

# Arbitrary Precision Numbers

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# 1 User's Documentation

This macro file `apnum.tex` implements addition, subtraction, multiplication, division, power to an integer and other calculation ( $\sqrt{x}$ ,  $e^x$ ,  $\ln x$ ,  $\sin x$ ,  $\arctan x$ , ...) with “large numbers” with arbitrary number of decimal digits. The numbers are in the form:

```
<sign><digits>.<digits>
```

where optional  $\langle sign \rangle$  is the sequence of + and/or -. The nonzero number is treated as negative if and only if there is odd number of - signs. The first part or second part of decimal  $\langle digits \rangle$  (but no both) can be empty. The decimal point is optional if second part of  $\langle digits \rangle$  is empty.

There can be unlimited number of digits in the operands. Only TeX main memory or your patience during calculation with very large numbers are your limits. Note, that the `apnum.tex` implementation includes a lot of optimization and it is above 100 times faster (on large numbers) than the implementation of the similar task in the package `fltpoint.sty`. And the `fp.sty` doesn't implements arbitrary number of digits. The extensive technical documentation can serve as an inspiration how to do TeX macro programming.

## 1.1 Evaluation of Expressions

After `\input apnum` in your document you can use the macro `\evaldef<sequence>\{<expression>\}`. It gives the possibility for comfortable calculation. The  $\langle expression \rangle$  can include numbers (in the form described above) combined by +, -, \*, / and ^ operators and by possible brackets () in an usual way. The result is stored to the  $\langle sequence \rangle$  as a “literal macro”. Examples:

```
\evaldef\A {2+4*(3+7)}
% ... the macro \A includes 42
\evaldef\B {\the\pageno * \A}
% ... the macro \B includes 84
\evaldef\C {123456789000123456789 * -123456789123456789123456789}
% ... \C includes -15241578765447341344197531849955953099750190521
\evaldef\D {1.23456789 + 12345678.9 - \A}
% ... the macro \D includes 12345596.13456789
```

```
\evaldef\X {1/3}
% ... the macro \X includes .333333333333333333
```

The limit of the number of digits of the division result can be set by `\apTOT` and `\apFRAC` registers. First one declares maximum calculated digits in total and second one declares maximum of digits after decimal point. The result is limited by both those registers. If the `\apTOT` is negative, then its absolute value is treated as a “soft limit”: all digits before decimal point are calculated even if this limit is exceeded. The digits after decimal point are not calculated when this limit is reached. The special value `\apTOT=0` means that the calculation is limited only by `\apFRAC`. Default values are `\apTOT=0` and `\apFRAC=20`.

The operator `^` means the powering, i.e.  $2^8$  is 256. The exponent have to be an integer (no decimal point is allowed) and a relatively small integer is assumed.

The scanner of the `\evaldef` macro reads (roughly speaking) the *<expression>* in the form “operand binary-operator operand binary-operator etc.” without expansion. The spaces are not significant in the *<expression>*. The operands are:

- numbers (in the format *<sign><digits>.<digits>*) or
- numbers in scientific notation (see the section 1.2) or
- sequences *<sign>\the<token>* or *<sign>\number<token>* or
- any other single *<token>* optionally preceded by *<sign>* and optionally followed by a sequence of parameters enclosed in braces, for example `\A` or `\B{<text>}` or `-\C{<textA>}{<textB>}`. This case has two meanings:
  - numeric constant defined in a “literal macro” (`\def\A{42}`, `\evaldef\A{13/15}`) or
  - “function-like” macro which returns a value after processing.

The `apnum.tex` macro file provides the following “function-like” macros allowed to use them as an operand in the *<expression>*:

- `\ABS {<value>}` for absolute value,
- `\SGN {<value>}` returns sign of the *<value>*,
- `\iDIV {<dividend>}{<divisor>}` for integer division,
- `\iMOD {<dividend>}{<divisor>}` for integer remainder,
- `\iROUND {<value>}` for rounding the number to the integer,
- `\iFRAC {<value>}` for fraction part of the `\iROUND`,
- `\FAC {<integer value>}` for factorial,
- `\BINOM {<integer above>}{<integer below>}` for binomial coefficient,
- `\SQRT {<value>}` for square root of the *<value>*,
- `\EXP {<value>}` applies exponential function to *<value>*,
- `\LN {<value>}` for natural logarithm of the *<value>*,
- `\SIN {<value>}`, `\COS {<value>}`, `\TAN {<value>}` for  $\sin x$ ,  $\cos x$  and  $\tan x$  functions,
- `\ASIN {<value>}`, `\ACOS {<value>}`, `\ATAN {<value>}` for  $\arcsin x$ ,  $\arccos x$  and  $\arctan x$  functions,
- `\PI`, `\PIhalf` for constants  $\pi$  and  $\pi/2$ .

The arguments of all these functions can be a nested *<expressions>* with the syntax like in the `\evaldef` macro. Example:

```
\def\A{20}
\evaldef\B{ 30*\SQRT{ 100 + 1.12*\the\widowpenalty } / (4-\A) }
```

Note that the arguments of the “function-like” macros are enclosed by normal T<sub>E</sub>X braces `{}` but the round brackets `()` are used for re-arranging of the common priority of the `+`, `-`, `*`, `/` and `^` operators. The macros `\SQRT`, `\EXP`, `\LN`, `\SIN`, `\COS`, `\TAN`, `\ASIN`, `\ACOS`, `\ATAN` use `\apTOT` and `\apFRAC` registers similar like during division.

The `\PI` and `\PIhalf` are “function-like” macros without parameters. They returns the constant with `\apFRAC` digits after decimal point.

Users can define their own “function-like” macros, see the section 1.3.

The output of `\evaldef\foo{<expression>}` processing is stored, of course, to the “literal macro” `\foo`. But there are another outputs like side effect of the processing:

- The `\OUT` macro includes exactly the same result as `\foo`.

- The `\apSIGN` register includes the value 1 or 0 or -1 respectively dependent on the fact that the output is positive, zero or negative.
- The `\apE` register is equal to the decimal exponent when scientific number format is used (see the next section 1.2).

For example, you can compare long numbers using `\apSIGN` register (where direct usage of `\ifnum` primitive may cause arithmetic overflow):

```
\TEST {123456789123456789} > {123456789123456788} \iftrue OK \else KO \fi
```

The `\TEST` macro is defined like:

```
\def\TEST#1#2#3#4{\evaldef\tmp{#1-(#3)}\ifnum\apSIGN #2 0 }
```

The `apnum.tex` macros do not provide the evaluation of the  $\langle expression \rangle$  at the expansion level only. There are two reasons. First, the macros can be used in classical  $\text{\TeX}$  only with Knuth's plain  $\text{\TeX}$  macro. No  $\text{\eTeX}$  is needed. And the expansion-only evaluation of any expression isn't possible in classical  $\text{\TeX}$ . Second reason is the speed optimization (see the section 1.4). Anyway, users needn't expansion-only evaluation. They can write `\evaldef\alpha{\langle expression \rangle} \edef\foo{... \alpha ...}` instead of `\edef\foo{... \langle expression \rangle ...}`. There is only one case when this “pre-processing” trick cannot be used: while expansion of the parameters of asynchronous `\write` commands. But you can save the  $\langle expression \rangle$  unexpanded into the file and you can read the file again in the second step and do `\evaldef` during reading the file.

## 1.2 Scientific Notation of Numbers

The macro `\evaldef` is able to operate with the numbers written in the notation:

```
<sign><digits>.<digits>E<sign><digits>
```

For example `1.234E9` means  $1.234 \cdot 10^9$ , i.e. 1234000000 or the text `1.234E-3` means .001234. The decimal exponent (after the E letter) have to be in the range  $\pm 2147483647$  because we store this value in normal  $\text{\TeX}$  register.

The `\evaldef\langle sequence \rangle\{\langle expression \rangle\}` operates by “normal way” if there are no operands with E syntax in the  $\langle expression \rangle$ . But if an operand is expressed in scientific form then `\evaldef` provide the calculation with the mantissa and the exponent separately. Only the mantissa of the result is found in the  $\langle sequence \rangle$  and `\OUT` macros. The exponent of the result is stored in the `\apE` register. You can define the macro which shows the complete result of the calculation, for example:

```
\def\showE#1{\message{#1\ifnum\apE=0 \else*10^{\the\apE}\fi}}
```

Suppose `\evaldef\foo{\langle expression \rangle}` is processed and the complete result is  $R = \text{\foo} * 10^{\text{\apE}}$ . There are two possibilities how to save such complete result  $R$  to the `\foo` macro: use `\apEadd\foo` or `\apEnum\foo`. Both macros do nothing if `\apE=0`. Else the `\apEadd\langle sequence \rangle` macro adds  $E\langle exponent \rangle$  to the  $\langle sequence \rangle$  macro and `\apEnum\langle sequence \rangle` moves the decimal point to the new right position in the  $\langle sequence \rangle$  macro or appends zeros. The `\apE` register is set to zero after the macro `\apEadd` or `\apEnum` is finished. Example:

```
\evaldef\foo{ 3 * 4E9 }      % \foo is 12, \apE=9
\apEadd\foo                  % \foo is 12E+9
\evaldef\foo{ 7E9 + 5E9 }    % \foo is 12, \apE=9
\apEnum\foo                  % \foo is 12000000000
```

There are another usable macros for operations with scientific numbers.

- `\apROLL \langle sequence \rangle\{\langle shift \rangle\}` ... the  $\langle sequence \rangle$  is assumed to be a macro with the number. The decimal point of this number is shifted right by  $\langle shift \rangle$  parameter, i.e. the result is multiplied by  $10^{\langle shift \rangle}$ . The  $\langle sequence \rangle$  is redefined by this result. For example the `\apEnum\A` does `\apROLL\A{\apE}`.
- `\apNORM \langle sequence \rangle\{\langle num \rangle\}` ... the  $\langle sequence \rangle$  is supposed to be a macro with  $\langle mantissa \rangle$  and it will be redefined. The number  $\langle mantissa \rangle * 10^{\text{\apE}}$  (with current value of the `\apE` register) is assumed. The new mantissa saved in the  $\langle sequence \rangle$  is the “normalized mantissa” of the same number. The `\apE` register is corrected so the “normalized mantissa”  $* 10^{\text{\apE}}$  gives the

same number. The  $\langle num \rangle$  parameter is the number of non-zero digits before the decimal point in the outputted mantissa. If the parameter  $\langle num \rangle$  starts by dot following by integer (for example  $\{.2\}$ ), then the outputted mantissa has  $\langle num \rangle$  digits after decimal point. For example `\def\A{1.234}\apE=0 \apNORM\A{.0}` defines  $\A$  as 1234 and  $\apE=-3$ .

- The `\apROUND  $\langle sequence \rangle \{ \langle num \rangle \}$`  rounds the number, which is included in the macro  $\langle sequence \rangle$  and redefines  $\langle sequence \rangle$  as rounded number. The digits after decimal point at the position greater than  $\langle num \rangle$  are ignored in the rounded number. The decimal point is removed, if it is the right most character in the `\OUT`. The ignored part is saved to the `\XOUT` macro without trailing right zeros.

Examples of `\apROUND` usage:

```
\def\A{12.3456}\apROUND\A{1}    % \A is "12.3",    \XOUT is "456"
\def\A{12.3456}\apROUND\A{9}    % \A is "12.3456", \XOUT is empty
\def\A{12.3456}\apROUND\A{0}    % \A is "12",      \XOUT is "3456"
\def\A{12.0000}\apROUND\A{0}    % \A is "12",      \XOUT is empty
\def\A{12.0001}\apROUND\A{2}    % \A is "12",      \XOUT is "01"
\def\A{.000010}\apROUND\A{2}    % \A is "0",        \XOUT is "001"
\def\A{-12.3456}\apROUND\A{2}   % \A is "-12.34",  \XOUT is "56"
\def\A{12.3456}\apROUND\A{-1}   % \A is "10",      \XOUT is "23456"
\def\A{12.3456}\apROUND\A{-4}   % \A is "0",        \XOUT is "00123456"
```

The following example saves the result of the `\evaldef` in scientific notation with the mantissa with maximal three digits after decimal point and one digit before.

```
\evaldef\X{...}\apNORM\X{1}\apROUND\X{3}\apEadd\X
```

The macros `\apEadd`, `\apEnum`, `\apROLL`, `\apNORM` and `\apROUND` redefine the macro  $\langle sequence \rangle$  given as their first argument. They are not “function-like” macros and they cannot be used in an  $\langle expression \rangle$ . The macro  $\langle sequence \rangle$  must be the number in the format  $\langle simple sign \rangle \langle digits \rangle . \langle digits \rangle$  where  $\langle simple sign \rangle$  is one minus or none and the rest of number has the format described in the first paragraph of this documentation. The scientific notation isn't allowed here. This format of numbers is in accordance with the output of the `\evaldef` macro.

### 1.3 Notes for macro programmers

If you plan to create a “function-like” macro which can be used as an operand in the  $\langle expression \rangle$  then observe that first token in the macro body must be `\relax`. This tells to the  $\langle expression \rangle$  scanner that the calculation follows. The result of this calculation must be saved into the `\OUT` macro and into the `\apSIGN` and `\apE` registers.

Example. The `\ABS` macro for the absolute value is defined by:

```
704: \def\ABS#1{\relax           % mandatory \relax for "function-like" macros
705:   \evaldef\OUT{#1}%        % evaluation of the input parameter
706:   \ifnum\apSIGN<0          % if (input < 0)
707:     \apSIGN=1               % sign = 1
708:     \apREfirst\OUT          % remove first "minus" from OUT
709:   \fi                       % fi
```

apnum.tex

```
Usage: \evaldef\A{ 2 - \ABS{3-10} }% \A includes -5.
```

Note, that `\apSIGN` register is corrected by final routine of `\evaldef` according the `\OUT` value. But setting `\apSIGN` in your macro is recommended because user can use your macro directly outside of `\evaldef`.

The `\evaldef\foo{ $\langle expression \rangle$ }` is processed in two steps. The  $\langle expression \rangle$  scanner converts the input to the macro call of the `\apPLUS`, `\apMINUS`, `\apMUL`, `\apDIV` or `\apPOW` macros with two parameters. They do addition, subtraction, multiplication, division and power to the integer. These macros are processed in the second step. For example:

```
\evaldef\A{ 2 - 3*8 }      converts the input to:
\apMINUS{2}{\apMUL{3}{8}} and this is processed in the second step.
```

The macros `\apPLUS`, `\apMINUS`, `\apMUL`, `\apDIV` and `\apPOW` behave like normal “function-like” macros with one important exception: they don't accept general  $\langle expression \rangle$  in their parameters, only single operand (see section 1.1) is allowed.

If your calculation is processed in the loop very intensively than it is better to save time of such calculation and to avoid the  $\langle expression \rangle$  scanner processing (first step of the `\evaldef`). So, it is recommended to use directly the Polish notation of the expression as shown in the second line in the example above. See section 2.10 for more inspirations.

The output of the `\apPLUS`, `\apMINUS`, `\apMUL`, `\apDIV` and `\apPOW` macros is stored in `\OUT` macro and the registers `\apSIGN` and `\apE` are set accordingly.

The number of digits calculated by `\apDIV` macro is limited by the `\apTOT` and `\apFRAC` registers as described in the section 1.1. There is another result of `\apDIV` calculation stored in the `\XOUT` macro. It is the remainder of the division. Example:

```
\apTOT=0 \apFRAC=0 \apDIV{1234567892345}{2}\ifnum\XOUT=0 even \else odd\fi
```

You cannot apply `\ifodd` primitive on “large numbers” directly because the numbers may be too big.

If you set something locally inside your “function-like” macro, then such data are accessible only when your macro is called outside of `\evaldef`. Each parameter and the whole `\evaldef` is processed inside a  $\TeX$  group, so your locally set data are inaccessible when your macro is used inside another “function-like” parameter or inside `\evaldef`. The `\XOUT` output is set locally by `\apDIV` macro, so it serves as a good example of this feature:

```
{\apDIV{1}{3}          ... \XOUT is .00000000000000000001 }
... \XOUT is undefined
\evaldef{1/3}          ... \XOUT is undefined
\apPLUS{1}{\apDIV{1}{3}} ... \XOUT is undefined
```

The macro `\apPOW{\langle base \rangle}{\langle exponent \rangle}` calculates the power to the integer exponent. A slight optimization is implemented here so the usage of `\apPOW` is faster than repeated multiplication. The decimal non-integer exponents are not allowed. Use `\EXP` and `\LN` macros if you need to calculate non-integer exponent:

```
\def\POWER#1#2{\relax \EXP{(#2)*\LN{#1}}}
```

Note that both parameters are expected as an  $\langle expression \rangle$ . Thus the #2 is surrounded in the rounded brackets.

Examples of another common “function-like” macros:

```
\evaldef\degcoef{PI/180}
\def\SINdeg#1{\relax \SIN{\degcoef*(#1)}}
\def\COSdeg#1{\relax \COS{\degcoef*(#1)}}
\def\SINH#1{\relax \evaldef\myE{\EXP{#1}}\evaldef\OUT{(\myE-1/\myE)/2}}
\def\ASINH#1{\relax \LN{#1+\SQRT{(#1)^2+1}}}
\def\LOG#1{\relax \apLNtenexec \apDIV{\LN{#1}}{\apLNten}}
```

In another example, we implement the field  $\mathbb{F}\{\langle index \rangle\}$  as an “function-like” macro. User can set values by `\set\mathbb{F}\{\langle index \rangle\}=\{\langle value \rangle\}` and then these values can be used in an  $\langle expression \rangle$ .

```
\def\set#1#2#3#4{\evaldef\index{#2}\evaldef\value{#4}%
  \expandafter\edef\csname \string#1[\index]\endcsname{\value}}
\def\F#1{\relax % function-like macro
  \evaldef\index{#1}%
  \expandafter\ifx\csname \string\F[\index]\endcsname\relax
    \def\OUT{0}% undefined value
  \else
    \edef\OUT{\csname \string\F[\index]\endcsname}%
  \fi
}
\set \mathbb{F}{12/2} = {28+13}
\set \mathbb{F}{2*4} = {144^2}
\evaldef\test { 1 + \mathbb{F}{6} } \message{result=\test}
```



As an exercise, you can implement linear interpolation of known values.

The final example shows, how to implement the macro `\usedimen{⟨dimen⟩}{⟨unit⟩}`. It is “function-like” macro, it can be used in the  $\langle expression \rangle$  and it returns the  $\langle decimal number \rangle$  with the property  $\langle dimen \rangle = \langle decimal number \rangle \langle unit \rangle$ .

```
\def\usedimen #1#2{\relax % function-like macro
  \def\OUT{0}% % default value, if the unit isn't known
  \csname dimenX#2\endcsname{#1}}
\def\dimenXpt #1{\apDIV{\number#1}{65536}}
\def\dimenXcm #1{\apDIV{\number#1}{1864682.7}}
\def\dimenXmm #1{\apDIV{\number#1}{186468.27}}
%... etc.
\evaldef\af{\usedimen{\hsize}{cm}} % \a includes 15.91997501773358008845
```

Note that user cannot write `\usedimen\hsize{cm}` without braces because this isn't the syntactically correct operand (see section 1.1) and the  $\langle expression \rangle$  scanner is unable to read it.

## 1.4 Experiments

The following table shows the time needed for calculation of randomly selected examples. The comparison with the package `fltpoint.sty` is shown. The symbol  $\infty$  means that it is out of my patience.

input	# of digits in the result	time spent by <code>apnum.tex</code>	time spent by <code>fltpoint.sty</code>
200!	375	0.33 sec	173 sec
1000!	2568	29 sec	$\infty$
$5^{17^2}$	203	0.1 sec	81 sec
$5^{17^3}$	3435	2.1 sec	$\infty$
1/17	1000	0.13 sec	113 sec
1/17	100000	142 sec	$\infty$

## 2 The Implementation

### 2.1 Name Convention, Version, Counters

The internal control sequence names typically used in `apnum.tex` have the form `\apNAMEsuffix`, but there are exceptions. The control sequences mentioned in the section 1.1 (user's documentation) have typically more natural names. And the internal counter registers have names `\apnumA`, `\apnumB`, `\apnumC` etc.

The code starts by the greeting. The `\apVERSION` includes the version of this software.

```
7: \def\apVERSION{1.4a <Dec 2015>}
8: \message{The Arbitrary Precision Numbers, \apVERSION}
```

`apnum.tex`

We declare auxiliary counters and one Boolean variable.

```
12: \newcount\apnumA \newcount\apnumB \newcount\apnumC \newcount\apnumD
13: \newcount\apnumE \newcount\apnumF \newcount\apnumG \newcount\apnumH
14: \newcount\apnumO \newcount\apnumP \newcount\apnumL
15: \newcount\apnumX \newcount\apnumY \newcount\apnumZ
16: \newcount\apSIGNa \newcount\apSIGNb \newcount\apEa \newcount\apEb
17: \newif\ifapX
```

`apnum.tex`

The counters `\apSIGN`, `\apE`, `\apTOT` and `\apFRAC` are declared here:

```
19: \newcount\apSIGN
20: \newcount\apE
21: \newcount\apTOT \apTOT=0
22: \newcount\apFRAC \apFRAC=20
```

`apnum.tex`

---

`\apVERSION`: 7      `\apSIGN`: 4, 5–8, 11–12, 14–18, 21–23, 27–28, 32, 36–37, 39–40, 43–46  
`\apE`: 4, 5–9, 11–12, 14, 17–18, 21–23, 27–28, 32, 36–38, 40      `\apTOT`: 3, 6–7, 22, 36, 41–45  
`\apFRAC`: 3, 6–7, 22, 36–45

Somebody sometimes sets the @ character to the special catcode. But we need to be sure that there is normal catcode of the @ character.

```
24: \apnumZ=\catcode'\@ \catcode'\@=12
```

apnum.tex

## 2.2

## Evaluation of the Expression

Suppose the following expression  $A+B*(C+D)+E$  as an example.

