the value taking $\langle key \rangle$ in $\langle set \rangle$ to $\langle cs \rangle$, there are no checks on $\langle cs \rangle$ enforced.

\ekvletNoVal

 $\ensuremath{\ensuremath}\amb}\amb}\amb}}}}}}}}}}}}}}$

Let the no value taking $\langle key \rangle$ in $\langle set \rangle$ to $\langle cs \rangle$, it is not checked whether $\langle cs \rangle$ exists or that it takes no parameter.

\ekvletkv

 $\ensuremath{\ensuremath}\ensuremath{\ensuremath{\ens$

Let the $\langle key \rangle$ in $\langle set \rangle$ to $\langle key2 \rangle$ in $\langle set2 \rangle$, it is not checked whether that second key exists.

\ekvletkvNoVal

 $\verb|\ekvletkvNoVal{\(set\)}{(\key\)}{(\set2)}{(\key2)}|$

Let the $\langle key \rangle$ in $\langle set \rangle$ to $\langle key2 \rangle$ in $\langle set2 \rangle$, it is not checked whether that second key exists.

\ekvdefunknown

 $\ensuremath{\ensuremath{\mbox{\sc holdsymbol{\sc holdsymbol{\sc$

By default an error will be thrown if an unknown $\langle key \rangle$ is encountered. With this macro you can define $\langle code \rangle$ that will be executed for a given $\langle set \rangle$ when an unknown $\langle key \rangle$ with a $\langle value \rangle$ was encountered instead of throwing an error. You can refer to the given $\langle value \rangle$ with #1 and to the unknown $\langle key \rangle$'s name with #2 in $\langle code \rangle$.

\ekvdefunknownNoVal

 $\verb|\ekvdefunknownNoVal{$\langle set \rangle$} {\langle code \rangle$}$

As already explained for \ekvdefunknown, expkv would throw an error when encountering an unknown \(\key \). With this you can instead let it execute \(\cdot code \) if an unknown NoVal \(\key \) was encountered. You can refer to the given \(\key \) with #1 in \(\cdot code \).

1.2 Parsing Keys

\ekvset

 $\left(\left(set \right) \right) \left(\left(key \right) = \left(value \right), \ldots \right)$

Splits $\langle key \rangle = \langle value \rangle$ pairs on commas. From both $\langle key \rangle$ and $\langle value \rangle$ up to one space is stripped from both ends, if then only a braced group remains the braces are stripped as well. So $\langle kvset\{foo\}\{bar=baz\}\}$ and $\langle kvset\{foo\}\{\{bar\}=\{baz\}\}\}$ will both do $\langle foobarcode \rangle\{baz\}$, so you can hide commas, equal signs and spaces at the ends of either $\langle key \rangle$ or $\langle value \rangle$ by putting braces around them. If you omit the equal sign the code of the key created with the NoVal variants described in subsection 1.1 will be executed. If $\langle key \rangle = \langle value \rangle$ contains more than a single unhidden equal sign, it will be split at the first one and the others are considered part of the value. $\langle kvset \rangle$ should be nestable.

\ekvsetdef

 $\verb|\ekvsetdef| \langle cs \rangle \{ \langle set \rangle \}|$

With this function you can define a shorthand macro $\langle cs \rangle$ to parse keys of a specified $\langle set \rangle$. It is always defined $\lceil long \rceil$, but if you need to you can also prefix it with $\lceil long \rceil$. The resulting macro is a bit faster than the idiomatic definition:

 $\label{eq:cs} $$ \log \left(cs \right) #1{\ekvset{\langle set \rangle}} $$$

¹That order is correct, this way the code is faster.

\ekvparse

 $\ensuremath{\ensuremath{\mbox{\sc key}}\mbox{\sc key}}\=\ensuremath{\mbox{\sc key}}\mbox{\sc key}\=\ensuremath{\mbox{\sc key}}\mbox{\sc key}\$

This macro parses the $\langle key \rangle = \langle value \rangle$ pairs and provides those list elements which are only keys as the argument to $\langle cs1 \rangle$, and those which are a $\langle key \rangle = \langle value \rangle$ pair to $\langle cs2 \rangle$ as two arguments. It is fully expandable as well and returns the parsed list in \unexpanded, which has no effect outside of an \expanded or \edge def context². If you need control over the necessary steps of expansion you can use \expanded around it.

\ekvbreak, \ekvsneak, and \ekvchangeset and their relatives don't work in \ekvparse. It is analogue to expl3's \keyval_parse:NNn, but not with the same parsing rules - \keyval_parse:NNn throws an error on multiple equal signs per $\langle key \rangle = \langle value \rangle$ pair and on empty $\langle key \rangle$ names in a $\langle key \rangle = \langle value \rangle$ pair, both of which \ekvparse doesn't deal with.

As a small example:

 $\ensuremath{\mbox{\sc key}\mbox{\sc handlekey}\mbox{\sc handleke$

\handlekeyval{foo}{bar}\handlekey{key}\handlekeyval{baz}{zzz}

and afterwards \handlekey and \handlekeyval would have to further handle the \(\lambda key \rangle \). There are no macros like these two contained in expkv, you have to set them up yourself if you want to use \ekvparse (of course the names might differ). If you need the results of \ekvparse as the argument for another macro, you should use \expanded as only then the input stream will contain the output above:

\handle\kv{foo}{bar}\k{key}\kv{baz}{zzz}

•

1.3 Miscellaneous

1.3.1 Other Macros

exply provides some other macros which might be of interest.

\ekvVersion \ekvDate

These two macros store the version and date of the package.

\ekvifdefined \ekvifdefinedNoVal

```
\label{eq:continuity} $$ \ekvifdefined{\langle set \rangle} {\langle true \rangle} {\langle tr
```

These two macros test whether there is a $\langle key \rangle$ in $\langle set \rangle$. It is false if either a hash table entry doesn't exist for that key or its meaning is $\$ relax.

²This is a change in behaviour, previously (vo.3 and before) \ekvparse would expand in exactly two steps. This isn't always necessary, but makes the parsing considerably slower. If this is necessary for your application you can put an \expanded around it and will still be faster since you need only a single \expandafter this way.

\ekvifdefinedset

 $\ensuremath{\ensuremath}\amb}\amb}\amb}}}}}}}}}}}}}}$

This macro tests whether $\langle set \rangle$ is defined (which it is if at least one key was defined for it). If it is $\langle true \rangle$ will be run, else $\langle false \rangle$.

\ekvbreak \ekvbreak PreSneak \ekvbreak Post Sneak $\ensuremath{\ensuremath{\mbox{\ensuremath}\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath}\ens$

Gobbles the remainder of the current \ekvset macro and its argument list and reinserts \(\after \rangle \). So this can be used to break out of \ekvset. The first variant will also gobble anything that has been sneaked out using \ekvsneak or \ekvsneakPre, while \ekvbreakPreSneak will put \(\after \rangle \) before anything that has been smuggled and \ekvbreakPostSneak will put \(\after \rangle \) after the stuff that has been sneaked out.

\ekvsneak \ekvsneakPre $\ensuremath{\ensuremath{\mbox{\ensuremath}\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath}\ens$

Puts (after) after the effects of \ekvset. The first variant will put (after) after any other tokens which might have been sneaked before, while \ekvsneakPre will put (after) before other smuggled stuff. This reads and reinserts the remainder of the current \ekvset macro and its argument list to do its job. A small usage example is shown in subsubsection 1.4.2.

\ekvchangeset

 $\verb|\ekvchangeset{|\langle new-set \rangle|}|$

Replaces the current set with $\langle new-set \rangle$, so for the rest of the current $\langle ekvset call$, that call behaves as if it was called with $\langle ekvset \langle new-set \rangle \rangle$. Just like $\langle ekvsneak$ this reads and reinserts the remainder of the current $\langle ekvset \rangle$ macro to do its job. It is comparable to using $\langle key \rangle$.cd in pgfkeys.