The main task of the `\evaldef`  $x\{A+B*(C+D)+E\}$  is to prepare the macro `\tmpb` with the content (in this example) `\apPLUS{\apPLUS{A}{\apMUL{B}{\apPLUS{C}{D}}}}{E}` and to execute the `\tmpb` macro.

The expression scanner adds the `\limits` at the end of the expression and reads from left to right the couples “operand, operator”. For our example: `A+`, `B*`, `C+`, `D+` and `E\limits`. The `\limits` operator has the priority 0, plus, minus have priority 1, `*` and `/` have priority 2 and `^` has priority 3. The brackets are ignored, but each occurrence of the opening bracket `(` increases priority by 4 and each occurrence of closing bracket `)` decreases priority by 4. The scanner puts each couple including its current priority to the stack and does a test at the top of the stack. The top of the stack is executed if the priority of the top operator is less or equal the previous operator priority. For our example the stack is only pushed without execution until `D+` occurs. Our example in the stack looks like:

<code>\D + 1</code>	<code>1&lt;=5</code>	<code>exec:</code>	
<code>\C + 5</code>		<code>{\C+\D} + 1</code>	<code>1&lt;=2</code> <code>exec:</code>
<code>\B * 2</code>		<code>\B * 2</code>	<code>{\B*{\C+\D}} + 1</code> <code>1&lt;=1</code> <code>exec:</code>
<code>\A + 1</code>		<code>\A + 1</code>	<code>\A + 1</code> <code>{\A+{\B*{\C+\D}}}</code> <code>+ 1</code>
bottom	0	bottom	0

Now, the priority on the top is greater, then scanner pushes next couple and does the test on the top of the stack again.

<code>\E \limits 0</code>	<code>0&lt;=1</code>	<code>exec:</code>	
<code>{\A+{\B*{\C+\D}}}</code>	<code>+ 1</code>	<code>{\A+{\B*{\C+\D}}}+\E</code>	<code>\end 0</code> <code>0&lt;=0</code> <code>exec:</code>
bottom	0	bottom	0
			RESULT

Let  $p_t, p_p$  are the priority on the top and the previous priority in the stack. Let  $v_t, v_p$  are operands on the top and in the previous line in the stack, and the same notation is used for operators  $o_t$  and  $o_p$ . If  $p_t \leq p_p$  then: pop the stack twice, create composed operand  $v_n = v_p o_p v_t$  and push  $v_n, o_t, p_t$ . Else push new couple “operand, operator” from the expression scanner. In both cases try to execute the top of the stack again. If the bottom of the stack is reached then the last operand is the result.

The `\evaldef` macro is protected by `\relax`. It means that it can be used inside an  $\langle expression \rangle$  as a “function-like” macro, but I don’t imagine any usual application of this. The `\apEVALa` is executed.

```
28: \def\evaldef{\relax \apEVALa}
```

apnum.tex

The macro `\apEVALa`  $\langle sequence \rangle \{ \langle expression \rangle \}$  runs the evaluation of the expression in the group. The base priority is initialized by `\apnumA=0`, then `\apEVALb`  $\langle expression \rangle \backslash limits$  scans the expression and saves the result in the form `\apPLUS{A}{\apMUL{B}{C}}` (etc.) into the `\tmpb` macro. This macro is executed. The group is finished by `\apEND` macro, which keeps the `\OUT`, `\apSIGN` and `\ape` values unchanged. Finally the defined  $\langle sequence \rangle$  is set equivalent to the `\OUT` macro.

```
29: \def\apEVALa#1#2{\begingroup \apnumA=0 \apnumE=1 \apEVALb#2\limits \tmpb \apEND \let#1=\OUT}
```

apnum.tex

The scanner is in one of the two states: reading operand or reading operator. The first state is initialized by `\apEVALb` which follows to the `\apEVALc`. The `\apEVALc` reads one token and switches by its value. If the value is a `+` or `-` sign, it is assumed to be the part of the operand prefix. Plus sign is ignored (and `\apEVALc` is run again), minus signs are accumulated into `\tmpa`.

The auxiliary macro `\apEVALd` runs the following tokens to the `\fi`, but first it closes the conditional and skips the rest of the macro `\apEVALc`.

---

`\evaldef`: 2, 3–6, 8, 11, 27, 35–37, 39–40, 43–45    `\apEVALa`: 8, 10    `\OUT`: 3, 4–6, 8, 11–12, 14–24,  
26–29, 32, 34–46    `\apEVALb`: 8–10    `\apEVALc`: 8–9    `\apEVALd`: 9



```

30: \def\apEVALb{\def\tmpa{}\apEVALc}
31: \def\apEVALc#1{%
32:   \ifx+#1\apEVALd \apEVALc \fi
33:   \ifx-#1\edef\tmpa{\tmpa-}\apEVALd\apEVALc \fi
34:   \ifx(#1\apEVALd \apEVALe \fi
35:   \ifx\the#1\apEVALd \apEVALf\the\fi
36:   \ifx\number#1\apEVALd \apEVALf\number\fi
37:   \apTESTdigit#1\iftrue
38:     \ifx E#1\let\tmpb=\tmpa \expandafter\apEVALd\expandafter\apEVALk
39:     \else \edef\tmpb{\tmpa#1}\expandafter\apEVALd\expandafter\apEVALn\fi\fi
40:   \edef\tmpb{\tmpa\noexpand#1}\expandafter
41:   \futurelet\expandafter\apNext\expandafter\apEVALg\romannumeral-'\.%
42: }
43: \def\apEVALd#1\fi#2-'\.{\fi#1}

```

If the next token is opening bracket, then the global priority is increased by 4 using the macro `\apEVALe`. Moreover, if the sign before bracket generates the negative result, then the new multiplication (by  $-1$ ) is added using `\apEVALp` to the operand stack.

```

44: \def\apEVALe{%
45:   \ifx\tmpa\empty \else \ifnum\tmpa1<0 \def\tmpb{-1}\apEVALp \apMUL 4\fi\fi
46:   \advance\apnumA by4
47:   \apEVALb
48: }

```

If the next token is `\the` or `\number` primitives (see lines 35 and 36), then one following token is assumed as  $\TeX$  register and these two tokens are interpreted as an operand. This is done by `\apEVALf`. The operand is packed to the `\tmpb` macro.

```

49: \def\apEVALf#1#2{\expandafter\def\expandafter\tmpb\expandafter{\tmpa#1#2}\apEVALo}

```

If the next token is not a number (the `\apTESTdigit#1\iftrue` results like `\iffalse` at line 37) then we save the sign plus this token to the `\tmpb` at line 41 and we do check of the following token by `\futurelet`. The `\apEVALg` is processed after that. The test is performed here if the following token is open brace (a macro with parameter). If this is true then this parameter is appended to `\tmpb` by `\apEVALh` and the test about the existence of second parameter in braces is repeated by next `\futurelet`. The result of this loop is stored into `\tmpb` macro which includes  $\langle sign \rangle$  followed by  $\langle token \rangle$  followed by all parameters in braces. This is considered as an operand.

```

50: \def\apEVALg{\ifx\apNext \bgroup \expandafter\apEVALh \else \expandafter\apEVALo \fi}
51: \def\apEVALh#1{\expandafter\def\expandafter\tmpb\expandafter{\tmpb#1}}\expandafter

```

If the next token after the sign is a digit or a dot (tested in `\apEVALc` by `\apTESTdigit` at line 37), then there are two cases. The number includes the E symbol as a first symbol (this is allowed in scientific notation, mantissa is assumed to equal to one). The `\apEVALk` is executed in such case. Else the `\apEVALn` starts the reading the number.

The first case with E letter in the number is solved by macros `\apEVALk` and `\apEVALm`. The number after E is read by `\apE=` and this number is appended to the `\tmpb` and the expression scanner skips to `\apEVALo`.

```

53: \def\apEVALk{\afterassignment\apEVALm\apE=}
54: \def\apEVALm{\edef\tmpb{\tmpb E\the\apE}\apEVALo}

```

The second case (there is normal number) is processed by the macro `\apEVALn`. This macro reads digits (token per token) and saves them to the `\tmpb`. If the next token isn't digit nor dot then the second state of the scanner (reading an operator) is invoked by running `\apEVALo`. If the E is found then the exponent is read to `\apE` and it is processed by `\apEVALm`.

```

55: \def\apEVALn#1{\apTESTdigit#1%
56:   \iftrue \ifx E#1\afterassignment\apEVALm\expandafter\expandafter\expandafter\apE
57:   \else\edef\tmpb{\tmpb#1}\expandafter\expandafter\expandafter\apEVALn\fi
58:   \else \expandafter\apEVALo\expandafter#1\fi

```

---

```

\apEVALe: 9   \apEVALf: 9   \apEVALg: 9   \apEVALh: 9   \apEVALk: 9   \apEVALm: 9
\apEVALn: 9

```

59: }

The reading an operator is done by the `\apEVALo` macro. This is more simple because the operator is only one token. Depending on this token the macro `\apEVALp`  $\langle operator \rangle \langle priority \rangle$  pushes to the stack (by the macro `\apEVALpush`) the value from `\tmpb`, the  $\langle operator \rangle$  and the priority increased by `\apnumA` (level of brackets).

If there is a problem (level of brackets less than zero, level of brackets not equal to zero at the end of the expression, unknown operator) we print an error using `\apEVALerror` macro.

The `\apNext` is set to `\apEVALb`, i.e. scanner returns back to the state of reading the operand. But exceptions exist: if the `)` is found then priority is decreased and the macro `\apEVALo` is executed again. If the end of the  $\langle expression \rangle$  is found then the loop is ended by `\let\apNext=\relax`.

```

60: \def\apEVALo#1{\let\apNext=\apEVALb
61:   \ifx+#1\apEVALp \apPLUS 1\fi
62:   \ifx-#1\apEVALp \apMINUS 1\fi
63:   \ifx*#1\apEVALp \apMUL 2\fi
64:   \ifx/#1\apEVALp \apDIV 2\fi
65:   \ifx^#1\apEVALp \apPOWx 3\fi
66:   \ifx)#1\advance\apnumA by-4 \let\apNext=\apEVALo \let\tmpa=\relax
67:   \ifnum\apnumA<0 \apEVALerror{many brackets "}"\fi
68:   \fi
69:   \ifx\limits#1%
70:     \ifnum\apnumA>0 \apEVALerror{missing bracket "}"\let\tmpa=\relax
71:     \else \apEVALp\END 0\let\apNext=\relax \fi
72:   \fi
73:   \ifx\tmpa\relax \else \apEVALerror{unknown operator "\string#1"}\fi
74:   \apnumE=0 \apNext
75: }
76: \def\apEVALp#1#2{%
77:   \apnumB=#2 \advance\apnumB by\apnumA
78:   \toks0=\expandafter{\expandafter{\tmpb}{#1}}%
79:   \expandafter\apEVALpush\the\toks0\expandafter{\the\apnumB}% {value}{op}{priority}
80:   \let\tmpa=\relax
81: }
```

apnum.tex

The `\apEVALstack` macro includes the stack, three items  $\{\langle operand \rangle\}\{\langle operator \rangle\}\{\langle priority \rangle\}$  per level. Left part of the macro contents is the top of the stack. The stack is initialized with empty operand and operator and with priority zero. The dot here is only the “total bottom” of the stack.

82: \def\apEVALstack{{}{0}.}

apnum.tex

The macro `\apEVALpush`  $\{\langle operand \rangle\}\{\langle operator \rangle\}\{\langle priority \rangle\}$  pushes its parameters to the stack and runs `\apEVALdo`  $\langle whole stack \rangle @$  to do the desired work on the top of the stack.

```

83: \def\apEVALpush#1#2#3{% value, operator, priority
84:   \toks0={{#1}{#2}{#3}}%
85:   \expandafter\def\expandafter\apEVALstack\expandafter{\the\toks0\apEVALstack}%
86:   \expandafter\apEVALdo\apEVALstack@%
87: }
```

apnum.tex

Finally, the macro `\apEVALdo`  $\{\langle vt \rangle\}\{\langle ot \rangle\}\{\langle pt \rangle\}\{\langle vp \rangle\}\{\langle op \rangle\}\{\langle pp \rangle\}\langle rest of the stack \rangle @$  performs the execution described at the beginning of this section. The new operand  $\langle vn \rangle$  is created as  $\langle op \rangle \langle vp \rangle \langle vt \rangle$ , this means `\apPLUS`  $\{\langle vp \rangle\}\{\langle vt \rangle\}$  for example. The operand is not executed now, only the result is composed by the normal T<sub>E</sub>X notation. If the bottom of the stack is reached then the result is saved to the `\tmpb` macro. This macro is executed after group by the `\apEVALa` macro.

```

88: \def\apEVALdo#1#2#3#4#5#6#7@{%
89:   \apnumB=#3 \ifx#2\apPOWx \advance\apnumB by1 \fi
90:   \ifnum\apnumB>#6\else
91:     \ifnum#6=0 \def\tmpb{#1}%\toks0={#1}\message{RESULT: \the\toks0}
92:     \ifnum\apnumE=1 \def\tmpb{\apPn{#1}}\fi
93:   \else \def\apEVALstack{#7}\apEVALpush{#5{#4}{#1}}{#2}{#3}%
94:   \fi\fi
```

apnum.tex

---

`\apEVALo`: 9–10    `\apEVALp`: 9–10    `\apEVALstack`: 10    `\apEVALpush`: 10    `\apEVALdo`: 10

```
95: }
```

The macro `\apEVALerror`  $\langle string \rangle$  prints an error message. We decide to be better to print only `\message`, no `\errmessage`. The `\tmpb` is prepared to create `\OUT` as ?? and the `\apNext` macro is set in order to skip the rest of the scanned  $\langle expression \rangle$ .

```
96: \def\apEVALerror#1{\message{\noexpand\evaldef ERROR: #1.}%
97:   \def\OUT{0}\apE=0\apSIGN=0\def\apNext##1\apEND{\apEND}%
98: }
```

apnum.tex

The auxiliary macro `\apTESTdigit`  $\langle token \rangle$  `\iftrue` tests, if the given token is digit, dot or E letter.

```
99: \def\apTESTdigit#1#2{%
100:   \ifx E#1\apXtrue \else
101:     \ifcat.\noexpand#1%
102:       \ifx.#1\apXtrue \else
103:         \ifnum'#1<'0 \apXfalse\else
104:           \ifnum'#1>'9 \apXfalse\else \apXtrue\fi
105:       \fi\fi
106:     \else \apXfalse
107:   \fi\fi
108:   \ifapX
109: }
```

apnum.tex

### 2.3 Preparation of the Parameter

All operands of `\apPLUS`, `\apMINUS`, `\apMUL`, `\apDIV` and `\apPOW` macros are preprocessed by `\apPPa` macro. This macro solves (roughly speaking) the following tasks:

- It partially expands (by `\expandafter`) the parameter while  $\langle sign \rangle$  is read.
- The  $\langle sign \rangle$  is removed from parameter and the appropriate `\apSIGN` value is set.
- If the next token after  $\langle sign \rangle$  is `\relax` then the rest of the parameter is executed in the group and the results `\OUT`, `\apSIGN` and `\apE` are used.
- Else the number is read and saved to the parameter.
- If the read number has the scientific notation  $\langle mantissa \rangle E \langle exponent \rangle$  then only  $\langle mantissa \rangle$  is saved to the parameter and `\apE` is set as  $\langle exponent \rangle$ . Else `\apE` is zero.

The macro `\apPPa`  $\langle sequence \rangle \langle parameter \rangle$  calls `\apPPb`  $\langle parameter \rangle @ \langle sequence \rangle$  and starts reading the  $\langle parameter \rangle$ . The result will be stored to the  $\langle sequence \rangle$ .

Each token from  $\langle sign \rangle$  is processed by three `\expandafters` (because there could be `\csname...\endcsname`). It means that the parameter is partially expanded when  $\langle sign \rangle$  is read. The `\apPPb` macro sets the initial value of `\tmpc` and `\apSIGN` and executes the macro `\apPPc`  $\langle parameter \rangle @ \langle sequence \rangle$ .

```
113: \def\apPPa#1#2{\expandafter\apPPb#2@#1}
114: \def\apPPb{\def\tmpc{}\apSIGN=1 \apE=0 \apXfalse \expandafter\expandafter\expandafter\apPPc}
115: \def\apPPc#1{%
116:   \ifx+#1\apPPd \fi
117:   \ifx-#1\apSIGN=-\apSIGN \apPPd \fi
118:   \ifx\relax#1\apPPe \fi
119:   \apPPg#1%
120: }
121: \def\apPPd#1\apPPg#2{\fi\expandafter\expandafter\expandafter\apPPc}
```

apnum.tex

The `\apPPc` reads one token from  $\langle sign \rangle$  and it is called recursively while there are + or - signs. If the read token is + or - then the `\apPPd` closes conditionals and executes `\apPPc` again.

If `\relax` is found then the rest of parameter is executed by the `\apPPe`. The macro ends by `\apPPf`  $\langle result \rangle @$  and this macro reverses the sign if the result is negative and removes the minus sign from the front of the parameter.

---

`\apEVALerror`: 10–11    `\apTESTdigit`: 9, 11    `\apPPa`: 11–12    `\apPPb`: 11–12    `\apPPc`: 11  
`\apPPd`: 11    `\apPPe`: 11–12    `\apPPf`: 12

```

122: \def\apPpE#1\apPpG#2#3@{\fi\apXtrue
123:   \begingroup#3% execution of the parameter in the group
124:   \edef\tmpb{\apE=\the\apE\relax\noexpand\apPpF\OUT@}\expandafter\endgroup\tmpb
125: }
126: \def\apPpF#1{\ifx-#1\apSIGN=-\apSIGN \expandafter\apPpG\else\expandafter\apPpG\expandafter#1\fi}

```

The `\apPpG`  $\langle parameter \rangle @$  macro is called when the  $\langle sign \rangle$  was processed and removed from the input stream. The main reason of this macro is to remove trailing zeros from the left and to check, if there is the zero value written for example in the form 0000.000. When this macro is started then `\tmpc` is empty. This is a flag for removing trailing zeros. They are simply ignored before decimal point. The `\apPpG` is called again by `\apPpH` macro which removes the rest of `\apPpG` macro and closes the conditional. If the decimal point is found then next zeros are accumulated to the `\tmpc`. If the end of the parameter `@` is found and we are in the “removing zeros state” then the whole value is assumed to be zero and this is processed by `\apPpI @`. If another digit is found (say 2) then there are two situations: if the `\tmpc` is non-empty, then the digit is appended to the `\tmpc` and the `\apPpI`  $\langle expanded tmp \rangle$  is processed (say `\apPpI .002`) followed by the rest of the parameter. Else the digit itself is stored to the `\tmpc` and it is returned back to the input stream (say `\apPpI 2`) followed by the rest of the parameter.

```

127: \def\apPpG#1{%
128:   \ifx.#1\def\tmpc{.}\apPpH\fi
129:   \ifx\tmpc\empty\else\edef\tmpc{\tmpc#1}\fi
130:   \ifx0#1\apPpH\fi
131:   \ifx\tmpc\empty\edef\tmpc{#1}\fi
132:   \ifx@#1\def\tmpc{@}\apSIGN=0 \fi
133:   \expandafter\apPpI\tmpc
134: }
135: \def\apPpH#1\apPpI\tmpc{\fi\apPpG}

```

The macro `\apPpI`  $\langle parameter without trailing zeros \rangle @ \langle sequence \rangle$  switches to two cases: if the execution of the parameter was processed then the `\OUT` doesn't include E notation and we can simply define  $\langle sequence \rangle$  as the  $\langle parameter \rangle$  by the `\apPpJ` macro. This saves the copying of the (possible) long result to the input stream again.

If the executing of the parameter was not performed, then we need to test the existence of the E notation of the number by the `\apPpK` macro. We need to put the  $\langle parameter \rangle$  to the input stream and to use `\apPpL` to test these cases. We need to remove unwanted E letter by the `\apPpM` macro.

```

136: \def\apPpI{\ifapX \expandafter\apPpJ \else \expandafter\apPpK \fi}
137: \def\apPpJ#1@#2{\def#2{#1}}
138: \def\apPpK#1@#2{\ifx@#1@ \apSIGN=0 \def#2{0}\else \apPpL#1E@#2\fi}
139: \def\apPpL#1E#2@#3{%
140:   \ifx@#1@ \def#3{1}\else \def#3{#1}\fi
141:   \ifx@#2@ \else \afterassignment\apPpM \apE=#2\fi
142: }
143: \def\apPpM E{}

```

The `\apPpN`  $\langle param \rangle$  macro does the same as `\apPpA\OUT{\langle param \rangle}`, but the minus sign is returned back to the `\OUT` macro if the result is negative.

```

144: \def\apPpN#1{\expandafter\apPpB#1@\OUT}

```

The `\apPpAb`  $\langle macro \rangle \{ \langle paramA \rangle \} \{ \langle paramB \rangle \}$  is used for parameters of all macros `\apPLUS`, `\apMUL` etc. It prepares the  $\langle paramA \rangle$  to `\tmpa`,  $\langle paramB \rangle$  to `\tmpb`, the sign and  $\langle decimal exponent \rangle$  of  $\langle paramA \rangle$  to the `\apSIGNa` and `\apEa`, the same of  $\langle paramB \rangle$  to the `\apSIGNb` and `\apEb`. Finally, it executes the  $\langle macro \rangle$ .

```

148: \def\apPpAb#1#2#3{%
149:   \expandafter\apPpB#2@\tmpa \apSIGNa=\apSIGN \apEa=\apE
150:   \expandafter\apPpB#3@\tmpb \apSIGNb=\apSIGN \apEb=\apE
151:   #1%
152: }

```

---

`\apPpG`: 11–12    `\apPpH`: 12    `\apPpI`: 12    `\apPpJ`: 12    `\apPpK`: 12    `\apPpL`: 12    `\apPpM`: 12  
`\apPpN`: 10, 12    `\apPpAb`: 12–14, 17, 22, 27–28, 32

The `\apPPs`  $\langle macro \rangle \langle sequence \rangle \{ \langle param \rangle \}$  prepares parameters for `\apROLL`, `\apROUND` and `\apNORM` macros. It saves the  $\langle param \rangle$  to the `\tmpc` macro, expands the  $\langle sequence \rangle$  and runs the macro `\apPpt`  $\langle macro \rangle \langle expanded sequence \rangle . @ \langle sequence \rangle$ . The macro `\apPpt` reads first token from the  $\langle expanded sequence \rangle$  to `#2`. If `#2` is minus sign, then `\apnumG=-1`. Else `\apnumG=1`. Finally the  $\langle macro \rangle \langle expanded sequence \rangle . @ \langle sequence \rangle$  is executed (but without the minus sign in the input stream). If `#2` is zero then `\apPPu`  $\langle macro \rangle \langle rest \rangle . @ \langle sequence \rangle$  is executed. If the  $\langle rest \rangle$  is empty, (i.e. the parameter is simply zero) then  $\langle macro \rangle$  isn't executed because there is nothing to do with zero number as a parameter of `\apROLL`, `\apROUND` or `\apNORM` macros.

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```

153: \def\apPPs#1#2#3{\def\tmpc{#3}\expandafter\apPpt\expandafter#1#2.@#2}
154: \def\apPpt#1#2{%
155:   \ifx-#2\apnumG=-1 \def\apNext{#1}%
156:   \else \ifx0#2\apnumG=0 \def\apNext{\apPPu#1}\else \apnumG=1 \def\apNext{#1#2}\fi\fi
157:   \apNext
158: }
159: \def\apPPu#1#2.@#3{\ifx0#2\apnumG=0 \ifx#1\apROUNDa\def\XOUT{}\fi
160:   \else\def\apNext{\apPpt#1#2.@#3}\expandafter\apNext\fi
161: }
```

## 2.4 Addition and Subtraction

The significant part of the optimization in `\apPLUS`, `\apMUL`, `\apDIV` and `\apPOW` macros is the fact, that we don't treat with single decimal digits but with their quartets. This means that we are using the numeral system with the base 10000 and we calculate four decimal digits in one elementary operation. The base was chosen  $10^4$  because the multiplication of such numbers gives results less than  $10^8$  and the maximal number in the  $\text{\TeX}$  register is about  $2 \cdot 10^9$ . We'll use the word "Digit" (with capitalized D) in this documentation if this means the digit in the numeral system with base 10000, i.e. one Digit is four digits. Note that for addition we can use the numeral system with the base  $10^8$  but we don't do it, because the auxiliary macros `\apIV*` for numeral system of the base  $10^4$  are already prepared.

Suppose the following example (the spaces between Digits are here only for more clarity).

```

123 4567 8901 9999      \apnumA=12 \apnumE=3 \apnumD=16
+          22.423      \apnumB=0  \apnumF=2 \apnumC=12
-----
sum in reversed order and without transmissions:
      {4230}{10021}{8901}{4567}{123}  \apnumD=-4
sum in normal order including transmissions:
123 4567 8902 0021.423
```

In the first pass, we put the number with more or equal Digits before decimal point above the second number. There are three Digits more in the example. The `\apnumC` register saves this information (multiplied by 4). The first pass creates the sum in reversed order without transmissions between Digits. It simply copies the `\apnumC/4` Digits from the first number to the result in reversed order. Then it does the sums of Digits without transmissions. The `\apnumD` is a relative position of the decimal point to the edge of the calculated number.

The second pass reads the result of the first pass, calculates transmissions and saves the result in normal order.

The first Digit of the operands cannot include four digits. The number of digits in the first Digit is saved in `\apnumE` (for first operand) and in `\apnumF` (for second one). The rule is to have the decimal point between Digits in all circumstances.

The `\apPLUS` and `\apMINUS` macros prepare parameters using `\apPPab` and execute `\apPLUSa`:

apnum.tex

```

165: \def\apPLUS{\relax \apPPab\apPLUSa}
166: \def\apMINUS#1#2{\relax \apPPab\apPLUSa{#1}{-#2}}
```

The macro `\apPLUSa` does the following work:

---

```

\apPPs: 13, 17, 29, 31–32  \apPpt: 13  \apPPu: 13  \apPLUS: 5, 6, 8, 10–13, 37–42, 44–46
\apMINUS: 5, 6, 10–11, 13, 44, 46  \apPLUSa: 13–14
```

- It gets the operands in `\tmpa` and `\tmpb` macros using the `\apPPab`.
- If the scientific notation is used and the decimal exponents `\apEa` and `\apEb` are not equal then the decimal point of one operand have to be shifted (by the macro `\apPLUSxE` at line 168).
- The digits before decimal point are calculated for both operands by the `\apDIG` macro. The first result is saved to `\apnumA` and the second result is saved to `\apnumB`. The `\apDIG` macro removes decimal point (if exists) from the parameters (lines 169 and 170).
- The number of digits in the first Digit is calculated by `\apIVmod` for both operands. This number is saved to `\apnumE` and `\apnumF`. This number is subtracted from `\apnumA` and `\apnumB`, so these registers now includes multiply of four (lines 171 and 172).
- The `\apnumC` includes the difference of Digits before the decimal point (multiplied by four) of given operands (line 173).
- If the first operand is negative then the minus sign is inserted to the `\apPLUSxA` macro else this macro is empty. The same for the second operand and for the macro `\apPLUSxB` is done (lines 174 and 175).
- If both operands are positive, then the sign of the result `\apSIGN` is set to one. If both operands are negative, then the sign is set to  $-1$ . But in both cases mentioned above we will do (internally) addition, so the macros `\apPLUSxA` and `\apPLUSxB` are set to empty. If one operand is negative and second positive then we will do subtraction. The `\apSIGN` register is set to zero and it will set to the right value later (lines 176 to 178).
- The macro `\apPLUSb<first op><first dig><second op><second dig><first Dig>` does the calculation of the first pass. The `<first op>` has to have more or equal Digits before decimal point than `<second op>`. This is reason why this macro is called in two variants dependent on the value `\apnumC`. The macros `\apPLUSxA` and `\apPLUSxB` (with the sign of the operands) are exchanged (by the `\apPLUSg`) if the operands are exchanged (lines 179 to 180).
- The `\apnumG` is set by the macro `\apPLUSb` to the sign of the first nonzero Digit. It is equal to zero if there are only zero Digits after first pass. The result is zero in such case and we do nothing more (line 182).
- The transmission calculation is different for addition and subtraction. If the subtraction is processed then the sign of the result is set (using the value `\apnumG`) and the `\apPLUSm` for transmissions is prepared. Else the `\apPLUSp` for transmissions is prepared as the `\apNext` macro (line 183).
- The result of the first pass is expanded in the input stream and the `\apNext` (i.e. transmissions calculation) is activated at line 184.
- if the result is in the form `.000123`, then the decimal point and the trailing zeros have to be inserted. Else the trailing zeros from the left side of the result have to be removed by `\apPLUSy`. This macro adds the sign of the result too (lines 185 to 191)

apnum.tex

```

167: \def\apPLUSa{%
168:   \ifnum\apEa=\apEb \apE=\apEa \else \apPLUSxE \fi
169:   \apDIG\tmpa\relax \apnumA=\apnumD % digits before decimal point
170:   \apDIG\tmpb\relax \apnumB=\apnumD
171:   \apIVmod \apnumA \apnumE \advance\apnumA by-\apnumE % digits in the first Digit
172:   \apIVmod \apnumB \apnumF \advance\apnumB by-\apnumF
173:   \apnumC=\apnumB \advance\apnumC by-\apnumA % difference between Digits
174:   \ifnum\apSIGNa<0 \def\apPLUSxA{-}\else \def\apPLUSxA{}\fi
175:   \ifnum\apSIGNb<0 \def\apPLUSxB{-}\else \def\apPLUSxB{}\fi
176:   \apSIGN=0 % \apSIGN=0 means that we are doing subtraction
177:   \ifx\apPLUSxA\empty \ifx\apPLUSxB\empty \apSIGN=1 \fi\fi
178:   \if\apPLUSxA\relax \if\apPLUSxB\relax \apSIGN=-1 \def\apPLUSxA{}\def\apPLUSxB{}\fi\fi
179:   \ifnum\apnumC>0 \apPLUSg \apPLUSb \tmpb\apnumF \tmpa\apnumE \apnumB % first pass
180:   \else \apnumC=-\apnumC \apPLUSb \tmpa\apnumE \tmpb\apnumF \apnumA
181:   \fi
182:   \ifnum\apnumG=0 \def\OUT{0}\apSIGN=0 \apE=0 \else
183:     \ifnum\apSIGN=0 \apSIGN=\apnumG \let\apNext=\apPLUSm \else \let\apNext=\apPLUSp \fi
184:     \apnumX=0 \edef\OUT{\expandafter}\expandafter \apNext \OUT@% second pass
185:     \ifnum\apnumD<1 % result in the form .000123

```

---

`\apPLUSxA`: 14–16    `\apPLUSxB`: 14–16



```

186:         \apnumZ=-\apnumD
187:         \def\tmpa{.}%
188:         \ifnum\apnumZ>0 \apADDzeros\tmpa \fi % adding dot and left zeros
189:         \edef\OUT{\ifnum\apSIGN<0-\fi\tmpa\OUT}%
190:     \else
191:         \edef\OUT{\expandafter}\expandafter\apPLUSy \OUT@% removing left zeros
192:     \fi\fi
193: }

```

The macro `\apPLUSb`  $\langle first\ op \rangle \langle first\ dig \rangle \langle second\ op \rangle \langle second\ dig \rangle \langle first\ Dig \rangle$  starts the first pass. The  $\langle first\ op \rangle$  is the first operand (which have more or equal Digits before decimal point). The  $\langle first\ dig \rangle$  is the number of digits in the first Digit in the first operand. The  $\langle second\ op \rangle$  is the second operand and the  $\langle second\ dig \rangle$  is the number of digits in the first Digit of the second operand. The  $\langle first\ Dig \rangle$  is the number of Digits before decimal point of the first operand, but without the first Digit and multiplied by 4.

The macro `\apPLUSb` saves the second operand to `\tmpd` and appends the 4 – *(second dig)* empty parameters before this operand in order to read desired number of digits to the first Digit of this operand. The macro `\apPLUSb` saves the first operand to the input queue after `\apPLUSc` macro. It inserts the appropriate number of empty parameters (in `\tmpc`) before this operand in order to read the right number of digits in the first attempt. It appends the `\apNL` marks to the end in order to recognize the end of the input stream. These macros expands simply to zero but we can test the end of input stream by `\ifx`.

The macro `\apPLUSb` calculates the number of digits before decimal point (rounded up to multiply by 4) in `\apnumD` by advancing `\first DIG` by 4. It initializes `\apnumZ` to zero. If the first nonzero Digit will be found then `\apnumZ` will be set to this Digit in the `\apPLUSc` macro.

```

194: \def\apPLUSb#1#2#3#4#5{%
195:   \edef\tmpd{\ifcase#4\or{}{}\or{}{}\or{}{}\fi#3}%
196:   \edef\tmpc{\ifcase#2\or{}{}\or{}{}\or{}{}\fi}%
197:   \let\apNext=\apPLUSc \apnumD=#5\advance\apnumD by4 \apnumG=0 \apnumZ=0 \def\OUT{}%
198:   \expandafter\expandafter\expandafter\apPLUSc\expandafter\tmpc#1\apNL\apNL\apNL\apNL@%
199: }

```

The macro `\apPLUSc` is called repeatedly. It reads one Digit from input stream and saves it to the `\apnumY`. Then it calls the `\apPLUSe`, which reads (if it is allowed, i.e. if `\apnumC<=0`) one digit from second operand `\tmpd` by the `\apIVread` macro. Then it does the addition of these digits and saves the result into the `\OUT` macro in reverse order.

Note, that the sign `\apPLUSxA` is used when `\apnumY` is read and the sign `\apPLUSxB` is used when advancing is performed. This means that we are doing addition or subtraction here.

If the first nonzero Digit is reached, then the macro `\apPLUSh` sets the sign of the result to the `\apnumG` and (maybe) exchanges the `\apPLUSxA` and `\apPLUSxB` macros (by the `\apPLUSg` macro) in order to the internal result of the subtraction will be always non-negative.

If the end of input stream is reached, then `\apNext` (used at line 212) is reset from its original value `\apPLUSc` to the `\apPLUSd` where the `\apnumY` is simply set to zero. The reading from input stream is finished. This occurs when there are more Digits after decimal point in the second operand than in the first one. If the end of input stream is reached and the `\tmpd` macro is empty (all data from second operand was read too) then the `\apPLUSf` macro removes the rest of input stream and the first pass of the calculation is done.

```

200: \def\apPLUSc#1#2#3#4{\apnumY=\apPLUSxA#1#2#3#4\relax
201:   \ifx\apNL#4\let\apNext=\apPLUSd\fi
202:   \ifx\apNL#1\relax \ifx\tmpd\empty \expandafter\expandafter\expandafter\apPLUSf \fi\fi
203:   \apPLUSE
204: }
205: \def\apPLUSd{\apnumY=0 \ifx\tmpd\empty \expandafter\apPLUSf \else\expandafter \apPLUSE\fi}
206: \def\apPLUSE{%
207:   \ifnum\apnumC>0 \advance\apnumC by-4
208:   \else \apIVread\tmpd \advance\apnumY by\apPLUSxB\apnumX \fi
209:   \ifnum\apnumZ=0 \apPLUSH \fi
210:   \edef\OUT{{\the\apnumY}\OUT}%

```

$\backslash\text{apPLUSb: } 14\text{--}15$       $\backslash\text{apPLUSc: } 15$       $\backslash\text{apPLUSe: } 15$       $\backslash\text{apPLUSh: } 15\text{--}16$       $\backslash\text{apPLUSg: } 14, 16$   
 $\backslash\text{apPLUSd: } 15$       $\backslash\text{apPLUSf: } 15\text{--}16$

```

211: \advance\apnumD by-4
212: \apNext
213: }
214: \def\apPLUSf#1@{
215: \def\apPLUSg{\let\tmpc=\apPLUSxA \let\apPLUSxA=\apPLUSxB \let\apPLUSxB=\tmpc}
216: \def\apPLUSh{\apnumZ=\apnumY

```

Why there is a complication about reading one parameter from input stream but second one from the macro `\tmpd`? This is more faster than to save both parameters to the macros and using `\apIVread` for both because the `\apIVread` must redefine its parameter. You can examine that this parameter is very long.

The `\apPLUSm`  $\langle data \rangle$  macro does transmissions calculation when subtracting. The  $\langle data \rangle$  from first pass is expanded in the input stream. The `\apPLUSm` macro reads repeatedly one Digit from the  $\langle data \rangle$  until the stop mark is reached. The Digits are in the range  $-9999$  to  $9999$ . If the Digit is negative then we need to add 10000 and set the transmission value `\apnumX` to one, else `\apnumX` is zero. When the next Digit is processed then the calculated transmission value is subtracted. The macro `\apPLUSw` writes the result for each Digit `\apnumA` in the normal (human readable) order.

```

219: \def\apPLUSm#1{%
220: \ifx@#1\else
221: \apnumA=#1 \advance\apnumA by-\apnumX
222: \ifnum\apnumA<0 \advance\apnumA by\apIVbase \apnumX=1 \else \apnumX=0 \fi
223: \apPLUSw
224: \expandafter\apPLUSm
225: \fi
226: }

```

apnum.tex

The `\apPLUSp`  $\langle data \rangle$  macro does transmissions calculation when addition is processed. It is very similar to `\apPLUSm`, but Digits are in the range 0 to 19998. If the Digit value is greater then 9999 then we need to subtract 10000 and set the transmission value `\apnumX` to one, else `\apnumX` is zero.

```

227: \def\apPLUSp#1{%
228: \ifx@#1\ifnum\apnumX>0 \apnumA=1 \apPLUSw \fi %.5+.5=.1 bug fixed
229: \else
230: \apnumA=\apnumX \advance\apnumA by#1
231: \ifnum\apnumA<\apIVbase \apnumX=0 \else \apnumX=1 \advance\apnumA by-\apIVbase \fi
232: \apPLUSw
233: \expandafter\apPLUSp
234: \fi
235: }

```

apnum.tex

The `\apPLUSw` writes the result with one Digit (saved in `\apnumA`) to the `\OUT` macro. The `\OUT` is initialized as empty. If it is empty (it means we are after decimal point), then we need to write all four digits by `\apIVwrite` macro (including left zeros) but we need to remove right zeros by `\apREMzerosR`. If the decimal point is reached, then it is saved to the `\OUT`. But if the previous `\OUT` is empty (it means there are no digits after decimal point or all such digits are zero) then `\def\OUT{\empty}` ensures that the `\OUT` is non-empty and the ignoring of right zeros are disabled from now.

```

236: \def\apPLUSw{%
237: \ifnum\apnumD=0 \ifx\OUT\empty \def\OUT{\empty}\else \edef\OUT{.\OUT}\fi \fi
238: \advance\apnumD by4
239: \ifx\OUT\empty \edef\tmpa{\apIVwrite\apnumA}\edef\OUT{\apREMzerosR\tmpa}%
240: \else \edef\OUT{\apIVwrite\apnumA\OUT}\fi
241: }

```

apnum.tex

The macro `\apPLUSy`  $\langle expanded OUT \rangle$  removes left trailing zeros from the `\OUT` macro and saves the possible minus sign by the `\apPLUSz` macro.

```

242: \def\apPLUSy#1{\ifx0#1\expandafter\apPLUSy\else \expandafter\apPLUSz\expandafter#1\fi}
243: \def\apPLUSz#1@{\edef\OUT{\ifnum\apSIGN<0-\fi#1}}

```

apnum.tex

---

`\apPLUSm`: 14, 16    `\apPLUSp`: 14, 16    `\apPLUSw`: 16    `\apPLUSy`: 14–16    `\apPLUSz`: 16

The macro `\apPLUSxE` uses the `\apROLLa` in order to shift the decimal point of the operand. We need to set the same decimal exponent in scientific notation before the addition or subtraction is processed.