\ekv@name \ekv@name@set \ekv@name@key $\label{eq:constraint} $\ekv@name{\langle set \rangle}_{\langle key \rangle}$$ \ekv@name@set{\langle set \rangle}$$ \ekv@name@key{\langle key \rangle}$$$

The names of the macros that correspond to a key in a set are build with these macros. The default definition of $\ensuremath{\mbox{\mbox{$\rm ekv@name@set}}}$ and the default of $\ensuremath{\mbox{\mbox{$\rm ekv@name@key}}}$. The complete name is build using $\ensuremath{\mbox{\mbox{$\rm ekv@name@key}}}$. For NoVal keys an additional N gets appended irrespective of these macros' definition, so their name is $\ensuremath{\mbox{\mbox{$\rm ekv@name@key}}}$. You might redefine $\ensuremath{\mbox{\mbox{$\rm ekv@name@key}}}$ and $\ensuremath{\mbox{\mbox{$\rm ekv@name@key}}}$ but $\ensuremath{\mbox{\mbox{$\rm don't\mbox{$\rm redefine}$}}}$ ekv@name!

1.3.2 Bugs

Just like keyval, expkv is bug free. But if you find bugshidden features³ you can tell me about them either via mail (see the first page) or directly on GitHub if you have an account there: https://github.com/Skillmon/tex_expkv

³Thanks, David!

1.3.3 Comparisons

Comparisons of speed are done with a very simple test key and the help of the l3benchmark package. The key and its usage should be equivalent to

```
\protected\ekvdef{test}{height}{\def\myheight{#1}}
\ekvset{test}{ height = 6 }
```

and only the usage of the key, not its definition, is benchmarked. For the impatient, the essence of these comparisons regarding speed and buggy behaviour is contained in Table 1.

As far as I know expkv is the only fully expandable $\langle key \rangle = \langle value \rangle$ parser. I tried to compare expkv to every $\langle key \rangle = \langle value \rangle$ package listed on CTAN, however, one might notice that some of those are missing from this list. That's because I didn't get the others to work due to bugs, or because they just provide wrappers around other packages in this list.

In this subsubsection is no benchmark of \ekvparse and \keyval_parse:NNn contained, as most other packages don't provide equivalent features to my knowledge. \ekvparse is slightly faster than \ekvset, but keep in mind that it does less. The same is true for \keyval_parse:NNn compared to \keys_set:nn of expl3 (where the difference is much bigger).

keyval is about 1.6 times faster and has a comparable feature set just a slightly different way how it handles keys without values. That might be considered a drawback, as it limits the versatility, but also as an advantage, as it might reduce doubled code. Keep in mind that as soon as someone loads xkeyval the performance of keyval gets replaced by xkeyval's.

Also keyval has a bug, which unfortunately can't really be resolved without breaking backwards compatibility for *many* documents, namely it strips braces from the argument before stripping spaces if the argument isn't surrounded by spaces, also it might strip more than one set of braces. Hence all of the following are equivalent in their outcome, though the last two lines should result in something different than the first two:

```
\setkeys{foo}{bar=baz}
\setkeys{foo}{bar= {baz}}
\setkeys{foo}{bar={ baz}}
\setkeys{foo}{bar={{baz}}}
```

xkeyval is roughly seventeen times slower, but it provides more functionality, e.g., it has choice keys, boolean keys, and so on. It contains the same bug as keyval as it has to be compatible with it by design (it replaces keyval's frontend), but also adds even more cases in which braces are stripped that shouldn't be stripped, worsening the situation.

Itxkeys is over 370 times slower – which is funny, because it aims to be "[...] faster [...] than these earlier packages [referring to keyval and xkeyval]." It needs more time to parse zero keys than four of the packages in this comparison need to parse 100 keys. Since it aims to have a bigger feature set than xkeyval, it most definitely also has a bigger feature set than expkv. Also, it can't parse \long input, so as soon as your values contain a \par, it'll throw errors. Furthermore, ltxkeys doesn't strip outer braces at all by design, which, imho, is a weird design choice. In addition ltxkeys loads catoptions which is known to introduce bugs (e.g., see https://tex.stackexchange.com/questions/461783).

l3keys is around six times slower, but has an, imho, great interface to define keys. It strips *all* outer spaces, even if somehow multiple spaces ended up on either end. It offers more features, but is pretty much bound to expl3 code. Whether that's a drawback is up to you. Note that this comparison uses the version contained in TEXLive 2019 (frozen) which is a bit slower than versions starting with TEXLive 2020.

pgfkeys is around 2.7 times slower for one key, but has an *enormous* feature set. It has the same or a very similar bug keyval has. The brace bug (and also the category fragility) can be fixed by pgfkeyx, but this package was last updated in 2012 and it slows down \pgfkeys by factor 8. Also I don't know whether this might introduce new bugs.

kvsetkeys with kvdefinekeys is about 3.7 times slower, but it works even if commas and equals have category codes different from 12 (just as some other packages in this list). Else the features of the keys are equal to those of keyval, the parser has more features, though.

options is 1.5 times slower for only a single value. It has a much bigger feature set. Unfortunately it also suffers from the premature unbracing bug keyval has.

simplekv is hard to compare because I don't speak French (so I don't understand the documentation). There was an update released on 2020-04-27 which greatly improved the package's performance and adds functionality so that it can be used more like most of the other $\langle key \rangle = \langle value \rangle$ packages. It has problems with stripping braces and spaces in a hard to predict manner just like keyval. Also, while it tries to be robust against category code changes of commas and equal signs, the used mechanism fails if the $\langle key \rangle = \langle value \rangle$ list already got tokenized. Regarding unknown keys it got a very interesting behaviour. It doesn't throw an error, but stores the $\langle value \rangle$ in a new entry accessible with \useKV. Also if you omit $\langle value \rangle$ it stores true for that $\langle key \rangle$. It is around 10% faster than $\exp_{\mathbf{k}\mathbf{v}}$.

yax is over twenty times slower. It has a pretty strange syntax, imho, and again a direct equivalent is hard to define. It has the premature unbracing bug, too. Also somehow loading yax broke options for me. The tested definition was:

1.4 Examples

1.4.1 Standard Use-Case

Say we have a macro for which we want to create a $\langle key \rangle = \langle value \rangle$ interface. The macro has a parameter, which is stored in the dimension \ourdim having a default value from its initialization. Now we want to be able to change that dimension with the width key to some specified value. For that we'd do

```
\newdimen\ourdim
\ourdim=150pt
\protected\ekvdef{our}{width}{\ourdim=#1\relax}
```

Table 1: Comparison of $\langle key \rangle = \langle value \rangle$ packages. The packages are ordered from fastest to slowest for one $\langle key \rangle = \langle value \rangle$ pair. Benchmarking was done using l3benchmark and the scripts in the Benchmarks folder of the git repository. The columns p_i are the polynomial coefficients of a linear fit to the run-time, p_0 can be interpreted as the overhead for initialisation and p_1 the cost per key. The T_0 column is the actual mean ops needed for an empty list argument, as the linear fit doesn't match that point well in general. The column "BB" lists whether the parsing is affected by some sort of brace bug, "CF" stands for category code fragile and lists whether the parsing breaks with active commas or equal signs.