```

244: \def\apPLUSxE{%
245:   \apnumE=\apEa \advance\apnumE by-\apEb
246:   \ifnum\apEa>\apEb \apPPs\apROLLa\tmpb{-\apnumE}\apE=\apEa
247:   \else \apPPs\apROLLa\tmpa{\apnumE}\apE=\apEb \fi
248: }

```

apnum.tex

## 2.5 Multiplication

Suppose the following multiplication example:  $1234 \times 567 = 699678$ .

Normal format:	Mirrored format:
1 2 3 4 *	4 3 2 1 *
5 6 7	7 6 5
-----	-----
*7: 7 14 21 28	*7: 28 21 14 7
*6: 6 12 18 24	*6: 24 18 12 6
*5: 5 10 15 20	*5: 20 15 10 5
-----	-----
6 9 9 6 7 8	8 7 6 9 9 6

This example is in numeral system of base 10 only for simplification, the macros work really with base 10000. Because we have to do the transmissions between Digit positions from right to left in the normal format and because it is more natural for  $\text{\TeX}$  to put the data into the input stream and read it sequentially from left to right, we use the mirrored format in our macros.

The macro `\apMUL` prepares parameters using `\apPPab` and executes `\apMULa`

```

252: \def\apMUL{\relax \apPPab\apMULa}

```

apnum.tex

The macro `\apMULa` does the following:

- It gets the parameters in `\tmpa` and `\tmpb` preprocessed using the `\apPPab` macro.
- It evaluates the exponent of ten `\apE` which is usable when the scientific notation of numbers is used (line 254).
- It calculates `\apSIGN` of the result (line 255).
- If `\apSIGN=0` then the result is zero and we will do nothing more (line 256).
- The decimal point is removed from the parameters by `\apDIG<param><register>`. The `\apnumD` includes the number of digits before decimal point (after the `\apDIG` is used) and the `<register>` includes the number of digits in the rest. The `\apnumA` or `\apnumB` includes total number of digits in the parameters `\tmpa` or `\tmpb` respectively. The `\apnumD` is re-calculated: it saves the number of digits after decimal point in the result (lines 257 to 259).
- Let  $A$  is the number of total digits in the `<param>` and let  $F = A \bmod 4$ , but if  $F = 0$  then reassign it to  $F = 4$ . Then  $F$  means the number of digits in the first Digit. This calculation is done by `\apIVmod<A><F>` macro. All another Digits will have four digits. The `\apMULb<param>@@@` is able to read four digits, next four digits etc. We need to insert appropriate number of empty parameters before the `<param>`. For example `\apMULb{}{}{}<param>@@@` reads first only one digit from `<param>`, next four digits etc. The appropriate number of empty parameters are prepared in the `\tmpc` macro (lines 260 to 261).
- The `\apMULb` reads the `<paramA>` (all Digits) and prepares the `\OUT` macro in the special interleaved format (described below). The format is finished by `*`. in the line 263.
- Analogical work is done with the second parameter `<paramB>`. But this parameter is processed by `\apMULc`, which reads Digits of the parameter and inserts them to the `\tmpa` in the reversed order (lines 264 to 266).
- The main calculation is done by `\apMULd<paramB>@`, which reads Digits from `<paramB>` (in reversed order) and does multiplication of the `<paramA>` (saved in the `\OUT`) by these Digits (line 267).

`\apPLUSxE`: 14, 17    `\apMUL`: 5, 6, 8–13, 17, 25, 34, 36–37, 39–40, 42–46    `\apMULa`: 17–18, 27

- The `\apMULg` macro converts the result `\OUT` to the human readable form (line 268).
- The possible minus sign and the trailing zeros of results of the type .00123 is prepared by `\apADDzeros\tpa` to the `\tpa` macro. This macro is appended to the result in the `\OUT` macro (lines 269 to 271).

apnum.tex

We need to read the two data streams when the multiplication of the  $\langle paramA \rangle$  by one Digit from  $\langle paramB \rangle$  is performed and the partial sum is actualized. First: the digits of the  $\langle paramA \rangle$  and second: the partial sum. We can save these streams to two macros and read one piece of information from such macros at each step, but this is not effective because the whole stream has to be read and redefined at each step. For  $\text{\TeX}$  is more natural to put one data stream to the input queue and to read pieces of information thereof. Thus we interleave both data streams into one `\OUT` in such a way that one element of data from first stream is followed by one element from second stream and it is followed by second element from first stream etc. Suppose that we are at the end of  $i - th$  line of the multiplication scheme where we have the partial sums  $s_n, s_{n-1}, \dots, s_0$  and the Digits of  $\langle paramA \rangle$  are  $d_k, d_{k-1}, \dots, d_0$ . The zero index belongs to the most right position in the mirrored format. The data will be prepared in the form:

$$\cdot \{s_n\} \{s_{(n-1)}\} \dots \{s_{(k+1)}\} * \{s_k\} \{d_{(k-1)}\} \dots \{s_1\} \{d_1\} \{s_0\} \{d_0\} *$$

For our example (there is a simplification: numeral system of base 10 is used and no transmissions are processed), after second line (multiplication by 6 and calculation of partial sums) we have in `\OUT`:

$$, \{28\} * \{45\} \{4\} \{32\} \{3\} \{19\} \{2\} \{6\} \{1\} *$$

and we need to create the following line during calculation of next line of multiplication scheme:

$$\cdot \{28\} \{45\} * \{5*4+32\} \{4\} \{5*3+19\} \{3\} \{5*2+6\} \{2\} \{5*1\} \{1\} *$$

This special format of data includes two parts. After the starting dot, there is a sequence of sums which are definitely calculated. This sequence is ended by first \* mark. The last definitely calculated sum follows this mark. Then the partial sums with the Digits of  $\langle paramA \rangle$  are interleaved and the data are finalized by second \*. If the calculation processes the the second part of the data then the general task is to read two data elements (partial sum and the Digit) and to write two data elements (the new partial sum and the previous Digit). The line calculation starts by copying of the first part of data until the first \* and appending the first data element after \*. Then the \* is written and the middle processing described above is started.

The macro `\apMULb`  $\langle paramA \rangle$  prepares the special format of the macro `\OUT` described above where the partial sums are zero. It means:

$$* \quad , \quad \{d_k\} \quad 0 \quad \{d_{(k-1)}\} \quad 0 \quad \dots \quad 0 \quad \{d_0\} \quad *$$

**\apMULb:** 17–19, 27–28

where  $d_i$  are Digits of  $\langle paramA \rangle$  in reversed order.

The first “sum” is only dot. It will be moved before  $*$  during the first line processing. Why there is such special “pseudo-sum”? The `\apMULe` with the parameter delimited by the first  $*$  is used in the context `\apMULe.{\sum}*` during the third line processing and the dot here protects from removing the braces around the first real sum.

```
274: \def\apMULb#1#2#3#4{\ifx@#4\else
275:   \ifx\OUT\empty \edef\OUT{{#1#2#3#4}*}\else\edef\OUT{{#1#2#3#4}0\OUT}\fi
276:   \expandafter\apMULb\fi
277: }
```

The macro `\apMULc`  $\langle paramB \rangle @@@@$  reads Digits from  $\langle paramB \rangle$  and saves them in reversed order into `\tmpa`. Each Digit is enclosed by  $\TeX$  braces  $\{\}$ .

```
278: \def\apMULc#1#2#3#4{\ifx@#4\else \edef\tmpa{{#1#2#3#4}\tmpa}\expandafter\apMULc\fi}
```

The macro `\apMULd`  $\langle paramB \rangle @$  reads the Digits from  $\langle paramB \rangle$  (in reversed order), uses them as a coefficient for multiplication stored in `\tmpnumA` and processes the `\apMULe`  $\langle special data format \rangle$  for each such coefficient. This corresponds with one line in the multiplication scheme.

```
279: \def\apMULd#1{\ifx@#1\else
280:   \apnumA=#1 \expandafter\apMULe \OUT
281:   \expandafter\apMULd
282:   \fi
283: }
```

The macro `\apMULE`  $\langle special data format \rangle$  copies the first part of data format to the `\OUT`, copies the next element after first  $*$ , appends  $*$  and does the calculation by `\apMULf`. The `\apMULf` is recursively called. It reads the Digit to `#1` and the partial sum to the `#2` and writes  $\{\apnumA*#1+#2\}$  to the `\OUT` (lines 295 to 299). If we are at the end of data, then `#2` is  $*$  and we write the  $\{\apnumA*#1\}$  followed by ending  $*$  to the `\OUT` (lines 288 to 290).

```
284: \def\apMULE#1#2{\apnumX=0 \def\OUT{{#1#2}*}\def\apOUTl{}\apnum0=1 \apnumL=0 \apMULf}
285: \def\apMULf#1#2{%
286:   \advance\apnum0 by-1 \ifnum\apnum0=0 \apOUTx \fi
287:   \apnumB=#1 \multiply\apnumB by\apnumA \advance\apnumB by\apnumX
288:   \ifx#2%
289:     \ifnum\apnumB<\apIVbase
290:       \edef\OUT{\OUT\expandafter\apOUTs\apOUTl.,\ifnum\the\apnumB#1=0 \else{\the\apnumB}\{#1\}\fi*}%
291:       \else \apIVtrans
292:         \expandafter \edef\csname apOUT:\apOUTn\endcsname
293:           {\csname apOUT:\apOUTn\endcsname{\the\apnumB}\{#1\}}%
294:         \apMULf0*\fi
295:     \else \advance\apnumB by#2
296:       \ifnum\apnumB<\apIVbase \apnumX=0 \else \apIVtrans \fi
297:       \expandafter
298:         \edef\csname apOUT:\apOUTn\endcsname{\csname apOUT:\apOUTn\endcsname{\the\apnumB}\{#1\}}%
299:       \expandafter\apMULf \fi
300: }
```

There are several complications in the algorithm described above.

- The result isn't saved directly to the `\OUT` macro, but partially into the macros `\apOUT:` $\langle num \rangle$ , as described in the section 2.9 where the `\apOUTx` macro is defined.
- The transmissions between Digit positions are calculated. First, the transmission value `\apnumX` is set to zero in the `\apMULE`. Then this value is subtracted from the calculated value `\apnumB` and the new transmission is calculated using the `\apIVtrans` macro if `\apnumB  $\geq$  10000`. This macro modifies `\apnumB` in order it is right Digit in our numeral system.
- If the last digit has nonzero transmission, then the calculation isn't finished, but the new pair  $\{\langle transmission \rangle\}\{0\}$  is added to the `\OUT`. This is done by recursively call of `\apMULf` at line 294.
- The another situation can be occurred: the last pair has both values zeros. Then we needn't to write this zero to the output. This is solved by the test `\ifnum\the\apnumB#1=0` at line 290.

---

`\apMULc`: 17–19    `\apMULd`: 17–19, 28    `\apMULE`: 19, 29    `\apMULf`: 19, 29

The macro `\apMULg` (*special data format*)@ removes the first dot (it is the #1 parameter) and prepares the `\OUT` to writing the result in reverse order, i.e. in human readable form. The next work is done by `\apMULh` and `\apMULi` macros. The `\apMULh` repeatedly reads the first part of the special data format (Digits of the result are here) until the first \* is found. The output is stored by `\apMULo` $\langle digits \rangle \{ \langle data \rangle \}$  macro. If the first \* is found then the `\apMULi` macro repeatedly reads the triple  $\{ \langle Digit \text{ of result} \rangle \} \{ \langle Digit \text{ of } A \rangle \} \{ \langle next \text{ Digit of result} \rangle \}$  and saves the first element in full (four-digits) form by the `\apIVwrite` if the third element isn't the stop-mark \*. Else the last Digit (first Digit in the human readable form) is saved by `\the`, because we needn't the trailing zeros here. The third element is put back to the input stream but it is ignored by `\apMULj` macro when the process is finished.

```

301: \def\apMULg#1{\def\OUT{}\apMULh}
302: \def\apMULh#1{\ifx#1\expandafter\apMULi
303:   \else \apnumA=#1 \apMULo4{\apIVwrite\apnumA}%
304:   \expandafter\apMULh
305:   \fi
306: }
307: \def\apMULi#1#2#3{\apnumA=#1
308:   \ifx#3\apMULo{\apNUMdigits\tmpa}\the\apnumA\expandafter\apMULj
309:   \else \apMULo4{\apIVwrite\apnumA}\expandafter\apMULi
310:   \fi{#3}%
311: }
312: \def\apMULj#1{}

```

apnum.tex

The `\apMULo`  $\langle digits \rangle \{ \langle data \rangle \}$  appends  $\langle data \rangle$  to the `\OUT` macro. The number of digits after decimal point `\apnumD` is decreased by the number of actually printed digits  $\langle digits \rangle$ . If the decimal point is to be printed into  $\langle data \rangle$  then it is performed by the `\apMULT` macro.

```

313: \def\apMULo#1#2{\edef\tmpa{#2}%
314:   \advance\apnumD by-#1
315:   \ifnum\apnumD<1 \ifnum\apnumD>-4 \apMULT\fi\fi
316:   \edef\OUT{\tmpa\OUT}%
317: }
318: \def\apMULT{\edef\tmpa{\apIVdot{-\apnumD}\tmpa}\edef\tmpa{\tmpa}}

```

apnum.tex

## 2.6 Division

Suppose the following example:

$\langle paramA \rangle$	:	$\langle paramB \rangle$		$\langle output \rangle$
		12345:678 = [12:6=2]		2 (2->1)
2*678		-1356		
		-1215 <0 correction!		1
		12345		
1*678		-678		
		5565 [55:6=8]		9 (9->8)
9*678		-6102		
		-537 <0 correction!		8
		5565		
8*678		-5424		
		1410 [14:6=2]		2
2*678		-1356		
		0540 [05:6=0]		0
0*678		-0		
		5400 [54:6=8]		9 (2x correction: 9->8, 8->7)
		...		
		12345:678 = 182079...		

We implement the division similar like pupils do it in the school (only the numeral system with base 10000 instead 10 is actually used, but we keep with base 10 in our illustrations). At each step of the

---

`\apMULg`: 18, 20   `\apMULh`: 20   `\apMULi`: 20   `\apMULj`: 20   `\apMULo`: 20   `\apMULT`: 20



operation, we get first two Digits from the dividend or remainder (called partial dividend or remainder) and do divide it by the first nonzero Digit of the divisor (called partial divisor). Unfortunately, the resulted Digit cannot be the definitive value of the result. We are able to find out this after the whole divisor is multiplied by resulted Digit and compared with the whole remainder. We cannot do this test immediately but only after a lot of following operations (imagine that the remainder and divisor have a huge number of Digits).

We need to subtract the remainder by the multiple of the divisor at each step. This means that we need to calculate the transmissions from the Digit position to the next Digit position from right to left (in the scheme illustrated above). Thus we need to reverse the order of Digits in the remainder and divisor. We do this reversion only once at the preparation state of the division and we interleave the data from the divisor and the dividend (the dividend will be replaced by the remainder, next by next remainder etc.).

The number of Digits of the dividend can be much greater than the number of Digits of the divisor. We need to calculate only with the first part of dividend/remainder in such case. We need to get only one new Digit from the rest of dividend at each calculation step. The illustration follows:

```
...used dividend.. | ... rest of dividend ... | .... divisor ....
1234567890123456789 7890123456789012345678901234 : 1231231231231231
xxxxxxxxxxxxxxxxxxx 7 <- calculated remainder
xxxxxxxxxxxxxxxxxxx x8 <- new calculated remainder
xxxxxxxxxxxxxxxxxxx xx9 <- new calculated remainder etc.
```

We'll interleave only the “used dividend” part with the divisor at the preparation state. We'll put the “rest of dividend” to the input stream in the normal order. The macros do the iteration over calculation steps and they can read only one new Digit from this input stream if they need it. This approach needs no manipulation with the (potentially long) “rest of the dividend” at each step. If the divisor has only one Digit (or comparable small Digits) then the algorithm has only linear complexity with respect to the number of Digits in the dividend.

The numeral system with the base 10000 brings a little problem: we are simply able to calculate the number of digits which are multiple of four. But user typically wishes another number of calculated decimal digits. We cannot simply strip the trailing digits after calculation because the user needs to read the right remainder. This is a reason why we calculate the number of digits for the first Digit of the result. All another calculated Digits will have four digits. We need to prepare the first “partial dividend” in order to the  $F$  digits will be calculated first. How to do it? Suppose the following illustration of the first two Digits in the “partial remainder” and “partial divisor”:

```
0000 7777 : 1111 = 7    .. one digit in the result
0007 7778 : 1111 = 70   .. two digits in the result
0077 7788 : 1111 = 700  .. three digits in the result
0777 7888 : 1111 = 7000 .. four digits in the result
7777 8888 : 1111 = ???  .. not possible in the numeral system of base 10000
```

We need to read  $F-1$  digits to the first Digit and four digits to the second Digit of the “partial dividend”. But this is true only if the dividend is “comparably greater or equal to” divisor. The word “comparably greater” means that we ignore signs and the decimal point in compared numbers and we assume the decimal points in the front of both numbers just before the first nonzero digit. It is obvious that if the dividend is “comparably less” than divisor then we need to read  $F$  digits to the first Digit.

The macro `\apDIV` runs `\apDIVa` macro which uses the `\tmpa` (dividend) and `\tmpb` (divisor) macros and does the following work:

- If the divisor `\tmpb` is equal to zero, print error and do nothing more (line 324).
- The `\apSIGN` of the result is calculated (line 325).
- If the dividend `\tmpa` is equal to zero, then `\OUT` and `\XOUT` are zeros and do nothing more (line 326).
- Calculate the exponent of ten `\apE` when scientific notation is used (Line 326).
- The number of digits before point are counted by `\apDIG` macro for both parameters. The difference is saved to `\apnumD` and this is the number of digits before decimal point in the result

---

`\apDIV`: 5, 6, 10–11, 13, 22, 34, 36–46    `\apDIVa`: 22, 28

(the exception is mentioned later). The `\apDIG` macro removes the decimal point and (possible) left zeros from its parameter and saves the result to the `\apnumD` register (lines 328 to 330).

- The macro `\apDIVcomp⟨paramA⟩⟨paramB⟩` determines if the `⟨paramA⟩` is “comparably greater or equal” to `⟨paramB⟩`. The result is stored in the boolean value `apX`. We can ask to this by the `\ifapX⟨true⟩\else⟨false⟩\fi` construction (line 331).
- If the dividend is “comparably greater or equal” to the divisor, then the position of decimal point in the result `\apnumD` has to be shifted by one to the right. The same is completed with `\apnumH` where the position of decimal point of the remainder will be stored (line 332).
- The number of desired digits in the result `\apnumC` is calculated (lines 333 to 339).
- If the number of desired digits is zero or less than zero then do nothing more (line 339).
- Finish the calculation of the position of decimal point in the remainder `\apnumH` (line 332).
- Calculate the number of digits in the first Digit `\apnumF` (line 343).
- Read first four digits of the divisor by the macro `\apIVread⟨sequence⟩`. Note that this macro puts trailing zeros to the right if the data stream `⟨param⟩` is shorter than four digits. If it is empty then the macro returns zero. The returned value is saved in `\apnumX` and the `⟨sequence⟩` is redefined by new value of the `⟨param⟩` where the read digits are removed (line 344).
- We need to read only `\apnumF` (or `\apnumF - 1`) digits from the `\tmpa`. This is done by the `\apIVreadX` macro at line 346. The second Digit of the “partial dividend” includes four digits and it is read by `\apIVread` macro at line 348.
- The “partial dividend” is saved to the `\apDIVxA` macro and the “partial divisor” is stored to the `\apDIVxB` macro. Note, that the second Digit of the “partial dividend” isn’t expanded by simply `\the`, because when `\apnumX=11` and `\apnumA=2222` (for example), then we need to save 22220011. These trailing zeros from left are written by the `\apIVwrite` macro (lines 349 to 350).
- The `\XOUT` macro for the currently computed remainder is initialized. The special interleaved data format of the remainder `\XOUT` is described below (line 351).
- The `\OUT` macro is initialized. The `\OUT` is generated as literal macro. First possible `⟨sign⟩`, then digits. If the number of effective digits before decimal point `\apnumD` is negative, the result will be in the form `.000123` and we need to add the zeros by the `\apADDzeros` macro (lines 352 to 353).
- The registers for main loop are initialized. The `\apnumE` signalizes that the remainder of the partial step is zero and we can stop the calculation. The `\apnumZ` will include the Digit from the input stream where the “rest of dividend” will be stored (line 353).
- The main calculation loop is processed by the `\apDIVg` macro (line 355).
- If the division process stops before the position of the decimal point in the result (because there is zero remainder, for example) then we need to add the rest of zeros by `\apADDzeros` macro. This is actual for the results of the type 1230000 (line 356).
- If the remainder isn’t equal to zero, we need to extract the digits of the remainder from the special data format to the human readable form. This is done by the `\apDIVv` macro. The decimal point is inserted to the remainder by the `\apROLLa` macro (lines 358 to 359).

apnum.tex

```

322: \def\apDIV{\relax \apPPab\apDIVa}
323: \def\apDIVa{%
324:   \ifnum\apSIGNb=0 \apERR{Dividing by zero}\else
325:     \apSIGN=\apSIGNa \multiply\apSIGN by\apSIGNb
326:     \ifnum\apSIGNa=0 \def\OUT{0}\def\XOUT{0}\apE=0 \apSIGN=0 \else
327:       \apE=\apEa \advance\apE by-\apEb
328:       \apDIG\tmpb\relax \apnumB=\apnumD
329:       \apDIG\tmpa\relax \apnumH=\apnumD
330:       \advance\apnumD by-\apnumB % \apnumD = num. of digits before decimal point in the result
331:       \apDIVcomp\tmpa\tmpb % apXtrue <=> A>=B, i.e 1 digit from A/B
332:       \ifapX \advance\apnumD by1 \advance\apnumH by1 \fi
333:       \apnumC=\apTOT
334:       \ifnum\apTOT<0 \apnumC=-\apnumC
335:         \ifnum\apnumD>\apnumC \apnumC=\apnumD \fi
336:       \fi
337:       \ifnum\apTOT=0 \apnumC=\apFRAC \advance\apnumC by\apnumD
338:       \else \apnumX=\apFRAC \advance\apnumX by\apnumD
339:       \ifnum\apnumC>\apnumX \apnumC=\apnumX \fi

```

```

340:      \fi
341:      \ifnum\apnumC>0 % \apnumC = the number of digits in the result
342:      \advance\apnumH by-\apnumC % \apnumH = the position of decimal point in the remainder
343:      \apIVmod \apnumC \apnumF % \apnumF = the number of digits in the first Digit
344:      \apIVread\tmpb \apnumB=\apnumX % \apnumB = partial divisor
345:      \apnumX=\apnumF \ifapX \advance\apnumX by-1 \fi
346:      \apIVreadX\apnumX\tmpa
347:      \apnumA=\apnumX % \apnumA = first Digit of the partial dividend
348:      \apIVread\tmpa % \apnumX = second Digit of the partial dividend
349:      \edef\apDIVxA{\the\apnumA\apIVwrite\apnumX}% first partial dividend
350:      \edef\apDIVxB{\the\apnumB}% partial divisor
351:      \edef\XOUT{\apDIVxB}{\the\apnumX}@{\the\apnumA}}% the \XOUT is initialized
352:      \edef\OUT{\ifnum\apSIGN<0-\fi}%
353:      \ifnum\apnumD<0 \edef\OUT{\OUT.}\apnumZ=-\apnumD \apADDzeros\OUT \fi
354:      \apnumE=1 \apnumZ=0
355:      \let\apNext=\apDIVg \apNext % <--- the main calculation loop is here
356:      \ifnum\apnumD>0 \apnumZ=\apnumD \apADDzeros\OUT \fi
357:      \ifnum\apnumE=0 \def\XOUT{0}\else % extracting remainder from \XOUT
358:      \edef\XOUT{\expandafter}\expandafter\apDIVv\XOUT
359:      \def\tmpc{\apnumH}\apnumG=\apSIGNa \expandafter\apROLLa\XOUT.@\XOUT
360:      \fi
361:      \else
362:      \def\OUT{0}\def\XOUT{0}\apE=0 \apSIGN=0
363:      \fi\fi\fi
364:  }

```

The macro `\apDIVcomp`  $\langle paramA \rangle \langle paramB \rangle$  provides the test if the  $\langle paramA \rangle$  is “comparably greater or equal” to  $\langle paramB \rangle$ . Imagine the following examples:

```

123456789 : 123456789 = 1
123456788 : 123456789 = .99999999189999992628

```

The example shows that the last digit in the operands can be important for the first digit in the result. This means that we need to compare whole operands but we can stop the comparison when the first difference in the digits is found. This is lexicographic ordering. Because we don’t assume the existence of eTeX (or another extensions), we need to do this comparison by macros. We set the  $\langle paramA \rangle$  and  $\langle paramB \rangle$  to the `\tmpc` and `\tmpd` respectively. The trailing `\apNLs` are appended. The macro `\apDIVcompA` reads first 8 digits from first parameter and the macros `\apDIVcompB` reads first 8 digits from second parameter and does the comparison. If the numbers are equal then the loop is processed again.

```

365: \def\apDIVcomp#1#2{%
366:   \expandafter\def\expandafter\tmpc\expandafter{#1\apNL\apNL\apNL\apNL\apNL\apNL\apNL}%
367:   \expandafter\def\expandafter\tmpd\expandafter{#2\apNL\apNL\apNL\apNL\apNL\apNL\apNL}%
368:   \def\apNext{\expandafter\expandafter\expandafter\apDIVcompA\expandafter\tmpc\tmpd}%
369:   \apXtrue \apNext
370: }
371: \def\apDIVcompA#1#2#3#4#5#6#7#8#9@{%
372:   \ifx#8\apNL \def\tmpc{0000000\apNL@}\else\def\tmpc{#9@}\fi
373:   \apnumX=#1#2#3#4#5#6#7#8\relax
374:   \apDIVcompB
375: }
376: \def\apDIVcompB#1#2#3#4#5#6#7#8#9@{%
377:   \ifnum\apnumX<#1#2#3#4#5#6#7#8 \let\apNext=\relax \apXfalse \else
378:   \ifnum\apnumX>#1#2#3#4#5#6#7#8 \let\apNext=\relax \apXtrue
379:   \fi\fi
380:   \ifx\apNext\relax\else
381:     \ifx#8\apNL \def\tmpd{0000000\apNL@}\ifx\tmpc\tmpd\let\apNext=\relax\fi
382:     \else\def\tmpd{#9@}\fi
383:   \fi
384:   \apNext
385: }

```

apnum.tex

---

`\apDIVcomp`: 22–23    `\apDIVcompA`: 23    `\apDIVcompB`: 23

The format of interleaved data with divisor and remainder is described here. Suppose this partial step of the division process:

R0	R1	R2	R3	...	Rn	:	d1	d2	d3	...	dn	=	...	A...
@	-A*d1	-A*d2	-A*d3	...	-A*dn							[	R0 R1 : d1 = A ]	
0	N0	N1	N2	...	N(n-1) Nn									

The  $R_k$  are Digits of the remainder,  $d_k$  are Digits of the divisor. The  $A$  is calculated Digit in this step. The calculation of the Digits of the new remainder is hinted here. We need to do this from right to left because of the transmissions. This implies, that the interleaved format of `\XOUT` is in the reverse order and looks like

dn	Rn	...	d3	R3	d2	R2	d1	R1	@	R0
----	----	-----	----	----	----	----	----	----	---	----

for example for  $\langle paramA \rangle = 1234567893$ ,  $\langle paramB \rangle = 454502$  (in the human readable form) the `\XOUT` should be `{200}{9300}{4545}{5678}@{1234}` (in the special format). The Digits are separated by `\TeX` braces `{}`. The resulted digit for this step is  $A = 12345678/1415 = 2716$ .

The calculation of the new remainder takes  $d_k$ ,  $R_k$ ,  $d_{k-1}$  for each  $k$  from  $n$  to 0 and creates the Digit of the new remainder  $N_{k-1} = R_k - A \cdot d_k$  (roughly speaking, actually it calculates transmissions too) and adds the new couple  $d_{k-1}$   $N_{k-1}$  to the new version of `\XOUT` macro. The zero for  $N_{-1}$  should be reached. If it is not completed then a correction of the type  $A := A - 1$  have to be done and the calculation of this step is processed again.

The result in the new `\XOUT` should be (after one step is done):

dn	Nn	...	d3	N3	d2	N2	d1	N1	@	N0
----	----	-----	----	----	----	----	----	----	---	----

where  $N_n$  is taken from the “rest of the dividend” from the input stream.

The initialization for the main loop is done by `\apDIVg` macro. It reads the Digits from `\tmpa` (dividend) and `\tmpb` macros (using `\apIVread`) and appends them to the `\XOUT` in described data format. This initialization is finished when the `\tmpb` is empty. If the `\tmpa` is not empty in such case, we put it to the input stream using `\expandafter\apDIVh\tmpa` followed by four `\apNLs` (which simply expands zero digit) followed by stop-mark. The `\apDIVh` reads one Digit from input stream. Else we put only the stop-mark to the input stream and run the `\apDIVi`. The `\apNexti` is set to the `\apDIVi`, so the macro `\apDIVh` will be skipped forever and no new Digit is read from input stream.

apnum.tex

```

386: \def\apDIVg{%
387:   \ifx\tmpb\empty
388:     \ifx\tmpa\empty \def\apNext{\apDIVi!}\let\apNexti=\apDIVi
389:     \else \def\apNext{\expandafter\apDIVh\tmpa\apNL\apNL\apNL\apNL!}\let\apNexti=\apDIVh
390:   \fi\fi
391:   \ifx\apNext\apDIVg
392:     \apIVread\tmpa \apnumA=\apnumX
393:     \apIVread\tmpb
394:     \edef\XOUT{{\the\apnumX}{\the\apnumA}\XOUT}}%
395:   \fi
396:   \apNext
397: }
```

The macro `\apDIVh` reads one Digit from data stream (from the rest of the dividend) and saves it to the `\apnumZ` register. If the stop-mark is reached (this is recognized that the last digit is the `\apNL`), then `\apNexti` is set to `\apDIVi`, so the `\apDIVh` is never processed again.

apnum.tex

```

398: \def\apDIVh#1#2#3#4{\apnumZ=#1#2#3#4
399:   \ifx\apNL#4\let\apNexti=\apDIVi\fi
400:   \apDIVi
401: }
```

The macro `\apDIVi` contains the main loop for division calculation. The core of this loop is the macro call `\apDIVp<data>` which adds next digit to the `\OUT` and recalculates the remainder.

The macro `\apDIVp` decreases the `\apnumC` register (the desired digits in the output) by four, because four digits will be calculated in the next step. The loop is processed while `\apnumC` is positive. The `\apnumZ` (new Digit from the input stream) is initialized as zero and the `\apNexti` runs the next

step of this loop. This step starts from `\apDIVh` (reading one digit from input stream) or directly the `\apDIVi` is repeated. If the remainder from the previous step is calculated as zero (`\apnumE=0`), then we stop prematurely. The `\apDIVj` macro is called at the end of the loop because we need to remove the “rest of the dividend” from the input stream.

apnum.tex

```

402: \def\apDIVi{%
403:   \ifnum\apnumE=0 \apnumC=0 \fi
404:   \ifnum\apnumC>0
405:     \expandafter\apDIVp\XOUT
406:     \advance\apnumC by-4
407:     \apnumZ=0
408:     \expandafter\apNexti
409:   \else
410:     \expandafter\apDIVj
411:   \fi
412: }
413: \def\apDIVj#1!{}

```

The macro `\apDIVp` *<interleaved data>*@ does the basic setting before the calculation through the expanded `\XOUT` is processed. The `\apDIVxA` includes the “partial dividend” and the `\apDIVxB` includes the “partial divisor”. We need to do `\apDIVxA` over `\apDIVxB` in order to obtain the next digit in the output. This digit is stored in `\apnumA`. The `\apnumX` is the transmission value, the `\apnumB`, `\apnumY` will be the memory of the last two calculated Digits in the remainder. The `\apnumE` will include the maximum of all digits of the new remainder. If it is equal to zero, we can finish the calculation.

The new interleaved data will be stored to the `\apOUT:<num>` macros in similar way as in the `\apMUL` macro. This increases the speed of the calculation. The data `\apnumO`, `\apnumL` and `\apOUTl` for this purpose are initialized.

The `\apDIVq` is started and the tokens `0\apnumZ` are appended to the input stream (i.e to the expanded `\XOUT`. This zero will be ignored and the `\apnumZ` will be used as a new  $N_n$ , i.e. the Digit from the “rest of the dividend”.

apnum.tex

```

414: \def\apDIVp{%
415:   \apnumA=\apDIVxA \divide\apnumA by\apDIVxB
416:   \def\apOUTl{}\apnumO=1 \apnumL=0
417:   \apnumX=0 \apnumB=0 \apnumE=0
418:   \let\apNext=\apDIVq \apNext 0\apnumZ
419: }

```

The macro `\apDIVq`  $\langle d_k \rangle \langle R_k \rangle \langle d_{k-1} \rangle$  calculates the Digit of the new remainder  $N_{k-1}$  by the formula  $N_{k-1} = -A \cdot d_k + R_k - X$  where  $X$  is the transmission from the previous Digit. If the result is negative, we need to add minimal number of the form  $X \cdot 10000$  in order the result is non-negative. Then the  $X$  is new transmission value. The digit  $N_k$  is stored in the `\apnumB` register and then it is added to `\apOUT:<num>` in the order  $d_{k-1} N_{k-1}$ . The `\apnumY` remembers the value of the previous `\apnumB`. The  $d_{k-1}$  is put to the input stream back in order it would be read by the next `\apDIVq` call.

If  $d_{k-1} = @$  then we are at the end of the remainder calculation and the `\apDIVr` is invoked.

apnum.tex

```

420: \def\apDIVq#1#2#3{% B A B
421:   \advance\apnumO by-1 \ifnum\apnumO=0 \apOUTx \fi
422:   \apnumY=\apnumB
423:   \apnumB=#1\multiply\apnumB by-\apnumA
424:   \advance\apnumB by#2\advance\apnumB by-\apnumX
425:   \ifnum\apnumB<0 \apnumX=\apnumB \advance\apnumX by1
426:     \divide\apnumX by-\apIVbase \advance\apnumX by1
427:     \advance\apnumB by\the\apnumX 0000
428:   \else \apnumX=0 \fi
429:   \expandafter
430:     \edef\csname apOUT:\apOUTn\endcsname{\csname apOUT:\apOUTn\endcsname{#3}\the\apnumB}}%
431:   \ifnum\apnumE<\apnumB \apnumE=\apnumB \fi
432:   \ifx@#3\let\apNext=\apDIVr \fi
433:   \apNext{#3}%
434: }

```

---

`\apDIVj`: 25      `\apDIVp`: 24–25      `\apDIVxA`: 22–23, 25–26      `\apDIVxB`: 22–23, 25  
`\apDIVq`: 25–26



The `\apDIVr` macro does the final work after the calculation of new remainder is done. It tests if the remainder is OK, i.e. the transmission from the  $R_1$  calculation is equal to  $R_0$ . If it is true then new Digit `\apnumA` is added to the `\OUT` macro else the `\apnumA` is decreased (the correction) and the calculation of the remainder is run again.

If the calculated Digit and the remainder are OK, then we do following:

- The new `\XOUT` is created from `\apOUT:⟨num⟩` macros using `\apOUTs` macro.
- The `\apnumA` is saved to the `\OUT`. This is done with care. If the `\apnumD` (where the decimal point is measured from the actual point in the `\OUT`) is in the interval  $[0, 4)$  then the decimal point have to be inserted between digits into the next Digit. This is done by `\apDIVt` macro. If the remainder is zero (`\apnumE=0`), then the right trailing zeros are removed from the Digit by the `\apDIVu` and the shift of the `\apnumD` register is calculated from the actual digits. All this calculation is done in `\tmpa` macro. The last step is adding the contents of `\tmpa` to the `\OUT`.
- The `\apnumD` is increased by the number of added digits.
- The new “partial dividend” is created from `\apnumB` and `\apnumY`.

```

435: \def\apDIVr#1#2{%
436:   \ifnum\apnumX=#2 % the calculated Digit is OK, we save it
437:     \edef\XOUT{\expandafter\apOUTs\apOUTl.}%
438:     \edef\tmpa{\ifnum\apnumF=4 \expandafter\apIVwrite\else \expandafter\the\fi\apnumA}%
439:     \ifnum\apnumD<\apnumF \ifnum\apnumD>-1 \apDIVt \fi\fi %adding dot
440:     \ifx\apNexti\apDIVh \apnumE=1 \fi
441:     \ifnum\apnumE=0 \apDIVu % removing zeros
442:       \advance\apnumD by-\apNUMdigits\tmpa \relax
443:     \else \advance\apnumD by-\apnumF \apnumF=4 \fi
444:     \edef\OUT{\OUT\tmpa}% save the Digit
445:     \edef\apDIVxA{\the\apnumB\apIVwrite\apnumY}% next partial dividend
446:   \else % we need do correction and run the remainder calculation again
447:     \advance\apnumA by-1 \apnumX=0 \apnumB=0 \apnumE=0
448:     \def\apOUTl{}\apnumO=1 \apnumL=0
449:     \def\apNext{\let\apNext=\apDIVq
450:       \expandafter\apNext\expandafter0\expandafter\apnumZ\XOUT}%
451:     \expandafter\apNext
452:   \fi
453: }
```

apnum.tex

The `\apDIVt` macro inserts the dot into digits quartet (less than four digits are allowed too) by the `\apnumD` value. This value is assumed in the interval  $[0, 4)$ . The expandable macro `\apIVdot⟨shift⟩⟨data⟩` is used for this purpose. The result from this macro has to be expanded twice.

```

454: \def\apDIVt{\edef\tmpa{\apIVdot\apnumD\tmpa}\edef\tmpa{\tmpa}}
```

apnum.tex

The `\apDIVu` macro removes trailing zeros from the right and removes the dot, if it is the last token of the `\tmpa` after removing zeros. It uses expandable macros `\apREMzerosR⟨data⟩` and `\apREMDotR⟨data⟩`.

```

455: \def\apDIVu{\edef\tmpa{\apREMzerosR\tmpa}\edef\tmpa{\apREMDotR\tmpa}}
```

apnum.tex

The rest of the code concerned with the division does an extraction of the last remainder from the data and this value is saved to the `\XOUT` macro in human readable form. The `\apDIVv` macro is called repeatedly on the special format of the `\XOUT` macro and the new `\XOUT` is created. The trailing zeros from right are ignored by the `\apDIVw`.

```

456: \def\apDIVv#1#2{\apnumX=#2
457:   \ifx#1\apDIVw{.\apIVwrite\apnumX}\else\apDIVw{\apIVwrite\apnumX}\expandafter\apDIVv\fi
458: }
459: \def\apDIVw#1{%
460:   \ifx\XOUT\empty \ifnum\apnumX=0
461:     \else \edef\tmpa{#1}\edef\XOUT{\apREMzerosR\tmpa\XOUT}%
462:   \fi
463:   \else \edef\XOUT{#1\XOUT}\fi
```

apnum.tex

---

`\apDIVr`: 25–26    `\apDIVt`: 26    `\apDIVu`: 26    `\XOUT`: 6, 5, 13, 21–26, 31, 36–37, 44  
`\apDIVv`: 22–23, 26    `\apDIVw`: 26



464: }

## 2.7 Power to the Integer

The `\apPOW` macro does the power to the integer exponent only. The `\apPOWx` is equivalent to `\apPOW` and it is used in `\evaldef` macro for the  $\wedge$  operator. If you want to redefine the meaning of the  $\wedge$  operator then redefine the `\apPOWx` sequence.

468: `\def\apPOW{\relax \apPPab\apPOWa} \let\apPOWx=\apPOW % for usage as ^ operator`

apnum.tex

We can implement the power to the integer as repeated multiplications. This is simple but slow. The goal of this section is to present the power to the integer with some optimizations.

Let  $a$  is the base of the powering computation and  $d_1, d_2, d_3, \dots, d_n$  are binary digits of the exponent (in reverse order). Then

$$p = a^{1d_1 + 2d_2 + 2^2d_3 + \dots + 2^{n-1}d_n} = (a^1)^{d_1} \cdot (a^2)^{d_2} \cdot (a^{2^2})^{d_3} \cdot (a^{2^{n-1}})^{d_n}.$$

If  $d_i = 0$  then  $z^{d_i}$  is one and this can be omitted from the queue of multiplications. If  $d_i = 1$  then we keep  $z^{d_i}$  as  $z$  in the queue. We can see from this that the  $p$  can be computed by the following algorithm:

```
(* "a" is initialized as the base, "e" as the exponent *)
p := 1;
while (e>0) {
  if (e%2) p := p*a;
  e := e/2;
  if (e>0) a := a*a;
}
(* "p" includes the result *)
```

The macro `\apPOWa` does the following work.

- After using `\apPPab` the base parameter is saved in `\tmpa` and the exponent is saved in `\tmpb`.
- In trivial cases, the result is set without any computing (lines 470 and 471).
- If the exponent is non-integer or it is too big then the error message is printed and the rest of the macro is skipped by the `\apPOWe` macro (lines 473 to 476).
- The `\apE` is calculated from `\apEa` (line 477).
- The sign of the result is negative only if the `\tmpb` is odd and base is negative (line 479).
- The number of digits after decimal point for the result is calculated and saved to `\apnumD`. The total number of digits of the base is saved to `\apnumC`. (line 480).
- The first Digit of the base needn't to include all four digits, but other Digits do it. The similar trick as in `\apMULa` is used here (lines 482 to 483).
- The base is saved in interleaved reversed format (like in `\apMULa`) into the `\OUT` macro by the `\apMULb` macro. Let it be the  $a$  value from our algorithm described above (lines 484 and 485).
- The initial value of  $p = 1$  from our algorithm is set in interleaved format into `\tmpc` macro (line 486).
- The main loop described above is processed by `\apPOWb` macro. (line 487).
- The result in `\tmpc` is converted into human readable form by the `\apPOWg` macro and it is stored into the `\OUT` macro (line 488).
- If the result is negative or decimal point is needed to print then use simple conversion of the `\OUT` macro (adding minus sign) or using `\apROLLa` macro (lines 489 and 490).
- If the exponent is negative then do the  $1/r$  calculation, where  $r$  is previous result (line 491).

469: `\def\apPOWa{%`  
 470: `\ifnum\apSIGNa=0 \def\OUT{0}\apSIGN=0 \apE=0 \else`  
 471: `\ifnum\apSIGNb=0 \def\OUT{1}\apSIGN=1 \apE=0 \else`  
 472: `\apDIG\tmpb\apnumB`  
 473: `\ifnum\apnumB>0 \apERR{POW: non-integer exponent is not implemented yet}\apPOWe\fi`

apnum.tex

`\apPOW`: 5, 6, 11, 13, 27, 34, 39–40, 44–46    `\apPOWx`: 10, 27    `\apPOWa`: 27–28

```

474:         \ifnum\apEb=0 \else \apERR{POW: the E notation of exponent isn't allowed}\apPOWE\fi
475:         \ifnum\apnumD>8 \apERR{POW: too big exponent.
476:             Do you really need about 10\the\apnumD space digits in output?}\apPOWE\fi
477:         \apE=\apEa \multiply\apE by\tmpb\relax
478:         \apSIGN=\apSIGNa
479:         \ifodd\tmpb \else \apSIGN=1 \fi
480:         \apDIG\tmpa\apnumA \apnumC=\apnumA \advance\apnumC by\apnumD
481:         \apnumD=\apnumA \multiply\apnumD by\tmpb
482:         \apIVmod \apnumC \apnumA
483:         \edef\tmpc{\ifcase\apnumA\or{}\or{}\or{}\or{}\or{}\or{}\or{}\or{}\or{}\or{}\or{}%
484:         \expandafter\expandafter\expandafter \apMULb \expandafter \tmpc \tmpa @@@@}%
485:         \edef\OUT{*\OUT}%      \OUT := \tmpa in interleaved format
486:         \def\tmpc{*.1*}%
487:         \apnumE=\tmpb\relax \apPOWb
488:         \expandafter\apPOWg \tmpc % \OUT := \tmpc in human raedable form
489:         \ifnum\apnumD=0 \ifnum \apSIGN<0 \edef\OUT{-\OUT}\fi
490:         \else \def\tmpc{-\apnumD}\apnumG=\apSIGN \expandafter\apROLLa\OUT.@\OUT\fi
491:         \ifnum\apSIGNb<0 \apPPab\apDIVa 1\OUT \fi
492:         \relax
493:     \fi\fi
494: }

```

The macro `\apPOwb` is the body of the loop in the algorithm described above. The code part after `\ifodd\apnumE` does  $p := p \cdot a$ . In order to do this, we need to convert `\OUT` (where  $a$  is stored) into normal format using `\apPOwd`. The result is saved in `\tmpb`. Then the multiplication is done by `\apMULd` and the result is normalized by the `\apPOwn` macro. Because `\apMULd` works with `\OUT` macro, we temporary set `\tmpc` to `\OUT`.

The code part after `\ifnum\apnumE<0` does `a := a*a` using the `\apPOWt` macro. The result is normalized by the `\apPOWn` macro.