Package	p_1	p_0	T_0	BB	CF	Date
keyval	13.7	1.5	7.0	yes	yes	2014-10-28
simplekv	18.7	5.3	17.7	yes	yes	2020-04-27
exp _k v	22.0	3.1	10.1	no	no	2020-06-21
options	24.4	12.0	20.4	yes	yes	2015-03-01
pgfkeys	24.6	45.5	53.3	yes	yes	2020-06-17
kvsetkeys	*	*	40.0	no	no	2019-12-15
l3keys	92.1	32.7	38.1	no	no	2020-06-18
xkeyval	257.1	173.7	164.5	yes	yes	2014-12-03
yax	440.2	76.3	113.9	yes	yes	2010-01-22
ltxkeys	3448.3	4470.0	5241.0	no	no	2012-11-17

^{*}For kvsetkeys the linear model used for the other packages is a poor fit, kvsetkeys seems to have approximately quadratic run-time, the coefficients of the second degree polynomial fit are $p_2 = 8.8$, $p_1 = 36.0$, and $p_0 = 81.8$. Of course the other packages might not really have linear run-time, but at least from 1 to 20 keys the fits don't seem too bad. If one extrapolates the fits for 100 $\langle key \rangle = \langle value \rangle$ pairs one finds that most of them match pretty well, the exception being ltxkeys, which behaves quadratic as well with $p_2 = 29.5$, $p_1 = 2828.9$, and $p_0 = 6741.2$.

as you can see, we use the set our here. We want the key to behave different if no value is specified. In that case the key should not use its initial value, but be smart and determine the available space from \hsize, so we also define

```
\protected \earline{$\protected \earline{$\protec
```

Now we set up our macro to use this $\langle key \rangle = \langle value \rangle$ interface

\protected\def\ourmacro#1{\begingroup\ekvset{our}{#1}\the\ourdim\endgroup}

Finally we can use our macro like in the following

```
\ourmacro{}\par 150.opt
\ourmacro{width}\par 192.85382pt
\ourmacro{width=5pt}\par 5.opt
```

The same key using explyidef Using explyidef we can set up the equivalent key using a $\langle key \rangle = \langle value \rangle$ interface, after the following we could use \ourmacro in the same way as above. explyidef will allocate and initialise \ourdim and define the width key \protected for us, so the result will be exactly the same – with the exception that the default will use \ourdim=.9\hsize\relax instead.

1.4.2 An Expandable (key)=(value) Macro Using \ekvsneak

Let's set up an expandable macro, that uses a $\langle key \rangle = \langle value \rangle$ interface. The problems we'll face for this are:

- 1. ignoring duplicate keys
- 2. default values for keys which weren't used
- 3. providing the values as the correct argument to a macro (ordered)

First we need to decide which \(\lambda \equiv \rangle \rangle \alpha \alpha \text{lue} \) parsing macro we want to do this with, \(\ekvset \text{ or \ekvsneak}, \text{ hence we'll choose \ekvset.} \) And we'll have to use \(\ekvset \text{ such that it builds a parsable list for our macro internals.} \) To gain back control after \(\ekvset \text{ is done we have to put an internal of our macro at the start of that list, so we use an internal key that uses \\ \ekvsneak \text{Pre after any user input.} \)

To ignore duplicates will be easy if the value of the key used last will be put first in the list, so the following will use \ekvsneakPre for the user-level keys. If we wanted some key for which the first usage should be the binding one we would use \ekvsneak instead for that key.

Providing default values can be done in different ways, we'll use a simple approach in which we'll just put the outcome of our keys if they were used with default values before the parsing list terminator.

Ordering the keys can be done simply by searching for a specific token for each argument which acts like a flag, so our sneaked out values will include specific tokens acting as markers.

Now that we have answers for our technical problems, we have to decide what our example macro should do. How about we define a macro that calculates the sine of a number and rounds that to a specified precision? As a small extra this macro should understand input in radian and degree and the used trigonometric function should be selectable as well. For the hard part of this task (expandably evaluating trigonometric functions) we'll use the xfp package.

First we set up our keys according to our earlier considerations and set up the user facing macro \sine. The end marker of the parsing list will be a \sine@stop token, which we don't need to define and we put our defaults right before it.

```
\RequirePackage {xfp}
\makeatletter
\ekvdef{expex}{f}{\ekvsneakPre{\f{\fin}}}
\ekvdef{expex}{fround}{\ekvsneakPre{\rnd{\fin}}}
\ekvdefNoVal{expex}{degree}{\ekvsneakPre{\deg{d}}}
\ekvdefNoVal{expex}{radian}{\ekvsneakPre{\deg{\fin}}}
\ekvdefNoVal{expex}{internal}{\ekvsneakPre{\sine@rnd}}
\newcommand*\sine[2]
{\ekvset{expex}{\fin},internal}\rnd{\fin}\deg{\fin}\sine@stop{\fin}
```

For the sake of simplicity we defined the macro \sine with two mandatory arguments, the first being the $\langle key \rangle = \langle value \rangle$ list, the second the argument to the trigonometric function. We could've used xparse's facilities here to define an expandable macro which takes an optional argument instead.

Now we need to define some internal macros to extract the value of each key's last usage (remember that this will be the group after the first special flag-token). For that we use one delimited macro per key.

After the macros \sine@rnd, \sine@deg, and \sine@f the macro \sine@final will see \sine@final $\{\langle f \rangle\}\{\langle degree/radian \rangle\}\{\langle round \rangle\}\{\langle num \rangle\}\}$. Now \sine@final has to expandably deal with those arguments such that the \fpeval macro of xfp gets the correct input. Luckily this is pretty straight forward in this example. In \fpeval the trigonometric functions have names such as sin or cos and the degree taking variants sind or cosd. And since the degree key puts a d in #2 and the radian key leaves #2 empty all we have to do to get the correct function name is stick the two together.

Let's test our macro:

The same macro using expkvics Using expkvics we can set up something equivalent with a bit less code. The implementation chosen in expkvics is more efficient than the example above and way easier to code for the user.

```
\makeatletter
\ekvcSplitAndForward\sine\sine@
{
    f=sin,
    unit=d,
    round=3,
}
\ekvcSecondaryKeys\sine
{
    nmeta degree={unit=d},
    nmeta radian={unit={}},
}
\newcommand*\sine@[4]{\fpeval{round(#1#2(#4),#3)}}
\makeatother
```

The resulting macro will behave just like the one previously defined, but will have an additional unit key, since in expkvics every argument must have a value taking key which defines it.

1.5 Error Messages

expkv should only send messages in case of errors, there are no warnings and no info messages. In this subsection those errors are listed.

1.5.1 Load Time

expkv.tex checks whether ε -TeX is available. If it isn't, an error will be thrown using \errmessage:

```
! expkv Error: e-TeX required.
```

1.5.2 Defining Keys

If you get any error from expkv while you're trying to define a key, the definition will be aborted and gobbled.

If you try to define a key with an empty set name you'll get:

```
! expkv Error: empty set name not allowed.
```

Similarly, if you try to define a key with an empty key name:

```
! expkv Error: empty key name not allowed.
```

Both of these messages are done in a way that doesn't throw additional errors due to \global, \long, etc., not being used correctly if you prefixed one of the defining macros.

1.5.3 Using Keys

This subsubsection contains the errors thrown during \ekvset. The errors are thrown in an expandable manner by providing an undefined macro. In the following messages $\langle key \rangle$ gets replaced with the problematic key's name, and $\langle set \rangle$ with the corresponding set. If any errors during $\langle key \rangle = \langle value \rangle$ handling are encountered, the entry in the comma separated list will be omitted after the error is thrown and the next $\langle key \rangle = \langle value \rangle$ pair will be parsed.