```

495: \def\apPOwb{%
496:   \ifodd\apnumE   \def\tmpb{}\expandafter\apPOwd\OUT
497:                   \let\tmpd=\OUT \let\OUT=\tmpc
498:                   \expandafter\apMULD \tmpb@\expandafter\apPOwn\OUT@%
499:                   \let\tmpc=\OUT \let\OUT=\tmpd
500:   \fi
501:   \divide\apnumE by2
502:   \ifnum\apnumE>0 \expandafter\apPOwt\OUT \expandafter\apPOwn\OUT@%
503:                   \expandafter\apPOwb
504:   \fi
505: }

```

The macro `\apPOwd` *`\apPOwd` `<initialized interleaved reversed format>`* extracts the Digits from its argument and saves them to the `\tmpb` macro.

```

506: \def\apPOwd#1#2{% \apPOwd <spec format> => \tmpb (in simple reverse format)
507:   \ifx##1\expandafter\apPOwd \else
508:     \edef\tmpb{\tmpb{#1}}%
509:     \ifx##2\else \expandafter\expandafter\expandafter\apPOwd\fi
510:   \fi
511: }

```

The `\apPOWe` macro skips the rest of the body of the `\apPOWa` macro to the `\relax`. It is used when `\errmessage` is printed.

```
512: \def\apPOWe#1\relax{\fi}
```

The `\apPOWg` macro provides the conversion from interleaved reversed format to the human readable form and save the result to the `\OUT` macro. It ignores the first two elements from the format and runs `\apPOWh`.

```

513: \def\apPOWg#1#2{\def\OUT{}\apPOWh} % conversion to the human readable form
514: \def\apPOWh#1#2{\apnumA=#1
515:   \ifx#2\edef\OUT{\the\apnumA\OUT}\else \edef\OUT{\apIVwrite\apnumA\OUT}\expandafter\apPOWh\fi

```

516: }

The normalization to the initialized interleaved format of the `\OUT` is done by the `\apPOWn`  $\langle data \rangle$  macro. The `\apPOWna` reads the first part of the  $\langle data \rangle$  (to the first \*, where the Digits are non-interleaved). The `\apPOWnn` reads the second part of  $\langle data \rangle$  where the Digits of the result are interleaved with the digits of the old coefficients. We need to set the result as a new coefficients and prepare zeros between them for the new calculation. The dot after the first \* is not printed (the zero is printed instead it) but it does not matter because this token is simply ignored during the calculation.

apnum.tex

```
517: \def\apPOWn#1{\def\OUT{*}\apPOWna}
518: \def\apPOWna#1{\ifx*#1\expandafter\apPOWnn\else \edef\OUT{\OUT0{#1}}\expandafter\apPOWna\fi}
519: \def\apPOWnn#1#2{\ifx*#1\edef\OUT{\OUT*}\else\edef\OUT{\OUT0{#1}}\expandafter\apPOWnn\fi}
```

The powering to two (`\OUT:=\OUT^2`) is provided by the `\apPOWt`  $\langle data \rangle$  macro. The macro `\apPOWu` is called repeatedly for each `\apnumA=Digit` from the  $\langle data \rangle$ . One line of the multiplication scheme is processed by the `\apPOWv`  $\langle data \rangle$  macro. We can call the `\apMULe` macro here but we don't do it because a slight optimization is used here. You can try to multiply the number with digits `abcd` by itself in the mirrored multiplication scheme. You'll see that first line includes `a^2 2ab 2ac 2ad`, second line is intended by two columns and includes `b^2 2bc 2bd`, next line is intended by next two columns and includes `c^2 2cd` and the last line is intended by next two columns and includes only `d^2`. Such calculation is slightly shorter than normal multiplication and it is implemented in the `\apPOWv` macro.

apnum.tex

```
520: \def\apPOWt#1#2{\apPOWu} % power to two
521: \def\apPOWu#1#2{\apnumA=#1
522:   \expandafter\apPOWv\OUT
523:   \ifx*#2\else \expandafter\apPOWu\fi
524: }
525: \def\apPOWv#1#2#3#4{\def\apOUT1{}\apnum0=1 \apnumL=0
526:   \apnumB=\apnumA \multiply\apnumB by\apnumB \multiply\apnumA by2
527:   \ifx*#4\else\advance\apnumB by#4 \fi
528:   \ifx\apnumB<\apIVbase \apnumX=0 \else \apIVtrans \fi
529:   \edef\OUT{#1{#2}{the\apnumB}*}%
530:   \ifx*#4\apMULf0*\else\expandafter\apMULf\fi
531: }
```

## 2.8 apROLL, apROUND and apNORM Macros

The macros `\apROLL`, `\apROUND` and `\apNORM` are implemented by `\apROLLa`, `\apROUNDa` and `\apNORMa` macros with common format of the parameter text:  $\langle expanded\ sequence \rangle.\langle sequence \rangle$  where  $\langle expanded\ sequence \rangle$  is the expansion of the macro  $\langle sequence \rangle$  (given as first parameter of `\apROLL`, `\apROUND` and `\apNORM`, but without optionally minus sign. If there was the minus sign then `\apnumG=-1` else `\apnumG=1`. This preparation of the parameter  $\langle sequence \rangle$  is done by the `\apPPs` macro. The second parameter of the macros `\apROLL`, `\apROUND` and `\apNORM` is saved to the `\tmpc` macro.

`\apROLLa`  $\langle param \rangle.\langle sequence \rangle$  shifts the decimal point of the  $\langle param \rangle$  by `\tmpc` positions to the right (or to the left, if `\tmpc` is negative) and saves the result to the  $\langle sequence \rangle$  macro. The `\tmpc` value is saved to the `\apnumA` register and the `\apROLLc` is executed if we need to shift the decimal point to left. Else `\apROLLg` is executed.

apnum.tex

```
535: \def\apROLL{\apPPs\apROLLa}
536: \def\apROLLa{\apnumA=\tmpc\relax \ifnum\apnumA<0 \expandafter\apROLLc\else \expandafter\apROLLg\fi}
```

The `\apROLLc`  $\langle param \rangle.\langle sequence \rangle$  shifts the decimal point to left by the `-\apnumA` decimal digits. It reads the tokens from the input stream until the dot is found using `\apROLLd` macro. The number of such tokens is set to the `\apnumB` register and tokens are saved to the `\tmpc` macro. If the dot is found then `\apROLLe` does the following: if the number of read tokens is greater then the absolute value of the  $\langle shift \rangle$ , then the number of positions from the most left digit of the number to the desired place of the dot is set to the `\apnumA` register a the dot is saved to this place by `\apROLLi`  $\langle parameter \rangle.\langle sequence \rangle$ .

---

```
\apPOWn: 28–29   \apPOWna: 29   \apPOWnn: 29   \apPOWt: 28–29   \apPOWu: 29   \apPOWv: 29
\apROLL: 4, 5, 13, 29, 32, 37–38, 40, 44, 46   \apROUND: 5, 13, 29, 31, 36–37, 39–40, 42, 44–46
\apNORM: 4, 5, 13, 29, 32, 46   \apROLLa: 17, 22–23, 27–29, 31–32   \apROLLc: 29–30   \apROLLd: 30
\apROLLe: 30
```

Else the new number looks like .000123 and the right number of zeros are saved to the  $\langle sequence \rangle$  using the `\apADDzeros` macro and the rest of the input stream (including expanded `\tmpc` returned back) is appended to the macro  $\langle sequence \rangle$  by the `\apROLLf`  $\langle param \rangle$ .@ macro.

```

537: \def\apROLLc{\edef\tmpc{\edef\tmpd{\ifnum\apnumG<0-\fi}\apnumB=0 \apROLLd}
538: \def\apROLLd#1{%
539:   \ifx.#1\expandafter\apROLLe
540:   \else \edef\tmpc{\tmpc#1}%
541:     \advance\apnumB by1
542:     \expandafter\apROLLd
543:   \fi
544: }
545: \def\apROLLe#1{\ifx@#1\edef\tmpc{\tmpc.@}\else\edef\tmpc{\tmpc#1}\fi
546:   \advance\apnumB by\apnumA
547:   \ifnum\apnumB<0
548:     \apnumZ=-\apnumB \edef\tmpd{\tmpd.}\apADDzeros\tmpd
549:     \expandafter\expandafter\expandafter\apROLLf\expandafter\tmpc
550:   \else
551:     \apnumA=\apnumB
552:     \expandafter\expandafter\expandafter\apROLLi\expandafter\tmpc
553:   \fi
554: }
555: \def\apROLLf#1.@#2{\edef#2{\tmpd#1}}

```

apnum.tex

The `\apROLLg`  $\langle param \rangle$ .@ $\langle sequence \rangle$  shifts the decimal point to the right by `\apnumA` digits starting from actual position of the input stream. It reads tokens from the input stream by the `\apROLLh` and saves them to the `\tmpd` macro where the result will be built. When dot is found the `\apROLLi` is processed. It reads next tokens and decreases the `\apnumA` by one for each token. It ends (using `\apROLLj``\apROLLk`) when `\apnumA` is equal to zero. If the end of the input stream is reached (the @ character) then the zero is inserted before this character (using `\apROLLj``\apROLLi0@`). This solves the situations like 123,  $\langle shift \rangle=2$ ,  $\rightarrow 12300$ .

```

556: \def\apROLLg#1{\edef\tmpd{\ifnum\apnumG<0-\fi}\ifx.#1\apnumB=0 \else\apnumB=1 \fi \apROLLh#1}
557: \def\apROLLh#1{\ifx.#1\expandafter\apROLLi\else \edef\tmpd{\tmpd#1}\expandafter\apROLLh\fi}
558: \def\apROLLi#1{\ifx.#1\expandafter\apROLLi\else
559:   \ifnum\apnumA>0 \else \apROLLj \apROLLk#1\fi
560:   \ifx@#1\apROLLj \apROLLi0@\fi
561:   \advance\apnumA by-1
562:   \ifx0#1\else \apnumB=1 \fi
563:   \ifnum\apnumB>0 \edef\tmpd{\tmpd#1}\fi
564:   \expandafter\apROLLi\fi
565: }

```

apnum.tex

The `\apROLLg` macro initializes `\apnumB=1` if the  $\langle param \rangle$  doesn't begin by dot. This is a flag that all digits read by `\apROLLi` have to be saved. If the dot begins, then the number can look like .000123 (before moving the dot to the right) and we need to ignore the trailing zeros. The `\apnumB` is equal to zero in such case and this is set to 1 if here is first non-zero digit.

The `\apROLLj` macro closes the conditionals and runs its parameter separated by `\fi`. It skips the rest of the `\apROLLi` macro too.

```

566: \def\apROLLj#1\fi#2\apROLLi\fi{\fi\fi#1}

```

apnum.tex

The macro `\apROLLk` puts the decimal point to the `\tmpd` at current position (using `\apROLLn`) if the input stream is not fully read. Else it ends the processing. The result is an integer without decimal digit in such case.

```

567: \def\apROLLk#1{\ifx@#1\expandafter\apROLLo\expandafter@\else
568:   \def\tmpc{\apnumB=0 \expandafter\apROLLn\expandafter#\fi
569: }

```

apnum.tex

The macro `\apROLLn` reads the input stream until the dot is found. Because we read now the digits after a new position of the decimal point we need to check situations of the type 123.000 which is

---

`\apROLLf`: 30    `\apROLLg`: 29–30    `\apROLLh`: 30    `\apROLLi`: 29–30    `\apROLLj`: 30    `\apROLLk`: 30  
`\apROLLn`: 30–31

needed to be written as 123 without decimal point. This is a reason of a little complication. We save all digits to the `\tmpc` macro and calculate the sum of such digits in `\apnumB` register. If this sum is equal to zero then we don't append the `\tmpc` to the `\tmpd`. The macro `\apROLLn` is finished by the `\apROLLo`

```

570: \def\apROLLn#1{%
571:   \ifx.#1\ifnum\apnumB>0 \edef\tmpd{\tmpd.\tmpc}\fi \expandafter\apROLLo
572:   \else \edef\tmpc{\tmpc#1}\advance\apnumB by#1 \expandafter\apROLLn
573:   \fi
574: }
575: \def\apROLLo@#1{\let#1=\tmpd}

```

The macro `\apROUNDa`  $\langle param \rangle . @ \langle sequence \rangle$  rounds the number given in the  $\langle param \rangle$ . The number of digits after decimal point `\tmpc` is saved to `\apnumD`. If this number is negative then `\apROUNDe` is processed else the `\apROUNDb` reads the  $\langle param \rangle$  to the decimal point and saves this part to the `\tmpc` macro. The `\tmpd` macro (where the rest after decimal point of the number will be stored) is initialized to empty and the `\apROUNDc` is started. This macro reads one token from input stream repeatedly until the number of read tokens is equal to `\apnumD` or the stop mark `@` is reached. All tokens are saved to `\tmpd`. Then the `\apROUNDd` macro reads the rest of the  $\langle param \rangle$ , saves it to the `\XOUT` macro and defines  $\langle sequence \rangle$  (i.e. #2) as the rounded number.

```

577: \def\apROUND{\apPPs\apROUNDa}
578: \def\apROUNDa{\apnumD=\tmpc\relax
579:   \ifnum\apnumD<0 \expandafter\apROUNDe
580:   \else \expandafter\apROUNDb
581:   \fi
582: }
583: \def\apROUNDb#1.{\edef\tmpc{#1}\apnumX=0 \def\tmpd{}\let\apNext=\apROUNDc \apNext}
584: \def\apROUNDc#1{\ifx@#1\def\apNext{\apROUNDd.}%
585:   \else \advance\apnumD by-1
586:   \ifnum\apnumD<0 \def\apNext{\apROUNDd#1}%
587:   \else \ifx.#1\else \advance\apnumX by#1 \edef\tmpd{\tmpd#1}\fi
588:   \fi
589:   \fi \apNext
590: }
591: \def\apROUNDd#1.@#2{\def\XOUT{#1}\edef\XOUT{\apREMzerosR\XOUT}%
592:   \ifnum\apnumX=0 \def\tmpd{}\fi
593:   \ifx\tmpd\empty \def#2{0}%
594:   \else \edef#2{\ifnum\apnumG<0-\fi\tmpc}\fi
595:   \else\edef#2{\ifnum\apnumG<0-\fi\tmpc.\tmpd}\fi
596:   \fi
597: }

```

The macro `\apROUNDe` solves the “less standard” problem when rounding to the negative digits after decimal point `\apnumD`, i.e. we need to set  $-\text{apnumD}$  digits before decimal point to zero. The solution is to remove the rest of the input stream, use `\apROLLa` to shift the decimal point left by  $-\text{apnumD}$  positions, use `\apROUNDa` to remove all digits after decimal point and shift the decimal point back to its previous place.

```

598: \def\apROUNDe#1.@#2{\apnumC=\apnumD
599:   \apPPs\apROLLa#2{\apnumC}\apPPs\apROUNDa#2{0}\apPPs\apROLLa#2{-\apnumC}%
600: }

```

The macro `\apNORMa` redefines the  $\langle sequence \rangle$  in order to remove minus sign because the `\apDIG` macro uses its parameter without this sign. Then the `\apNORMb`  $\langle sequence \rangle \langle parameter \rangle @$  is executed where the dot in the front of the parameter is tested. If the dot is here then the `\apDIG` macro measures the digits after decimal point too and the `\apNORMc` is executed (where the `\apROLLa` shifts the decimal point from the right edge of the number). Else the `\apDIG` macro doesn't measure the digits after decimal point and the `\apNORMd` is executed (where the `\apROLLa` shifts the decimal point from the left edge of the number).

---

<code>\apROLLo</code> : 30–31	<code>\apROUNDa</code> : 13, 29, 31	<code>\apROUNDb</code> : 31	<code>\apROUNDc</code> : 31	<code>\apROUNDd</code> : 31
<code>\apROUNDe</code> : 31	<code>\apNORMa</code> : 29, 32	<code>\apNORMb</code> : 32	<code>\apNORMc</code> : 32	<code>\apNORMd</code> : 32

```

601: \def\apNORM{\apPPs\apNORMa}
602: \def\apNORMa#1.0#2{\ifnum\apnumG<0 \def#2{#1}\fi \expandafter\apNORMb\expandafter#2\tmpc0}
603: \def\apNORMb#1#2#3@{%
604:   \ifx.#2\apnumC=#3\relax \apDIG#1\apnumA \apNORMc#1%
605:   \else \apnumC=#2#3\relax \apDIG#1\relax \apNORMd#1%
606:   \fi
607: }
608: \def\apNORMc#1{\advance\apE by-\apnumA \advance\apE by\apnumC
609:   \def\tmpc{-\apnumC}\expandafter\apROLLa#1.0#1%
610: }
611: \def\apNORMd#1{\advance\apE by\apnumD \advance\apE by-\apnumC
612:   \def\tmpc{\apnumC}\expandafter\apROLLa\expandafter.#1.0#1%
613: }

```

The macro `\apEadd`  $\langle sequence \rangle$  adds E in scientific format into  $\langle sequence \rangle$  macro and `\apEnum`  $\langle sequence \rangle$  normalizes the number in the  $\langle sequence \rangle$ . After processing these macros the `\apE` register is set to zero.

```

614: \def\apEadd#1{\ifnum\apE=0 \else\edef#1{#1E\ifnum\apE>0+\fi\the\apE}\apE=0 \fi}
615: \def\apEnum#1{\ifnum\apE=0 \else\apROLL#1\apE \apE=0 \fi}

```

## 2.9 Miscellaneous Macros

The macro `\apEND` closes the `\begingroup` group, but keeps the values of `\OUT` macro and `\apSIGN`, `\apE` registers.

```

619: \def\apEND{\global\let\apENDx=\OUT
620:   \edef\tmpb{\apSIGN=\the\apSIGN \apE=\the\apE}%
621:   \expandafter\endgroup \tmpb \let\OUT=\apENDx
622: }

```

The macro `\apDIG`  $\langle sequence \rangle \langle register \text{ or } relax \rangle$  reads the content of the macro  $\langle sequence \rangle$  and counts the number of digits in this macro before decimal point and saves it to `\apnumD` register. If the macro  $\langle sequence \rangle$  includes decimal point then it is redefined with the same content but without decimal point. The numbers in the form .00123 are replaced by 123 without zeros, but `\apnumD=-2` in this example. If the second parameter of the `\apDIG` macro is `\relax` then the number of digits after decimal point isn't counted. Else the number of these digits is stored to the given  $\langle register \rangle$ .

The macro `\apDIG` is developed in order to do minimal operations over a potentially long parameters. It assumes that  $\langle sequence \rangle$  includes a number without  $\langle sign \rangle$  and without left trailing zeros. This is true after parameter preparation by the `\apPPab` macro.

The macro `\apDIG` prepares an incrementation in `\tmpc` if the second parameter  $\langle register \rangle$  isn't `\relax`. It initializes `\apnumD` and  $\langle register \rangle$ . It runs `\apDIGa`  $\langle data \rangle . . @ \langle sequence \rangle$  which increments the `\apnumD` until the dot is found. Then the `\apDIGb` is executed (if there are no digits before dot) or the `\apDIGc` is called (if there is at least one digit before dot). The `\apDIGb` ignores zeros immediately after dot. The `\apDIGc` reads the rest of the  $\langle data \rangle$  to the `#1` and saves it to the `\tmpd` macro. It runs the counter over this  $\langle data \rangle$  `\apDIGd`  $\langle data \rangle @$  only if it is desired (`\tmpc` is non-empty). Else the `\apDIGe` is executed. The `\apDIGe`  $\langle dot \text{ or } nothing \rangle @ \langle sequence \rangle$  redefines  $\langle sequence \rangle$  if it is needed. Note, that `#1` is empty if and only if the  $\langle data \rangle$  include no dot (first dot was reached as the first dot from `\apDIG`, the second dot from `\apDIG` was a separator of `#1` in `\apDIGc` and there is nothing between the second dot and the `@` mark. The  $\langle sequence \rangle$  isn't redefined if it doesn't include a dot. Else the sequence is set to the `\tmpd` (the rest after dot) if there are no digits before dot. Else the sequence is redefined using expandable macro `\apDIGf`.