If you're using an undefined key you'll get:

```
! Undefined control sequence.
<argument> \! expkv Error:
unknown key ('<key>', set '<set>').
```

If you're using a key for which only a normal version and no NoVal version is defined, but don't provide a value, you'll get:

```
! Undefined control sequence.

<argument> \! expkv Error:

value required ('<key>', set '<set>').
```

If you're using a key for which only a NoVal version and no normal version is defined, but provide a value, you'll get:

```
! Undefined control sequence.

<argument> \! expkv Error:

value forbidden ('<key>', set '<set>').
```

If you're using a set for which you never executed one of the defining macros from subsection 1.1 you'll get a low level TeX error, as that isn't actively tested by the parser (and hence will lead to undefined behaviour and not be gracefully ignored). The error will look like

```
! Missing \endcsname inserted.
<to be read again>
\! expkv Error: Set '<set>' undefined.
```

1.6 License

Copyright © 2020 Jonathan P. Spratte

This work may be distributed and/or modified under the conditions of the LATEX Project Public License (LPPL), either version 1.3c of this license or (at your option) any later version. The latest version of this license is in the file:

```
http://www.latex-project.org/lppl.txt
This work is "maintained" (as per LPPL maintenance status) by
Jonathan P. Spratte.
```

2 Implementation

2.1 The LATEX Package

First we set up the LATEX package. That one doesn't really do much except \inputting the generic code and identifying itself as a package.

```
1 \def\ekv@tmp
2 {%
3  \ProvidesFile{expkv.tex}%
4      [\ekvDate\space v\ekvVersion\space an expandable key=val implementation]%
5 }
6 \input{expkv.tex}
7 \ProvidesPackage{expkv}%
8  [\ekvDate\space v\ekvVersion\space an expandable key=val implementation]
```

2.2 The Generic Code

The rest of this implementation will be the generic code.

We make sure that it's only input once:

```
9 \expandafter\ifx\csname ekvVersion\endcsname\relax
10 \else
11 \expandafter\endinput
12 \fi

Check whether ε-ΤΕΧ is available - expkv requires ε-ΤΕΧ.
13 \begingroup\expandafter\expandafter\expandafter\endgroup
14 \expandafter\ifx\csname numexpr\endcsname\relax
15 \errmessage{expkv requires e-TeX}
16 \expandafter\endinput
17 \fi
```

\ekvVersion \ekvDate

\ekvVersion We're on our first input, so lets store the version and date in a macro.

```
18 \def\ekvVersion{1.4}
19 \def\ekvDate{2020-07-04}
```

(End definition for \ekvVersion and \ekvDate. These functions are documented on page 4.)

If the LATEX format is loaded we want to be a good file and report back who we are, for this the package will have defined \ekv@tmp to use \ProvidesFile, else this will expand to a \relax and do no harm.

20 \csname ekv@tmp\endcsname

Store the category code of @ to later be able to reset it and change it to 11 for now.

- 21 \expandafter\chardef\csname ekv@tmp\endcsname=\catcode'\@
- 22 \catcode'\@=11

\ekv@tmp might later be reused to gobble any prefixes which might be provided to \ekvdef and similar in case the names are invalid, we just temporarily use it here as means to store the current category code of @ to restore it at the end of the file, we never care for the actual definition of it.

\@gobble
\@firstofone
\@firstoftwo
\@secondoftwo
\ekv@gobbleto@stop
\ekv@fi@gobble
\ekv@fi@firstoftwo
\ekv@fi@secondoftwo
\ekv@gobble@mark
\ekv@gobble@mark

Since branching tests are often more versatile than \if...\else...\fi constructs, we define helpers that are branching pretty fast. Also here are some other utility functions that just grab some tokens. The ones that are also contained in LATEX don't use the ekv prefix. Not all of the ones defined here are really needed by expkv but are provided because packages like expkvider or expkvider need them (and I don't want to define them in each package which might need them).

```
23 \long\def\@gobble#1{}
24 \long\def\@firstofone#1{#1}
25 \long\def\@firstoftwo#1#2{#1}
26 \long\def\@secondoftwo#1#2{#2}
27 \long\def\ekv@fi@gobble\fi\@firstofone#1{\fi}
28 \long\def\ekv@fi@firstoftwo\fi\@secondoftwo#1#2{\fi#1}
29 \long\def\ekv@fi@secondoftwo\fi\@firstoftwo#1#2{\fi#2}
30 \long\def\ekv@gobbleto@stop#1\ekv@stop{}
31 \def\ekv@gobble@mark\ekv@mark{}
32 \long\def\ekv@gobble@from@mark@to@stop\ekv@mark#1\ekv@stop{}
```

(End definition for \@gobble and others.)

As you can see \ekv@gobbleto@stop uses a special marker \ekv@stop. The package will use three such markers, the one you've seen already, \ekv@mark and \ekv@nil. Contrarily to how for instance expl3 does things, we don't define them, as we don't need them to have an actual meaning. This has the advantage that if they somehow get expanded – which should never happen if things work out – they'll throw an error directly.

\ekv@ifempty@
\ekv@ifempty@true
\ekv@ifempty@false
\ekv@ifempty@true@F
\ekv@ifempty@true@F@gobble
\ekv@ifempty@true@F@gobbletwo

We can test for a lot of things building on an if-empty test, so lets define a really fast one. Since some tests might have reversed logic (true if something is not empty) we also set up macros for the reversed branches.

(End definition for \ekv@ifempty and others.)

\ekv@ifblank \ekv@ifblank@

The obvious test that can be based on an if-empty is if-blank, meaning a test checking whether the argument is empty or consists only of spaces. Our version here will be tweaked a bit, as we want to check this, but with one leading \ekv@mark token that is to be ignored. The wrapper \ekv@ifblank will not be used by expkv for speed reasons but expkviopt uses it.

```
\long\def\ekv@ifblank#1%
48 \{%
49 \ekv@ifblank@#1\ekv@nil\ekv@ifempty@B\ekv@ifempty@true
```

\ekv@ifdefined

We'll need to check whether something is defined quite frequently, so why not define a macro that does this. The following test is expandable, slower than the typical expandable test for undefined control sequences, but faster for defined ones. Since we want to be as fast as possible for correct input, this is to be preferred.

```
63 \def\ekv@ifdefined#1%
64 {%
65  \expandafter
66  \ifx\csname\ifcsname #1\endcsname #1\else relax\fi\endcsname\relax
67  \ekv@fi@secondoftwo
68  \fi
69  \@firstoftwo
60 }
```

(End definition for \ekv@ifdefined.)

\ekv@name \ekv@name@set \ekv@name@key

\ekv@name The keys will all follow the same naming scheme, so we define it here.

```
61 \def\ekv@name#1#2{\ekv@name@set{#1}\ekv@name@key{\detokenize{#2}}}
62 \def\ekv@name@set#1{ekv#1(}
63 \def\ekv@name@key#1{#1)}
```

(End definition for \ekv@name, \ekv@name@set, and \ekv@name@key. These functions are documented on page 5.)

\ekv@undefined@set

We can misuse the macro name we use to expandably store the set-name in a single token — since this increases performance drastically, especially for long set-names — to throw a more meaningful error message in case a set isn't defined. The name of \ekv@undefined@set is a little bit misleading, as it is called in either case inside of \csname, but the result will be a control sequence with meaning \relax if the set is undefined, hence will break the \csname building the key-macro which will throw the error message.

```
64 \def\ekv@undefined@set#1{! expkv Error: Set '#1' undefined.}

(End definition for \ekv@undefined@set.)
```