```

623: \def\apDIG#1#2{\ifx\relax#2\def\tmpc{}\else #2=0 \def\tmpc{\advance#2 by1 }\fi
624:   \apnumD=0 \expandafter\apDIGa#1..0#1%
625: }
626: \def\apDIGa#1{\ifx.#1\csname apDIG\ifnum\apnumD>0 c\else b\fi\expandafter\endcsname
627:   \else \advance\apnumD by1 \expandafter\apDIGa\fi}
628: \def\apDIGb#1{%

```

---

`\apEadd`: 4, 5, 32, 36    `\apEnum`: 4, 5, 32, 36–37, 39, 43–45    `\apEND`: 8, 11, 32, 35–40, 43–46  
`\apDIG`: 14, 17–18, 21–22, 27–28, 31–32, 38, 40, 45    `\apDIGa`: 32    `\apDIGb`: 32–33    `\apDIGc`: 32–33  
`\apDIGd`: 33    `\apDIGe`: 32–33    `\apDIGf`: 33



```

629: \ifx0#1\advance\apnumD by-1 \tmpc \expandafter\apDIGb
630: \else \expandafter\apDIGc \expandafter#1\fi
631: }
632: \def\apDIGc#1.{\def\tmpd{#1}%
633: \ifx\tmpc\empty \let\apNext=\apDIGe
634: \else \def\apNext{\expandafter\apDIGd\tmpd0}%
635: \fi \apNext
636: }
637: \def\apDIGd#1{\ifx0#1\expandafter\apDIGe \else \tmpc \expandafter\apDIGd \fi}
638: \def\apDIGe#10#2{%
639: \ifx0#10\else % #1=empty <=> the param has no dot, we need to do nothing
640: \ifnum\apnumD>0 \edef#2{\expandafter\apDIGf#20}% the dot plus digits before dot
641: \else \let#2=\tmpd % there are only digits after dot, use \tmpd
642: \fi\fi
643: }
644: \def\apDIGf#1.#2@{#1#2}

```

The macro `\apIVread`  $\langle sequence \rangle$  reads four digits from the macro  $\langle sequence \rangle$ , sets `\apnumX` as the Digit consisting from read digits and removes the read digits from  $\langle sequence \rangle$ . It internally expands  $\langle sequence \rangle$ , adds the `\apNL` marks and runs `\apIVreadA` macro which sets the `\apnumX` and redefines  $\langle sequence \rangle$ .

The usage of the `\apNL` as a stop-marks has the advantage: they act as simply zero digits in the comparison but we can ask by `\ifx` if this stop mark is reached. The `#5` parameter of `\apIVreadA` is separated by first occurrence of `\apNL`, i.e. the rest of the macro  $\langle sequence \rangle$  is here.

```
646: \def\apNL{0}
647: \def\apIVread#1{\expandafter\apIVread#1\apNL\apNL\apNL\apNL\apNL@#1}
648: \def\apIVreadA#1#2#3#4#5\apNL#6#7{\apnumX=#1#2#3#4\relax \def#7{#5}}
```

The macro `\apIVreadX`  $\langle num \rangle \langle sequence \rangle$  acts similar as `\apIVread` $\langle sequence \rangle$ , but only  $\langle num \rangle$  digits are read. The  $\langle num \rangle$  is expected in the range 0 to 4. The macro prepares the appropriate number of empty parameters in `\tmpc` and runs `\apIVreadA` with these empty parameters inserted before the real body of the  $\langle sequence \rangle$ .

```
649: \def\apIVreadX#1#2{\edef\tmpc{\ifcase#1{}\{}0\or{}\{}1\or{}\{}2\or{}\{}3\or{}\{}4\or{}\{}5\or{}\{}6\or{}\{}7\or{}\{}8\or{}\{}9\or{}\{}A\or{}\{}B\or{}\{}C\or{}\{}D\or{}\{}E\or{}\{}F\fi}%
650:   \expandafter\expandafter\expandafter\apIVreadA\expandafter\tmpc#2\apNL\apNL\apNL\apNL\apNL\@%2%
651: }
```

The macro `\apIVwrite`  $\langle num \rangle$  expands the digits from  $\langle num \rangle$  register. The number of digits are four. If the  $\langle num \rangle$  is less than 1000 then left zeros are added.

```
652: \def\apIVwrite#1{\ifnum#1<1000 0\ifnum#1<100 0\ifnum#1<10 0\fi\fi\fi\the#1}
```

The macro `\apIVtrans` calculates the transmission for the next Digit. The value (greater or equal 10000) is assumed to be in `\apnumB`. The new value less than 10000 is stored to `\apnumB` and the transmission value is stored in `\apnumX`. The constant `\apIVbase` is used instead of literal 10000 because it is quicker.

```

654: \mathchardef\apIVbase=10000
655: \def\apIVtrans{\apnumX=\apnumB \divide\apnumB by\apIVbase \multiply\apnumB by-\apIVbase
656:   \advance\apnumB by\apnumX \divide\apnumX by\apIVbase
657: }

```

The macro `\apIVmod`  $\langle length \rangle \langle register \rangle$  sets  $\langle register \rangle$  to the number of digits to be read to the first Digit, if the number has  $\langle length \rangle$  digits in total. We need to read all Digits with four digits, only first Digit can be shorter.

```
658: \def\apIVmod#1#2{#2=#1\divide#2by4 \multiply#2by-4 \advance#2by#1\relax
```

The macro `\apIVdot`  $\langle num \rangle \langle param \rangle$  adds the dot into  $\langle param \rangle$ . Let  $K = \langle num \rangle$  and  $F$  is the number of digits in the  $\langle param \rangle$ . The macro expects that  $K \in [0, 4)$  and  $F \in (0, 4]$ . The macro inserts the dot after  $K$ -th digit if  $K < F$ . Else no dot is inserted. It is expandable macro, but two full expansions

```
\apIVread: 15–16, 22–24, 33      \apIVreadA: 33      \apNL: 15, 23–24, 33      \apIVreadX: 22–23, 33
\apIVwrite: 16, 20, 22–23, 26, 28, 33      \apIVtrans: 19, 29, 33      \apIVbase: 16, 19, 25, 29, 33
\apIVmod: 14, 17–18, 23, 28, 33      \apIVdot: 20, 26, 34
```

are needed. After first expansion the result looks like `\apIVdota<dots><param>...@` where `<dots>` are the appropriate number of dots. Then the `\apIVdota` reads the four tokens (maybe the generated dots), ignores the dots while printing and appends the dot after these four tokens, if the rest #5 is non-empty.

```

662: \def\apIVdota#1#2{\noexpand\apIVdota\ifcase#1...\or...\or...\or.\fi #2...@}
663: \def\apIVdota#1#2#3#4#5.#6{\ifx.#1\else#1\fi
664:   \ifx.#2\else#2\fi \ifx.#3\else#3\fi \ifx.#4\else#4\fi\ifx.#5.\else.#5\fi
665: }

```

apnum.tex

The expandable macro `\apNUMdigits {<param>}` expands (using the `\apNUMdigitsA` macro) to the number of digits in the `<param>`. We assume that maximal number of digits will be four.

```

666: \def\apNUMdigits#1{\expandafter\apNUMdigitsA#1@@@!}
667: \def\apNUMdigitsA#1#2#3#4#5!{\ifx#4\ifx#3\ifx#2\ifx#1\else1\fi \else2\fi \else3\fi \else4\fi}

```

apnum.tex

The macro `\apADDzeros <sequence>` adds `\apnumZ` zeros to the macro `<sequence>`.

```

669: \def\apADDzeros#1{\edef#1{#10}\advance\apnumZ by-1
670:   \ifnum\apnumZ>0 \expandafter\apADDzeros\expandafter#1\fi
671: }

```

apnum.tex

The expandable macro `\apREMzerosR {<param>}` removes right trailing zeros from the `<param>`. It expands to `\apREMzerosRa<param>@0@!`. The macro `\apREMzerosRa` reads all text terminated by `@` to #1. This termination zero can be the most right zero of the `<param>` (then #2 is non-empty) or `<param>` hasn't such zero digit (then #2 is empty). If #2 is non-empty then the `\apREMzerosRa` is expanded again in the recursion. Else `\apREMzerosRb` removes the stop-mark `@` and the expansion is finished.

```

672: \def\apREMzerosR#1{\expandafter\apREMzerosRa#1@0@!}
673: \def\apREMzerosRa#10@#2!{\ifx!#2!\apREMzerosRb#1\else\apREMzerosRa#1@0@!\fi}
674: \def\apREMzerosRb#1@{#1}

```

apnum.tex

The expandable macro `\apREMDotR {<param>}` removes right trailing dot from the `<param>` if exists. It expands to `\apREMDotRa` and works similarly as the `\apREMzerosR` macro.

```

675: \def\apREMDotR#1{\expandafter\apREMDotRa#1@.@!}
676: \def\apREMDotRa#1.@#2!{\ifx!#2!\apREMzerosRb#1\else#1\fi}

```

apnum.tex

The `\apREMfirst <sequence>` macro removes the first token from the `<sequence>` macro. It can be used for removing the “minus” sign from the “number-like” macros.

```

678: \def\apREMfirst#1{\expandafter\apREMfirsta#1@#1}
679: \def\apREMfirsta#1#2@#3{\def#3{#2}}

```

apnum.tex

The writing to the `\OUT` in the `\apMUL`, `\apDIV` and `\apPOW` macros is optimized, which decreases the computation time with very large numbers ten times and more. We can do simply `\edef\OUT{\OUT<something>}` instead of

```

\expandafter\edef\csname apOUT:\apOUTn\endcsname
  {\csname apOUT:\apOUTn\endcsname<something>}%

```

but `\edef\OUT{\OUT<something>}` is typically processed very often over possibly very long macro (many thousands of tokens). It is better to do `\edef` over more short macros `\apOUT:0`, `\apOUT:1`, etc. Each such macro includes only 7 Digits pairs of the whole `\OUT`. The macro `\apOUTx` is invoked each 7 digit (the `\apnum0` register is decreased). It uses `\apnumL` value which is the `<num>` part of the next `\apOUT:<num>` control sequence. The `\apOUTx` defines this `<num>` as `\apOUTn` and initializes `\apOUT:<num>` as empty and adds the `<num>` to the list `\apOUTl`. When the creating of the next `\OUT` macro is definitely finished, the `\OUT` macro is assembled from the parts `\apOUT:0`, `\apOUT:1` etc. by the macro `\apOUTs <list of numbers><dot><comma>`.

---

```

\apIVdota: 34   \apNUMdigits: 20, 26, 34   \apNUMdigitsA: 34   \apADDzeros: 15, 18, 22–23, 30, 34
\apREMzerosR: 16, 26, 31, 34   \apREMzerosRa: 34   \apREMzerosRb: 34   \apREMDotR: 26, 34
\apREMDotRa: 34   \apREMfirst: 5, 34, 36, 39, 43–45   \apOUTx: 19, 25, 34–35   \apOUTn: 19, 25, 35
\apOUTl: 19, 25–26, 29, 35   \apOUTs: 19, 26, 35

```

```

681: \def\apOUTx{\apnum0=7
682:   \edef\apOUTn{\the\apnumL}\edef\apOUTl{\apOUTl\apOUTn,}%
683:   \expandafter\def\csname apOUT:\apOUTn\endcsname{}%
684:   \advance\apnumL by1
685: }
686: \def\apOUTs#1,{\ifx.#1\else\csname apOUT:#1\expandafter\endcsname\expandafter\apOUTs\fi}

```

If a “function-like” macro needs a local counters then it is recommended to enclose all calculation into a group `\apINIT ... \apEND`. The `\apINIT` opens the group and prepares a short name `\do` and the macro `\localcounts<counters>`; . The typical usage is:

```

\def\MACRO#1{\relax \apINIT % function-like macro, \apINIT
  \evaldef\foo{#1}%          % preparing the parameter
  \localounts \N \M \K ;%    % local \newcount\N \newcount\M \newcount\K
  ...                        % calculation
  \apEND                     % end of \apINIT group
}

```

Note that `\localcounts` is used after preparing the parameter using `\evaldef` in order to avoid name conflict of local declared “variables” and “variables” used in #1 by user.

The `\apINIT` sets locally `\localcounts` to be equivalent to `\apCOUNTS`. This macro increases the top index of allocated counters `\count10` (used in plain  $\TeX$ ) locally and declares the counters locally. It means that if the group is closed then the counters are deallocated and top index of counters `\count10` is returned to its original value.

```

688: \def\apINIT{\begingroup \let\do=\apEVALxdo \let\localcounts=\apCOUNTS}
689: \def\apCOUNTS#1{\ifx.#1\else
690:   \advance\count10 by1 \countdef#1=\count10
691:   \expandafter\apCOUNTS\fi
692: }

```

The macro `\do <sequence>=<calculation>`; allows to write the calculation of Polish expressions more synoptic:

```

\do \X=\apPLUS{2}{\the\N};%    % is equivalent to:
\apPLUS{2}{\the\N}\let\X=\OUT

```

The `\do` macro is locally set to be equivalent to `\apEVALxdo` .

```

693: \def\apEVALxdo#1=#2;{\#2\let#1=\OUT}

```

The `\apRETURN` macro must be followed by `\fi`. It skips the rest of the block `\apINIT... \apEND` typically used in “function-like” macros. The `\apERR {<text>}` macro writes `<text>` as error message and returns the processing of the block enclosed by `\apINIT... \apEND`. User can redefine it if the `\errmessage` isn’t required.

```

695: \def\apRETURN#1\apEND{\fi\apEND}
696: \def\apERR#1{\errmessage{#1}}

```

The `\apNOPT` macro removes the `pt` letters after expansion of `<dimen>` register. This is usable when we do a classical `<dimen>` calculation, see TBN page 80. Usage: `\expandafter\apNOPT\the<dimen>`.

```

698: {\lccode'\?='p \lccode'\!='t \lowercase{\gdef\apNOPT#1?!{#1}}

```

The `\loop` macro from plain  $\TeX$  is redefined here in more convenient way. It does the same as original `\loop` by D. Knuth but moreover, it allows the construction `\if... \else... \repeat`.

```

700: \def\loop#1\repeat{\def\body{#1\relax\expandafter\body\fi}\body}

```

---

`\apINIT`: 35–37, 39–40, 42–45    `\localcounts`: 35–37, 39–40, 42, 44–45    `\apCOUNTS`: 35  
`\do`: 35, 37, 39–41, 44–45    `\apEVALxdo`: 35    `\apRETURN`: 35–37, 39–40, 44–46    `\apERR`: 22, 27–28,  
35–37, 40, 44–45    `\apNOPT`: 35, 38, 41    `\loop`: 35–40, 42, 44–45

## 2.10 Function-like Macros

The implementation of function-like macros `\ABS`, `\SGN`, `\iDIV`, `\iMOD`, `\iROUND`, `\iFRAC` are simple.

```

704: \def\ABS#1{\relax          % mandatory \relax for "function-like" macros
705:   \evaldef\OUT{#1}%        % evaluation of the input parameter
706:   \ifnum\apSIGN<0          % if (input < 0)
707:     \apSIGN=1              % sign = 1
708:     \apREMfirst\OUT        % remove first "minus" from OUT
709:   \fi                      % fi
710: }
711: \def\SGN#1{\relax \evaldef\OUT{#1}\edef\OUT{\the\apSIGN}\apE=0 }
712: \def\iDIV#1#2{\relax \apINIT          % calculation in group
713:   \evaldef\apAparam{#1}\apEadd\apAparam
714:   \evaldef\apBparam{#2}\apEadd\apBparam % evaluation of the parameters
715:   \apTOT=0 \apFRAC=0 \apDIV\apAparam\apBparam % integer division
716:   \apEND                          % end of group
717: }
718: \def\iMOD#1#2{\relax \apINIT          % calculation in group
719:   \evaldef\apAparam{#1}\apEadd\apAparam
720:   \evaldef\apBparam{#2}\apEadd\apBparam % evaluation of the parameters
721:   \apTOT=0 \apFRAC=0 \apDIV\apAparam\apBparam % integer division
722:   \let\OUT=\XOUT              % remainder is the output
723:   \apEND                      % end of group
724: }
725: \def\iROUND#1{\relax \evaldef\OUT{#1}\apEnum\OUT \apROUND\OUT0}
726: \def\iFRAC#1{\relax
727:   \evaldef\OUT{#1}\apEnum\OUT \apROUND\OUT0% % preparing the parameter
728:   \ifx\XOUT\empty \def\OUT{0}\apSIGN=0      % empty fraction part means zero
729:   \else \edef\OUT{\.\XOUT}\apSIGN=1          % else OUT = dot+fraction part
730:   \fi
731: }
```

The `\FAC` macro for **factorial** doesn't use recursive call because the  $\TeX$  group is opened in such case and the number of levels of  $\TeX$  group is limited (to 255 in my computer). But we want to calculate more factorial than only 255!.

```

733: \def\FAC#1{\relax \apINIT          % "function-like" in the group, FAC = factorial
734:   \evaldef\OUT{#1}\apEnum\OUT      % preparing the parameter
735:   \localcounts \N;%                % local \newcount
736:   \ifnum\apSIGN<0 \apERR{\string\FAC: argument {\OUT} cannot be negative}\apRETURN\fi
737:   \let\tmp=\OUT \apROUND\tmp0%      % test, if parameter is integer
738:   \ifx\XOUT\empty \else \apERR{\string\FAC: argument {\OUT} must be integer}\apRETURN\fi
739:   \N=\OUT\relax                    % N = param (error here if it is an big integer)
740:   \ifnum\N=0\def\OUT{1}\apSIGN=1 \fi % special definition for factorial(0)
741:   \loop \ifnum \N>2 \advance\N by-1 % loop if (N>2) N--
742:     \apMUL{\OUT}{\the\N}\repeat    % OUT = OUT * N , repeat
743:   \apEND                          % end of group
744: }
```

The `\BINOM`  $\{a\}\{b\}$  is **binomial coefficient** defined by

$$\binom{a}{b} = \frac{a!}{b!(a-b)!} = \frac{a(a-1)\cdots(a-b+1)}{b!} \quad \text{for integer } b > 0, \quad \binom{a}{0} = 1.$$

We use the formula where  $(a-b)!$  is missing in numerator and denominator (second fraction) because of time optimization. Second advantage of such formula is that  $a$  need not to be integer. That is the reason why the `\BINOM` isn't defined simply as

```
\def\BINOM#1#2{\relax \evaldef{ \FAC{#1} / (\FAC{#2} * \FAC{(#1)-(#2)}) }}
```

The macro `\BINOM` checks if  $a$  is integer. If it is true then we choose  $\backslash C$  as minimum of  $b$  and  $a-b$ . Then we calculate factorial of  $\backslash C$  in the denominator of the formula (second fraction). And

---

`\ABS:` 3, 5, 36    `\SGN:` 3, 36    `\iDIV:` 3, 36    `\iMOD:` 3, 36    `\iROUND:` 3, 36    `\iFRAC:` 3, 36  
`\FAC:` 3, 36–37    `\BINOM:` 3, 36–37

nominator includes  $\backslash C$  factors. If  $a$  is non-negative integer and  $a < b$  then the result is zero because one zero occurs between the factors in the nominator. Thus we give the result zero and we skip the rest of calculation. If  $a$  is non-integer, then  $\backslash C$  must be  $b$ . The  $\backslash\text{step}$  macro (it generates the factors in the nominator) is prepared in two versions: for  $a$  integer we use  $\backslash\text{advance}\backslash A$  by-1 which is much faster than  $\backslash\text{apPLUS}\backslash\text{paramA}\{-1\}$  used for  $a$  non-integer.

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```

745: \def\BINOM#1#2{\relax \apINIT          % BINOM = {#1 \choose #2} ...
746:   \evaldef\apAparam{#1}\apEnum\apAparam
747:   \evaldef\apBparam{#2}\apEnum\apBparam % preparation of the parameters
748:   \localcounts \A \B \C ;%           % local \newcounts
749:   \let\OUT=\apBparam \apROUND\OUT0%   % test if B is integer
750:   \ifx\XOUT\empty\else\apERR{\string\BINOM: second arg. {\apBparam} must be integer}\apRETURN\fi
751:   \let\OUT=\apAparam \apROUND\OUT0%   % test if A is integer
752:   \ifx\XOUT\empty                    % A is integer:
753:     \A=\apAparam \B=\apBparam         %   A = #1, B = #2
754:     \C=\A \advance\C by-\B            %   C = A - B
755:     \ifnum\C>\B \C=\B \fi             %   if (C > B) C = B fi
756:     \ifnum\A<0 \C=\B                  %   if (A < 0) C = B fi
757:     \else \ifnum\A<\B \def\OUT{0}\apSIGN=0 % if (0 <= A < B) OUT = 0 return
758:       \expandafter\expandafter\expandafter \apRETURN \fi\fi
759:     \def\step{\advance\A by-1 \apMUL\OUT{\the\A}}%
760:   \else \C=\apBparam                  % A is not integer
761:     \def\step{\let\apBparam\OUT \do\apAparam=\apPLUS\apAparam{-1}};%
762:     \let\OUT=\apBparam \apMUL\OUT\apAparam}%
763:   \fi
764:   \ifnum\C=0 \def\OUT{1}\apSIGN=1 \apRETURN\fi
765:   \do\D=\FAC{\the\C};%               %   D = C!
766:   \let\OUT=\apAparam                  %   OUT = #1
767:   \loop \advance\C by-1               %   loop C--
768:     \ifnum\C>0 \step \repeat          %   if (C > 0) A--, OUT = OUT * A, repeat
769:   \apDIV{\OUT}{\D}%                  %   OUT = OUT / D
770:   \apEND
771: }

```

The square root is computed in the macro  $\backslash\text{SQRT}\{a\}$  using Newton's approximation method. This method solves the equation  $f(x) = 0$  (in this case  $x^2 - a = 0$ ) by following way. Guess the initial value of the result  $x_0$ . Create tangent to the graph of  $f$  in the point  $[x_0, f(x_0)]$  using the knowledge about  $f'(x_0)$  value. The intersection of this line with the axis  $x$  is the new approximation of the result  $x_1$ . Do the same with  $x_1$  and find  $x_2$ , etc. If you apply the general Newton method to the problem  $x^2 - a = 0$  then you get the formula

$$\text{choose } x_0 \text{ as an initial guess, iterate: } x_{n+1} = \frac{1}{2} \left( x_n + \frac{a}{x_n} \right)$$

If  $|x_{n+1} - x_n|$  is sufficiently small we stop the processing. In practice, we stop the processing, if the  $\backslash\text{OUT}$  representation of  $x_{n+1}$  rounded to the  $\backslash\text{apFRAC}$  is the same as the previous representation of  $x_n$ , i.e.  $\backslash\text{ifx}\backslash\text{Xn}\backslash\text{OUT}$  in  $\text{T}_\text{E}\text{X}$  language. Amazingly, we need only about four iterations for 20-digits precision and about seven iterations for 50-digits precision, if the initial guess is good chosen.