\ekv@checkvalid

We place some restrictions on the allowed names, though, namely sets and keys are not allowed to be empty – blanks are fine (meaning set- or key-names consisting of spaces). The \def\ekv@tmp gobbles any TeX prefixes which would otherwise throw errors. This will, however, break the package if an \outer has been gobbled this way. I consider that good, because keys shouldn't be defined \outer anyways.

```
65 \protected\def\ekv@checkvalid#1#2%
66 {%
67  \ekv@ifempty{#1}%
68     {%
69      \def\ekv@tmp{}%
70      \errmessage{expkv Error: empty set name not allowed}%
71     }%
72     {%
73      \ekv@ifempty{#2}%
74     {%
```

```
\def\ekv@tmp{}%
                                     \errmessage{expkv Error: empty key name not allowed}%
                      76
                                   }%
                                   \@secondoftwo
                      78
                               }%
                      79
                             \@gobble
                      80
                          }
                     (End definition for \ekv@checkvalid.)
      \ekvifdefined And provide user-level macros to test whether a key is defined.
 \ekvifdefinedNoVal
                      %2 \def\ekvifdefined#1#2{\ekv@ifdefined{\ekv@name{#1}{#2}}}
                      83 \def\ekvifdefinedNoVal#1#2{\ekv@ifdefined{\ekv@name{#1}{#2}N}}
                     (End definition for \ekvifdefined and \ekvifdefinedNoVal. These functions are documented on page 4.)
            \ekvdef Set up the key defining macros \ekvdef etc.
       \ekvdefNoVal
                      84 \protected\long\def\ekvdef#1#2#3%
            \ekvlet
                          {%
       \ekvletNoVal
                             \ekv@checkvalid{#1}{#2}%
                               {%
          \ekvletkv
                                 \expandafter\def\csname\ekv@name{#1}{#2}\endcsname##1{#3}%
     \ekvletkvNoVal
                                 \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
     \ekvdefunknown
\ekvdefunknownNoVal
                          }
                        \protected\long\def\ekvdefNoVal#1#2#3%
                      92
                      93
                             \ekv@checkvalid{#1}{#2}%
                      94
                      95
                                 \expandafter\def\csname\ekv@name{#1}{#2}N\endcsname{#3}%
                                 \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                          }
                        \protected\def\ekvlet#1#2#3%
                             \ekv@checkvalid{#1}{#2}%
                                 \expandafter\let\csname\ekv@name{#1}{#2}\endcsname#3%
                                 \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                     106
                          }
                        \protected\def\ekvletNoVal#1#2#3%
                     108
                     109
                             \ekv@checkvalid{#1}{#2}%
                                 \expandafter\let\csname\ekv@name{#1}{#2}N\endcsname#3%
                                 \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                     114
                          }
                     116 \protected\def\ekvletkv#1#2#3#4%
                             \ekv@checkvalid{#1}{#2}%
                     118
                     119
                                 \expandafter\let\csname\ekv@name{#1}{#2}\expandafter\endcsname
                                 \csname\ekv@name{#3}{#4}\endcsname
```

```
\expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                            }%
                   123
                        }
                   124
                      \protected\def\ekvletkvNoVal#1#2#3#4%
                   126
                          \ekv@checkvalid{#1}{#2}%
                   128
                               \expandafter\let\csname\ekv@name{#1}{#2}N\expandafter\endcsname
                   129
                               \csname\ekv@name{#3}{#4}N\endcsname
                   130
                               \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                   132
                        }
                      \protected\def\ekvdefunknown#1#2%
                   134
                          \ekv@checkvalid{#1}{.}%
                   136
                               \expandafter\def\csname\ekv@name{#1}{}u\endcsname##1##2{#2}%
                   138
                               \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                        }
                   141
                      \protected\def\ekvdefunknownNoVal#1#2%
                   142
                          \ekv@checkvalid{#1}{.}%
                             {%
                   145
                               \expandafter\def\csname\ekv@name{#1}{}uN\endcsname##1{#2}%
                   146
                               \expandafter\ekv@defsetmacro\csname\ekv@undefined@set{#1}\endcsname{#1}%
                   147
                   148
                        }
                   149
                  (End definition for \ekvdef and others. These functions are documented on page 2.)
                  In order to enhance the speed the set name given to \ekvset will be turned into a control
                   sequence pretty early, so we have to define that control sequence.
                      \protected\def\ekv@defsetmacro#1#2%
                        {%
                             \edef#1##1{\ekv@name@set{#2}\ekv@name@key{\noexpand\detokenize{##1}}}%
                        }
                  (End definition for \ekv@defsetmacro.)
\ekvifdefinedset
                   156 \def\ekvifdefinedset#1%
                        {%
                          \ekv@ifdefined{\ekv@undefined@set{#1}}%
                   (End definition for \ekvifdefinedset. This function is documented on page 5.)
```

\extit{ekvset} Set up \extit{ekvset}, which should not be affected by active commas and equal signs. The equal signs are a bit harder to cope with and we'll do that later, but the active commas can be handled by just doing two comma-splitting loops one at actives one at others.

That's why we define \ekvset here with a temporary meaning just to set up the things with two different category codes. #1 will be a $_{13}$ and #2 will be a $_{13}$.

```
160 \def\ekvset#1#2{%
161 \endgroup
162 \long\def\ekvset##1##2%
163 {%
164 \expandafter\ekv@set\csname\ekv@undefined@set{##1}\endcsname
165 \ekv@mark##2#1\ekv@stop#1{}%
166 }
```

(End definition for \ekvset. This function is documented on page 3.)

\ekv@set

\ekv@set will split the $\langle key \rangle = \langle value \rangle$ list at active commas. Then it has to check whether there were unprotected other commas and resplit there.

```
167 \long\def\ekv@set##1##2#1%
168 {%
```

Test whether we're at the end, if so invoke \ekv@endset,

 $\label{lem:commark} $$ \ekv@gobble@from@mark@to@stop##2\ekv@endset\ekv@stop else go on with other commas,$

```
\ekv@set@other##1##2,\ekv@stop,%
```

and get the next active comma delimited $\langle key \rangle = \langle value \rangle$ pair.

(End definition for \ekv@set.)

\ekv@endset

\ekv@endset is a hungry little macro. It will eat everything that remains of \ekv@set and unbrace the sneaked stuff.

```
173 \long\def\ekv@endset
174 \ekv@stop\ekv@set@other##1,\ekv@stop,\ekv@set##2\ekv@mark
175 ##3%
176 {##3}
```

 $(End\ definition\ for\ \verb+\ekv@endset.)$

\ekv@eq@other \ekv@eq@active

Splitting at equal signs will be done in a way that checks whether there is an equal sign and splits at the same time. This gets quite messy and the code might look complicated, but this is pretty fast (faster than first checking for an equal sign and splitting if one is found). The splitting code will be adapted for \ekvset and \ekvparse to get the most speed, but some of these macros don't require such adaptions. \ekv@eq@other and \ekv@eq@other will split the argument at the first equal sign and insert the macro which comes after the first following \ekv@mark. This allows for fast branching based on TeX's argument grabbing rules and we don't have to split after the branching if the equal sign was there.

```
177 \long\def\ekv@eq@other##1=##2\ekv@mark##3##4\ekv@stop
178 {%
179 ##3##1\ekv@stop\ekv@mark##2%
180 }
181 \long\def\ekv@eq@active##1#2##2\ekv@mark##3##4\ekv@stop
182 {%
183 ##3##1\ekv@stop\ekv@mark##2%
184 }
```

(End definition for \ekv@eq@other and \ekv@eq@active.)

\ekv@set@other

The macro \ekv@set@other is guaranteed to get only single $\langle key \rangle = \langle value \rangle$ pairs.

```
185 \long\def\ekv@set@other##1##2,%
186 {%
```

First we test whether we're done.

87 \ekv@gobble@from@mark@to@stop##2\ekv@endset@other\ekv@stop

If not we split at the equal sign of category other.

And insert the set name and the next recursion step of \ekv@set@other.

```
190 ##1%
191 \ekv@set@other##1\ekv@mark
192 }
```

(End definition for \ekv@set@other.)

\ekv@set@eq@other@a \ekv@set@eq@other@b The first of these two macros runs the split-test for equal signs of category active. It will only be inserted if the $\langle key \rangle = \langle value \rangle$ pair contained at least one equal sign of category other and ##1 will contain everything up to that equal sign.

```
193 \long\def\ekv@set@eq@other@a##1\ekv@stop
194 {%
195 \ekv@eq@active##1\ekv@mark\ekv@set@eq@other@active@a
196 #2\ekv@mark\ekv@set@eq@other@b\ekv@stop
197 }
```

The second macro will have been called by \ekv@eq@active if no active equal sign was found. All it does is remove the excess tokens of that test and forward the $\langle key \rangle = \langle value \rangle$ pair to \ekv@set@pair.

```
198 \long\def\ekv@set@eq@other@b
199 ##1\ekv@nil\ekv@mark\ekv@set@eq@other@active@a\ekv@stop\ekv@mark
200 {%
201 \ekv@strip{##1}\ekv@set@pair
202 }
```

(End definition for \ekv@set@eq@other@a and \ekv@set@eq@other@b.)

\ekv@set@eq@other@active@a \ekv@set@eq@other@active@b \ekv@set@eq@other@active@a will be called if the $\langle key \rangle = \langle value \rangle$ pair was wrongly split on an equal sign of category other but has an earlier equal sign of category active. ##1 will be the contents up to the active equal sign and ##2 everything that remains until the first found other equal sign. It has to reinsert the equal sign and passes things on to \ekv@set@eq@other@active@b which calls \ekv@set@pair on the then correctly split $\langle key \rangle = \langle value \rangle$ pair.

```
long\def\ekv@set@eq@other@active@a##1\ekv@stop##2\ekv@nil\ekv@mark
{%
long\def\ekv@set@eq@other@active@b{##1}##2=%
long\def\ekv@set@eq@other@active@b##1%
{%
long\def\ekv@set@eq@other@active@b##1%
ekv@strip{##1}\ekv@set@pair
}
```

 $(\textit{End definition for } \verb|\ensuremath{|} ekv@set@eq@other@active@b.)$

\ekv@set@eq@active \ekv@set@eq@active@ \ekv@set@eq@active will be called when there was no equal sign of category other in the $\langle key \rangle = \langle value \rangle$ pair. It removes the excess tokens of the prior test and split-checks for an active equal sign.

```
long\def\ekv@set@eq@active
##1\ekv@nil\ekv@mark\ekv@set@eq@other@a\ekv@stop\ekv@mark
{%
lekv@eq@active##1\ekv@nil\ekv@mark\ekv@set@eq@active@
#2\ekv@mark\ekv@set@noeq\ekv@stop
}
```

If an active equal sign was found in $\ensuremath{\mbox{\mbox{$\sim$}}}$ derive we'll have to pass the now split $(\ensuremath{\mbox{$key$}}) = (\ensuremath{\mbox{\sim}})$ pair on to $\ensuremath{\mbox{$\sim$}}$ or to $\ensuremath{\mbox{$\sim$}}$ derive we'll have to pass the now split $(\ensuremath{\mbox{$\sim$}})$ pair on to $\ensuremath{\mbox{$\sim$}}$ or to $\ensuremath{\mbox{$\sim$}}$ derive we'll have to pass the now split $(\ensuremath{\mbox{$\sim$}})$ pair on to $\ensuremath{\mbox{$\sim$}}$ derive \ensure

```
217 \long\def\ekv@set@eq@active@##1\ekv@stop
218 {%
219 \ekv@strip{##1}\ekv@set@pair
220 }
```

(End definition for \ekv@set@eq@active and \ekv@set@eq@active@.)

\ekv@set@noeq

If no active equal sign was found by \ekv@set@eq@active there is no equal sign contained in the parsed list entry. In that case we have to check whether the entry is blank in order to ignore it (in which case we'll have to gobble the set-name which was put after these tests by \ekv@set@other). Else this is a NoVal key and the entry is passed on to \ekv@set@key.

```
221 \long\def\ekv@set@noeq##1\ekv@mark\ekv@set@eq@active@\ekv@stop\ekv@mark

222 {%
223 \ekv@ifblank@##1\ekv@nil\ekv@ifempty@B\ekv@ifempty@true@F@gobble

224 \ekv@ifempty@A\ekv@ifempty@B\@firstofone

225 {\ekv@strip{##1}\ekv@set@key}%

226 }
```

(End definition for \ekv@set@noeq.)

\ekv@endset@other

All that's left for \ekv@set@other is the macro which breaks the recursion loop at the end. This is done by gobbling all the remaining tokens.

(End definition for \ekv@endset@other.)

\ekvbreakPreSneak be useful for.

\ekvbreak Provide macros that can completely stop the parsing of \ekvset, who knows what it'll kPreSneak be useful for.

```
234 \long\def\ekvbreak##1##2\ekv@stop#1##3{##1}
235 \long\def\ekvbreakPreSneak ##1##2\ekv@stop#1##3{##1##3}
236 \long\def\ekvbreakPostSneak##1##2\ekv@stop#1##3{##3##1}
```

(End definition for \ekvbreak, \ekvbreakPreSneak, and \ekvbreakPostSneak. These functions are documented on page 5.)

\ekvsneak One last thing we want to do for \ekvset is to provide macros that just smuggle stuff \ekvsneakPre after \ekvset's effects.

```
237 \long\def\ekvsneak##1##2\ekv@stop#1##3%
       ##2\ekv@stop#1{##3##1}%
    }
241 \long\def\ekvsneakPre##1##2\ekv@stop#1##3%
242
       ##2\ekv@stop#1{##1##3}%
243
244
```

(End definition for \ekvsneak and \ekvsneakPre. These functions are documented on page 5.)

Additionally to the \ekvset macro we also want to provide an \ekvparse macro, that has the same scope as \keyval_parse:NNn from expl3. This is pretty analogue to the \ekvset implementation, we just put an \unexpanded here and there instead of other macros to stop the \expanded on our output.

```
245 \long\def\ekvparse##1##2##3%
                        {%
                           \ekv@parse##1##2\ekv@mark##3#1\ekv@stop#1%
                   (End definition for \ekvparse. This function is documented on page 4.)
      \ekv@parse
                   249 \long\def\ekv@parse##1##2##3#1%
                           \ekv@gobble@from@mark@to@stop##3\ekv@endparse\ekv@stop
                           \ekv@parse@other##1##2##3,\ekv@stop,%
                           \ekv@parse##1##2\ekv@mark
                   (End definition for \ekv@parse.)
   \ekv@endparse
                   255 \long\def\ekv@endparse
                           \ekv@stop\ekv@parse@other##1,\ekv@stop,\ekv@parse##2\ekv@mark
                   (End definition for \ekv@endparse.)
\ekv@parse@other
                   258 \long\def\ekv@parse@other##1##2##3,%
                   259
                           \ekv@gobble@from@mark@to@stop##3\ekv@endparse@other\ekv@stop
                           \ekv@eq@other##3\ekv@nil\ekv@mark\ekv@parse@eq@other@a
                   261
                             =\ekv@mark\ekv@parse@eq@active\ekv@stop
                           ##1##2%
                   263
                           \ekv@parse@other##1##2\ekv@mark
                   264
                   265
```

(End definition for \ekv@parse@other.)

```
\ekv@parse@eq@other@a
       \ekv@parse@eq@other@b
                                266 \long\def\ekv@parse@eq@other@a##1\ekv@stop
                                267
                                        \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@parse@eq@other@active@a
                                268
                                          #2\ekv@mark\ekv@parse@eq@other@b\ekv@stop
                                269
                                     }
                                   \long\def\ekv@parse@eq@other@b
                                        ##1\ekv@nil\ekv@mark\ekv@parse@eq@other@active@a\ekv@stop\ekv@mark
                                        \ekv@strip{##1}\ekv@parse@pair
                                274
                                      }
                                (End definition for \ekv@parse@eq@other@a and \ekv@parse@eq@other@b.)
\ekv@parse@eq@other@active@a
\ekv@parse@eq@other@active@b
                                276 \long\def\ekv@parse@eq@other@active@a##1\ekv@stop##2\ekv@nil\ekv@mark
                                        \ekv@parse@eq@other@active@b{##1}##2=%
                                278
                                279
                                280 \long\def\ekv@parse@eq@other@active@b##1%
                                281
                                        \ekv@strip{##1}\ekv@parse@pair
                                282
                                (\textit{End definition for } \verb|\ekv@parse@eq@other@active@a| and \verb|\ekv@parse@eq@other@active@b|)
        \ekv@parse@eq@active
       \ekv@parse@eq@active@
                                284 \long\def\ekv@parse@eq@active
                                        ##1\ekv@nil\ekv@mark\ekv@parse@eq@other@a\ekv@stop\ekv@mark
                                285
                                        \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@parse@eq@active@
                                287
                                          #2\ekv@mark\ekv@parse@noeq\ekv@stop
                                288
                                     }
                                289
                                   \long\def\ekv@parse@eq@active@##1\ekv@stop
                                291
                                        \ekv@strip{##1}\ekv@parse@pair
                                292
                                293
                                (End definition for \ekv@parse@eq@active and \ekv@parse@eq@active@.)
              \ekv@parse@noeq
                                   \long\def\ekv@parse@noeq
                                        ##1\ekv@nil\ekv@mark\ekv@parse@eq@active@\ekv@stop\ekv@mark
                                295
                                296
                                        \ekv@ifblank@##1\ekv@nil\ekv@ifempty@B\ekv@ifempty@true@F@gobbletwo
                                297
                                          \ekv@ifempty@A\ekv@ifempty@B\@firstofone
                                298
                                          {\ekv@strip{##1}\ekv@parse@key}%
                                299
                                      }
```

(End definition for \ekv@parse@noeq.)

```
\ekv@endparse@other
                        301 \long\def\ekv@endparse@other
                                \ekv@stop
                                \ekv@eq@other##1\ekv@nil\ekv@mark\ekv@parse@eq@other@a
                        303
                                =\ekv@mark\ekv@parse@eq@active\ekv@stop
                        304
                                \ekv@parse@other##3\ekv@mark
                        306
                              {}
                        307
                        (End definition for \ekv@endparse@other.)
    \ekv@parse@pair
   \ekv@parse@pair@
                        308 \long\def\ekv@parse@pair##1##2\ekv@nil
                                \ekv@strip{##2}\ekv@parse@pair@{##1}%
                           \label{longle} $$ \lceil ekv@parse@pair@#1##2##3##4\% $$
                                \unexpanded{##4{##2}{##1}}%
                        314
                        315
                        (End definition for \ekv@parse@pair and \ekv@parse@pair@.)
     \ekv@parse@key
                        316 \long\def\ekv@parse@key##1##2##3%
                                \displaystyle \sum_{\mu=0}^{\mu} \frac{1}{\mu}
                        318
                        319
                        (End definition for \ekv@parse@key.)
                             Finally really setting things up with \ekvset's temporary meaning:
                        320 }
                        321 \begingroup
                        322 \catcode'\,=13
                        323 \catcode'\==13
                        324 \ekvset,=
```

\ekvchangeset

Provide a macro that is able to switch out the current $\langle set \rangle$ in \ekvset. This operation is slow (by comparison, it should be slightly faster than \ekvsneak), but allows for something similar to pgfkeys's $\langle key \rangle$ /.cd mechanism. However this operation is more expensive than /.cd as we can't just redefine some token to reflect this, but have to switch out the set expandably, so this works similar to the \ekvsneak macros reading and reinserting the remainder of the $\langle key \rangle = \langle value \rangle$ list.

```
325 \def\ekvchangeset#1%
326 {%
327 \expandafter\ekv@changeset\csname\ekv@undefined@set{#1}\endcsname\ekv@mark
328 }
```

(End definition for \ekvchangeset. This function is documented on page 5.)

\ekv@changeset

This macro does the real change-out of \ekvchangeset. We introduced an \ekv@mark to not accidentally remove some braces which we have to remove again.

```
\label{eq:condition} $$ \log\left(\frac{e^{2\kappa}\left(\frac{1}{2\kappa}\right)}{{\kappa^{330}}} \left(\frac{\kappa^{2}}{2\kappa^{330}}\right)^{2\kappa} \left(\frac{1}{2\kappa^{330}}\right)^{2\kappa} \\ + \frac{1}{2\kappa^{330}} \left(\frac{1}{2\kappa^{330}}\right)^{2
```

(End definition for \ekv@changeset.)

\ekv@set@pair

\ekv@set@pair gets invoked with the space and brace stripped key-name as its first argument, the value as the second argument, and the set name as the third argument. It builds the key-macro name and provides everything to be able to throw meaningful error messages if it isn't defined. \ekv@set@pair@ will space and brace strip the value if the macro is defined and call the key-macro. Else it'll branch into the check whether an unknown key handler is defined for this set and that one will branch into the error messages provided by \ekv@set@pair if it isn't.

```
\long\def\ekv@set@pair#1#2\ekv@nil#3%
     {%
       \expandafter\ekv@set@pair@
335
          \csname
            \ifcsname #3{#1}\endcsname
              #3{#1}%
338
            \else
339
              relax%
            \fi
341
          \endcsname
342
          {#2}%
343
          {%
            \expandafter\ekv@set@pair@
              \csname
                 \ifcsname #3{}u\endcsname
                   #3{}u%
348
                 \else
                   relax%
350
                \fi
351
              \endcsname
352
              {#2}%
353
              {%
                 \ekv@ifdefined{#3{#1}N}%
355
                   \ekv@err@noarg
                   \ekv@err@unknown
                   #3%
              }%
              {#1}%
360
361
362
  \long\def\ekv@set@pair@#1#2%
363
364
       \int x#1\relax
365
         \ekv@fi@secondoftwo
366
       \fi
367
       \@firstoftwo
368
       {\ekv@strip{#2}#1}%
369
```

 $(End\ definition\ for\ \verb+\ekv@set@pair.)$

\ekv@set@key

Analogous to \ekv@set@pair, \ekv@set@key builds the NoVal key-macro and provides an error-branch. \ekv@set@key@ will test whether the key-macro is defined and if so call it, else the errors are thrown.

```
\long\def\ekv@set@key#1#2%
     {%
372
       \expandafter\ekv@set@key@
         \csname
            \ifcsname #2{#1}N\endcsname
375
              #2{#1}N%
376
            \else
377
              relax%
378
            \fi
379
         \endcsname
         {%
            \expandafter\ekv@set@key@
              \csname
                \ifcsname #2{}uN\endcsname
                  #2{}uN%
385
                \else
386
                  relax%
387
                \fi
              \endcsname
              {%
                \ekv@ifdefined{#2{#1}}%
                  \ekv@err@reqval
                  \ekv@err@unknown
                  #2%
394
395
              {#1}%
396
         }%
397
     }
398
   \def\ekv@set@key@#1%
399
       \ifx#1\relax
         \ekv@fi@secondoftwo
       \fi
       \@firstoftwo#1%
     }
405
```

(End definition for \ekv@set@key.)

\ekvsetdef Provide a macro to define a shorthand to use \ekvset on a specified \(set \). The first macro expands \ekvset twice, such that everything which can be done up to this point is done.

```
406 \protected\def\ekvsetdef#1#2%
    {%
       \expandafter\expandafter\expandafter
408
       \ekv@setdef\expandafter\expandafter\expandafter{\ekvset{#2}{##1}}#1%
```

(End definition for \ekvsetdef. This function is documented on page 3.)

\ekv@setdef

This auxiliary macro defines the shorthand macro after \ekvset got expanded as far as possible.

```
411 \protected\def\ekv@setdef#1#2%
412 {%
413 \long\def#2##1{#1}%
414 }

(End definition for \ekv@setdef.)
```

\ekv@err \ekv@err@ Since \ekvset is fully expandable as long as the code of the keys is (which is unlikely) we want to somehow throw expandable errors, in our case via undefined control sequences.

```
415 \begingroup
416 \edef\ekv@err
     {%
417
       \endgroup
418
       \unexpanded{\long\def\ekv@err}##1%
419
420
            \unexpanded{\expandafter\ekv@err@\@firstofone}%
421
            {\unexpanded\expandafter{\csname ! expkv Error:\endcsname}##1.}%
422
            \unexpanded{\ekv@stop}%
423
         }%
424
     }
425
426 \ekv@err
427 \def\ekv@err@{\expandafter\ekv@gobbleto@stop}
```

(End definition for \ekv@err and \ekv@err@.)

\ekv@err@common@ \ekv@err@unknown \ekv@err@noarg \ekv@err@reqval Now we can use \ekv@err to set up some error messages so that we can later use those instead of the full strings.

```
428 \long\def\ekv@err@common #1#2{\expandafter\ekv@err@common@\string#2{#1}}
429 \long\def\ekv@err@common@#1'#2' #3.#4#5{\ekv@err{#4 ('#5', set '#2')}}
430 \long\def\ekv@err@unknown#1#2{\ekv@err@common{unknown key}#1{#2}}
431 \long\def\ekv@err@noarg #1#2{\ekv@err@common{value forbidden}#1{#2}}
432 \long\def\ekv@err@reqval #1#2{\ekv@err@common{value required}#1{#2}}
```

(End definition for \ekv@err@common and others.)

\ekv@strip@a \ekv@strip@b \ekv@strip@c Finally we borrow some ideas of expl3's l3tl to strip spaces from keys and values. This \ekv@strip also strips one level of outer braces after stripping spaces, so an input of {abc} becomes abc after stripping. It should be used with #1 prefixed by \ekv@mark. Also this implementation at most strips one space from both sides.

```
433 \def\ekv@strip#1%
     {%
434
       \long\def\ekv@strip##1%
435
436
            \ekv@strip@a
437
              ##1%
438
              \ekv@nil
              \ekv@mark#1%
              #1\ekv@nil{}%
            \ekv@stop
         }%
443
       \long\def\ekv@strip@a##1\ekv@mark#1##2\ekv@nil##3%
444
445
```

```
\verb|\ekv@strip@b##3##1##2\\ekv@nil|
        }%
447
      \verb|\long\def\ekv@strip@b##1#1\ekv@nil| \\
448
         {%
          \ekv@strip@c##1\ekv@nil
450
451
       452
453
          ##3{##1}%
454
        }%
455
    }
_{457} \ekv@strip{ }
(End definition for \ekv@strip and others.)
    Now everything that's left is to reset the category code of @.
458 \catcode'\@=\ekv@tmp
```

Index

The italic numbers denote the pages where the corresponding entry is described, numbers underlined point to the definition, all others indicate the places where it is used.

E	\ekv@eq@other 177, 188, 229, 261, 303
\ekvbreak 5, <u>234</u>	\ekv@err 415, 429
\ekvbreakPostSneak 5, 234	\ekv@err@ 415
\ekvbreakPreSneak 5, 234	\ekv@err@common 428
\ekvchangeset 5, 325	\ekv@err@common@ 428
\ekvDate 4, 4, 8, <u>18</u>	\ekv@err@noarg 356, 428
\ekvdef	\ekv@err@reqval 392, 428
\ekvdefNoVal 2, <u>84</u>	\ekv@err@unknown 357, 393, 428
\ekvdefunknown $3, \frac{84}{}$	\ekv@fi@firstoftwo 23
$\ensuremath{\mbox{\sc kvdefunknownNoVal}}\ \dots \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $	\ekv@fi@gobble 23
\ekvifdefined 4, <u>82</u>	\ekv@fi@secondoftwo 23, 57, 366, 402
\ekvifdefinedNoVal	\ekv@gobble@from@mark@to@stop
\ekvifdefinedset 5, 156	
\ekvlet	\ekv@gobble@mark <u>23</u> , 331
\ekvletkv	\ekv@gobbleto@stop <u>23</u> , 427
\ekvletkvNoVal	\ekv@ifblank <u>47</u>
\ekvletNoVal	\ekv@ifblank@ <u>47</u> , 223, 297
\ekvparse	\ekv@ifdefined $\underline{53}$, 82, 83, 158, 355, 391
\ekvset	\ekv@ifempty <u>33</u> , 67, 73
\ekvsneak 5, 237	\ekv@ifempty@ <u>33</u> ,52
\ekvsneakPre 5, 237	\ekv@ifempty@A 35, 36,
\ekvVersion	38, 39, 40, 41, 42, 45, 50, 52, 224, 298 \ekv@ifempty@B 35, 36, 38, 39,
(onviolation	40, 41, 42, 45, 49, 50, 223, 224, 297, 298
N	\ekv@ifempty@false 33
\noexpand 153	\ekv@ifempty@true <u>33</u> , 49
T	\ekv@ifempty@true@F 33
T _E X and L ^A T _E X 2 $_{\mathcal{E}}$ commands:	\ekv@ifempty@true@F@gobble 33,223
\@firstofone 23, 41, 42, 45, 224, 298, 421	\ekv@ifempty@true@F@gobbletwo 33,297
\@firstoftwo <u>23</u> , 40, 59, 368, 404	\ekv@mark31,
\@gobble	32, 52, 165, 171, 174, 177, 179, 181,
$\c 23, 36, 39, 50, 78$	183, 188, 189, 191, 195, 196, 199, 203, 212, 214, 215, 221, 229, 230,
\ekv@changeset 327, 329	232, 247, 253, 256, 261, 262, 264,
\ekv@checkvalid 65,	268, 269, 272, 276, 285, 287, 288,
86, 94, 102, 110, 118, 127, 136, 144	295, 303, 304, 306, 327, 440, 444, 452
\ekv@defsetmacro 89,	\ekv@name 5, 61, 82, 83, 88, 96,
97, 105, 113, 122, 131, 139, 147, 150	104, 112, 120, 121, 129, 130, 138, 146 \ekv@name@key 5, 61, 153
\ekv@endparse 251, 255	\ekv@name@set
\ekv@endparse@other 260, 301	\ekv@nil 49, 188, 195, 199, 203, 212,
\ekv@endset 169, <u>173</u>	214, 221, 223, 229, 261, 268, 272,
\ekv@endset@other 187, 227	276, 285, 287, 295, 297, 303, 308,
\ekv@eq@active . $\frac{177}{195}$, 195, 214, 268, 287	333, 439, 441, 444, 446, 448, 450, 452

\ekv@parse 247, 249, 256	\ekv@set@key 225, 371
\ekv@parse@eq@active 262, 284, 304	\ekv@set@key@ 373, 382, 399
\ekv@parse@eq@active@ 284, 295	\ekv@set@noeq 215, 221
\ekv@parse@eq@other@a	\ekv@set@other
261, <u>266</u> , 285, 303	170, 174, <u>185,</u> 232, 329, 331
\ekv@parse@eq@other@active@a	\ekv@set@pair 201, 209, 219, 333
<u>268, 272, <u>276</u></u>	\ekv@set@pair@ 335, 345, 363
\ekv@parse@eq@other@active@b 276	\ekv@setdef 409, 411
\ekv@parse@eq@other@b <u>266</u>	\ekv@stop 30, 32, 165, 169, 170,
\ekv@parse@key 299, 316	174, 177, 179, 181, 183, 187, 189,
\ekv@parse@noeq 288, 294	193, 196, 199, 203, 212, 215, 217,
\ekv@parse@other 252, 256, 258, 306	221, 228, 230, 234, 235, 236, 237,
\ekv@parse@pair 274, 282, 292, 308	239, 241, 243, 247, 251, 252, 256,
\ekv@parse@pair@ 308	260, 262, 266, 269, 272, 276, 285, 288, 290, 295, 302, 304, 423, 442, 452
\ekv@set 164, 167, 174, 329, 331	\ekv@strip 201, 209, 219,
	225, 274, 282, 292, 299, 310, 369, 433
\ekv@set@eq@active 189, 211, 230 \ekv@set@eq@active@ 211, 221	\ekv@strip@a 433
\ekv@set@eq@other@a 188, 193, 212, 229	\ekv@strip@b 433
\ekv@set@eq@other@active@a	
	\ekv@strip@c 433
\ekv@set@eq@other@active@b 203	\ekv@tmp
	\ekv@undefined@set . 64, 89, 97, 105,
\ekv@set@eq@other@b 193	113, 122, 131, 139, 147, 158, 164, 327