The rest of the work in the  $\backslash\text{SQRT}$  macro is about the right choose of the initial guess (using  $\backslash\text{apSQRTtr}$  macro) and about shifting the decimal point in order to set the  $a$  value into the interval  $[1, 100)$ . The decimal point is shifted by  $-\backslash M$  value. After calculation is done, the decimal point is shifted by  $\backslash M/2$  value back. If user know good initial value then he/she can set it to  $\backslash\text{apSQRTxo}$  macro. The calculation of initial value  $x_0$  is skipped in such case.

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```

772: \def\SQRT#1{\relax \apINIT          % OUT = SQRT(#1) ...
773:   \evaldef\A{#1}%                  % parameter preparation
774:   \localcounts \M \E ;%           % local counters
775:   \E=\apE \apE=0
776:   \ifnum\apSIGN=0 \apRETURN\fi      % SQRT(0) = 0 (OUT is set to 0 by previous \evaldef)
777:   \ifnum\apSIGN<0 \apERR{\string\SQRT: argument {\A} is out of range}\apRETURN\fi
778:   \ifodd\E \apROLL\A{-1}\advance\E by1 \fi % we need the E representation with even exponent
779:   \let\B=\A \let\C=\A

```

$\backslash\text{SQRT}$ : 3, 37, 42, 45     $\backslash\text{apSQRTxo}$ : 38, 42



```

780: \apDIG\C\relax \M=\apnumD % M is the number of digits before decimal point
781: \advance\M by-2 \ifodd\M \advance\M by1 \fi % M = M - 2 , M must be even
782: \ifx\apSQRTxo\undefined % we need to calculate Xo
783: \ifnum\M=0 \else \apROLL\B{-\M}\divide\M by2 \fi % shift decimal point by -M, M = M / 2
784: \apSQRTTr\B \let\Xn=\OUT % Xn = estimate of SQRT
785: \ifnum\M<0 \let\A=\B \fi % if (A < 1) calculate with B where decimal point is shifted
786: \ifnum\M>0 \apROLL\Xn \M \fi % if (A >= 100) shift the decimal point of initial guess
787: \else \let\Xn=\apSQRTxo \fi
788: \loop % loop ... Newton's method
789: \apDIV{\apPLUS{\Xn}{\apDIV{\A}{\Xn}}}{2}% OUT = (Xn + A/Xn) / 2
790: \ifx\OUT\Xn \else % if (OUT != Xn)
791: \let\Xn=\OUT \repeat % Xn = OUT, repeat
792: \ifnum\M<0 \apROLL\OUT\M \fi % shift the decimal point by M back
793: \apE=\E \divide\apE by2 % correct the E exponent
794: \apEND
795: }

```

Note that if the input  $a < 1$ , then we start the Newton's method with  $b$ . It is the value  $a$  with shifted decimal point,  $b \in [1, 100)$ . On the other hand, if  $a \geq 1$  then we start the Newton's method directly with  $a$ , because the second derivative  $(x^2)''$  is constant so the speed of Newton's method is independent on the value of  $x$ . And we need to calculate the `\apFRAC` digits after the decimal point.

The macro `\apSQRTTr`  $\langle number \rangle$  excepts  $\langle number \rangle$  in the interval  $[1, 100]$  and makes a roughly estimation of square root of the  $\langle number \rangle$  in the `\OUT` macro. It uses only classical  $\langle dimen \rangle$  calculation, it doesn't use any `apnum.tex` operations. The result is based on the linear approximation of the function  $g(x) = \sqrt{x}$  with known exact points  $[1, 1], [4, 2], [9, 3], \dots, [100, 10]$ . Note, that the differences between  $x_i$  values of exact points are 3, 5, 7,  $\dots$ , 19. The inverted values of these differences are pre-calculated and inserted after `\apSQRTTra` macro call.

The `\apSQRTTra` macro operates repeatedly for  $i = 1, \dots, 10$  until  $\dimen0 = x < x_i$ . Then the `\apSQRTTrb` is executed. We are in the situation  $\dimen0 = x \in [x_{i-1}, x_i)$ ,  $g(x_i) = i$ ,  $g(x_{i-1}) = i - 1$  and the calculation of  $\OUT = g(x_{i-1}) + (x - x_{i-1}) / (x_i - x_{i-1})$  is performed. If  $x \in [1, 4)$  then the linear approximation is worse. So, we calculate additional linear correction in  $\dimen1$  using the pre-calculated value  $\sqrt{2} - 1.33333 \div 0.08088$  here.

```

796: \def\apSQRTTr#1{\dimen0=#1pt \apnumB=1 \apnumC=1 \apSQRTTra}
797: \def\apSQRTTra{\advance\apnumB by2 \advance\apnumC by\apnumB % B = difference, C = x_i
798: \ifnum\apnumC>100 \def\OUT{10}\else
799: \ifdim\dimen0<\apnumC pt \apSQRTTrb \else
800: \expandafter\expandafter\expandafter\apSQRTTra\fi\fi
801: }
802: \def\apSQRTTrb{% x = \dimen0, B = x_i - x_{i-1}, C = x_i = i
803: \ifdim\dimen0<4pt
804: \ifdim\dimen0>2pt \dimen1=4pt \advance\dimen1 by-\dimen0 \divide\dimen1 by2
805: \else \dimen1=\dimen0 \advance\dimen1 by-1pt \fi
806: \dimen1=.080884\dimen1 % \dimen1 = additional linear correction
807: \else \dimen1=0pt \fi
808: \advance\apnumC by-\apnumB % C = x_{i-1}
809: \advance\dimen0 by-\apnumC pt % \dimen0 = (x - x_{i-1})
810: \divide\dimen0 by\apnumB % \dimen0 = (x - x_{i-1}) / difference
811: \divide\apnumB by2 % B = i-1 = g(x_{i-1})
812: \advance\dimen0 by\apnumB pt % \dimen0 = g(x_{i-1}) + (x - x_{i-1}) / (x_i - x_{i-1})
813: \advance\dimen0 by\dimen1 % \dimen0 += additional linear correction
814: \edef\OUT{\expandafter\apNOPT\the\dimen0}% OUT = \dimen0
815: }

```

The exponential function  $e^x$  is implemented in the `\EXP` macro using Taylor series at zero point:

$$e^x = 1 + \frac{x}{1!} + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots$$

If  $x \in (0, 1)$  then this series converges relatively quickly.

The macro `\EXP` takes its argument. If it is negative, remember this fact, remove minus sign and do `\OUT=1/\OUT` in final step. Now, the argument is positive always. If the argument is greater than 1,



do argument = argument/2 and increase  $K$  register. Do this step in the loop until argument  $< 1$ . After  $\backslash\text{OUT}$  is calculated then we do  $\backslash\text{OUT}=\backslash\text{OUT}^2$  in the loop  $K$  times, because  $e^{2x} = (e^x)^2$ . Now we are ready to calculate the exponential of positive argument which is less than 1. This is done using loop of Taylor series.  $\backslash\text{S}$  is total sum,  $\backslash\text{Sn}$  is the new addition in the  $n$ -th step. If  $\backslash\text{Sn}$  is zero (in accordance to the  $\backslash\text{apFRAC}$  register) then we stop the calculation.

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```

816: \def\EXP#1{\relax\apINIT % OUT = EXP(#1) ...
817: \evaldef\OUT{#1}\apEnum\OUT % OUT = #1
818: \localcounts \N \K;%
819: \ifnum\apSIGN=0 \def\OUT{1}\apSIGN=1 \apRETURN \fi
820: \edef\digits{\the\apFRAC}\advance\apFRAC by3
821: \edef\signX{\the\apSIGN}%
822: \ifnum\apSIGN<0 \apSIGN=1 \apREMfirst\OUT \fi % remove "minus" sign
823: \K=0 \N=0 % K = 0, N = 0
824: \def\testDot ##1##2\relax##3{\ifx##1.}%
825: \loop \expandafter \testDot\OUT \relax % loop if (OUT >= 1)
826: \iftrue \else % OUT = OUT/2
827: \apDIV\OUT{2}% % K++
828: \advance\K by1 % repeat
829: \repeat % now: #1 = 2^K * OUT, OUT < 1
830: \advance\apFRAC by\K
831: \def\S{1}\def\Sn{1}\N=0 \let\X=\OUT % S = 1, Sn = 1, N = 0, X = OUT
832: \loop \advance\N by1 % loop N++
833: \do\Sn=\apDIV{\apMUL\Sn\X}{\the\N};% % Sn = Sn * X / N
834: \apTAYLOR\iftrue \repeat % S = S + Sn (... Taylor)
835: \N=0
836: \loop \ifnum\N < \K % loop if (N < K)
837: \apPOW\OUT{2}\apROUND\OUT\apFRAC % OUT = OUT^2
838: \advance\N by1 \repeat % N++
839: \ifnum\signX<0 \apDIV 1\OUT \fi % if (signX < 0) OUT = 1 / OUT
840: \apROUND\OUT\digits \apSIGN=1 % EXP is always positive
841: \apEND
842: }

```

The macro  $\backslash\text{apTAYLOR}$  is ready for general usage in the form:

```

\def\S{...}\def\Sn{...}\N=... % setting initial values for N=0
\loop
... % auxiliary calculation
\do\Sn=\apDIV{...}{...};% % calculation of new addition \Sn
% (division must be the last activity)
\apTAYLOR \iftrue \repeat % does S = S + Sn and finishes if Sn = 0

```

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```
843: \def\apTAYLOR#1{\ifnum\apSIGN=0 \let\OUT=\S \else \apPLUS\S\Sn \let\S=\OUT }
```

**The logarithm function**  $\ln x$  (inverse to  $e^x$ ) is implemented in  $\backslash\text{LN}$  macro by Taylor series in the point zero of the arg tanh function:

$$\ln x = 2 \arg \tanh \frac{x-1}{x+1} = 2 \left( \frac{x-1}{x+1} + \frac{1}{3} \left( \frac{x-1}{x+1} \right)^3 + \frac{1}{5} \left( \frac{x-1}{x+1} \right)^5 + \dots \right).$$

This series converges quickly when  $x$  is approximately equal to one. The idea of the macro  $\backslash\text{LN}$  includes the following steps:

- Whole calculation is in the group  $\backslash\text{apINIT}...\backslash\text{apEND}$ . Enlarge the  $\backslash\text{apFRAC}$  numeric precision by three digits in this group.
- Read the argument  $\backslash\text{X}$  using  $\backslash\text{evaldef}$ .
- If the argument is non positive, print error and skip the next processing.
- If the argument is in the interval  $(0, 1)$ , set new argument as  $1/\text{argument}$  and remember the “minus” sign for the calculated  $\backslash\text{OUT}$ , else the  $\backslash\text{OUT}$  remains to be positive. This uses the identity  $\ln(1/x) = -\ln x$ .

---

$\backslash\text{apTAYLOR}$ : 39–40, 43–45     $\backslash\text{LN}$ : 3, 6, 39–40

- shift the decimal point of the argument by  $M$  positions left in order to the new argument is in the interval  $[1, 10)$ .
- Let  $x \in [1, 10)$  be the argument calculated as mentioned before. Calculate roughly estimated  $\widetilde{\ln x}$  using `\apLNr` macro. This macro uses linear interpolation of the function  $\ln x$  in eleven points in the interval  $[1, 10]$ .
- Calculate  $A = x / \exp(\widetilde{\ln x})$ . The result is approximately equal to one, because  $\exp(\ln x) = x$ .
- Calculate  $\ln A$  using the Taylor series above.
- The result of  $\ln x$  is equal to  $\ln A + \widetilde{\ln x}$ , because  $x = A \cdot \exp(\widetilde{\ln x})$  and  $\ln(ab) = \ln a + \ln b$ .
- The real argument is in the form  $x \cdot 10^M$ , so `\OUT` is equal to  $\ln x + M \cdot \ln(10)$  because  $\ln(ab) = \ln a + \ln b$  and  $\ln(10^M) = M \ln(10)$ . The  $\ln(10)$  value with desired precision is calculated by `\apLNtenexec` macro. This macro saves its result globally when firstly calculated and use the calculated result when the `\apLNtenexec` is called again.
- Round the `\OUT` to the `\apFRAC` digits.
- Append “minus” to the `\OUT` if the input argument was in the interval  $(0, 1)$ .

```

845: \def\LN#1{\relax \apINIT           % OUT = LN(#1) ...
846:   \evaldef\X{#1}%                 % X = #1
847:   \localcounts \M \N \E;%
848:   \E=\apE
849:   \def\round{\apROUND\OUT\apFRAC}%
850:   \edef\digits{\the\apFRAC}\advance\apFRAC by4
851:   \ifnum\apSIGN>0 \else \apERR{\string\LN: argument {\X} is out of range}\apRETURN\fi
852:   \apDIG\OUT\relax \M=\apnumD      % find M: X = mantissa * 10^M
853:   \ifnum\M>-\E \def\sgnout{1}\else % if X in (0,1):
854:     \def\sgnout{-1}%               % sgnout = -1
855:     \do\X=\apDIV 1\X;\E=-\E        % X = 1/X
856:     \apDIG\OUT\relax \M=\apnumD      % find M: X = mantissa * 10^M
857:   \fi                               % else sgnout = 1
858:   \advance\M by-1                   % M = M - 1
859:   \ifnum\M=0 \else\apROLL\X{-\M}\fi % X = X * 10^(-M), now X in (1,10)
860:   \advance\M by\E                   % M = M + E (scientific format of numbers)
861:   \do\lnX=\apLNr\X;%               % lnX = LN(X) ... roughly estimate
862:   \do\A=\apDIV\X{\EXP\lnX};%       % A = X / EXP(lnX) ... A=approx= 1
863:   \apLNtaylor                      % OUT = LN(A)
864:   \do\LNOUT=\apPLUS\OUT\lnX;%       % LNOUT = OUT + LNrOUT
865:   \ifnum\M>0                       % if M > 0
866:     \apLNtenexec                   % LNtenOUT = ln(10)
867:     \apPLUS\LNOUT{\apMUL{\the\M}{\apLNten}}% OUT = LNOUT + M * LNten
868:   \fi
869:   \ifnum\apSIGN=0 \else \apSIGN=\sgnout \fi % if (OUT != 0) apSIGN = saved sign
870:   \apROUND\OUT\digits               % round result to desired precision
871:   \ifnum\apSIGN<0 \xdef\OUT{-\OUT}\else \global\let\OUT=\OUT \fi
872:   \apEND
873: }

```

The macro `\apLNtaylor` calculates  $\ln A$  for  $A \approx 1$  using Taylor series mentioned above.

```

874: \def\apLNtaylor{%
875:   \apDIV{\apPLUS{\A}{-1}}{\apPLUS{\A}{1}}% % OUT = (A-1) / (A+1)
876:   \ifnum\apSIGN=0 \def\OUT{0}\else          % ln 1 = 0 else:
877:     \let\Sn=\OUT \let\Kn=\OUT \let\S=\OUT    % Sn = OUT, Kn = OUT, S = OUT
878:     \apPOW\OUT{2}\round \let\XX=\OUT         % XX = OUT^2
879:     \N=1                                       % N = 1
880:     \loop \advance\N by2                     % loop N = N + 2
881:       \do\Kn=\apMUL\Kn\XX\round;%            % Kn = Kn * XX
882:       \do\Sn=\apDIV\Kn{\the\N};%            % Sn = Kn / N
883:       \apTAYLOR\iftrue \repeat              % S = S + Sn (Taylor)
884:       \apMUL\S{2}%                           % OUT = 2 * OUT
885:   \fi
886: }

```

The macro `\apLNr` finds an estimation  $\widetilde{\ln x}$  for  $x \in [1, 10)$  using linear approximation of  $\ln x$  function. Only direct  $\langle \text{dimen} \rangle$  and  $\langle \text{count} \rangle$  calculation with  $\text{T}_{\text{E}}\text{X}$  registers is used, no long numbers `apnum.tex` calculation. The  $\ln x_i$  is pre-calculated for  $x_i = i$ ,  $i \in \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  and the values are inserted after the `\apLNra` macro call. The input value  $x$  is set as `\dimen0`.

The `\apLNra`  $\{\langle \text{valueA} \rangle\}\{\langle \text{valueB} \rangle\}$  macro reads the pre-calculated values repeatedly in the loop. The loop ends if `\apnumC` (i. e.  $x_i$ ) is greater than  $x$ . Then we know that  $x \in [x_{i-1}, x_i)$ . The linear interpolation is

$$\widetilde{\ln x} = f(x_{i-1}) + (f(x_i) - f(x_{i-1}))(x - x_{i-1}),$$

where  $f(x_{i-1}) = \langle \text{valueA} \rangle$ ,  $f(x_i) = \langle \text{valueB} \rangle$  and  $x = \text{\dimen0}$ . The rest of the pre-calculated values is skipped by processing `\next` to `\relax`.

The pre-calculated approximation of  $\ln 10$  is saved in the macro `\apLNrten` because we use it at more places in the code.

```

887: \def\apLNr#1{\dimen0=#1pt \apnumC=1
888:   \apLNra {0}{.69}{1.098}{1.386}{1.609}{1.791}{1.9459}{2.079}{2.197}{\apLNrten}{\relax
889: }
890: \def\apLNra #1#2{\advance\apnumC by1
891:   \ifx\relax#2\relax \let\OUT=\apLNrten \let\apNext=\relax
892:   \else
893:     \ifdim\dimen0<\apnumC pt % linear interpolation:
894:       \advance\dimen0 by-\apnumC pt \advance\dimen0 by1pt % \dimen0 = x - x_{i-1}
895:       \dimen1=#2pt \advance\dimen1 by-#1pt % \dimen1 = f(x_i) - f(x_{i-1})
896:       \dimen1=\expandafter\apNOPT\the\dimen0 \dimen1 % \dimen1 = (x - x_{i-1}) * \dimen1
897:       \advance\dimen1 by#1pt % \dimen1 = f(x_{i-1}) + \dimen1
898:       \edef\OUT{\expandafter\apNOPT\the\dimen1}% % OUT = \dimen1
899:       \def\apNext#1\relax{}%
900:     \else \def\apNext{\apLNra{#2}}%
901:     \fi\fi \apNext
902: }
903: \def\apLNrten{2.302585} % \apLNrten = ln 10 (roughly)

```

The `\apLNtenexec` macro calculates the  $\ln 10$  value with the precision given by `\apFRAC`. The output is prepared to the `\apLNten` macro. The `\apLNtenexec` saves globally the result to the macro `\LNten`: $\langle \text{apFRAC} \rangle$  in order to use it if the value is needed again. This saves time.

```

904: \def\apLNtenexec{%
905:   \expandafter\ifx\csname LNten:\the\apFRAC\endcsname \relax
906:     \begingroup \apTOT=0
907:     \do\A=\apDIV{10}{\EXP\apLNrten};% % A = 10 / exp(LNrten)
908:     \apLNtaylor % OUT = ln A
909:     \apPLUS\OUT\apLNrten % OUT = OUT + LNrten
910:     \global\expandafter\let\csname LNten:\the\apFRAC\endcsname=\OUT
911:   \endgroup
912:   \fi
913:   \expandafter\let\expandafter \apLNten \csname LNten:\the\apFRAC\endcsname
914: }

```

The constant  $\pi$  is saved in the `\apPIvalue` macro initially with 30 digits. If user needs more digits (using `\apFRAC > 30`) then the `\apPIvalue` is recalculated and the `\apPIDigits` is changed appropriately.

```

915: \def\apPIvalue{3.141592653589793238462643383279}
916: \def\apPIDigits{30}

```

The macro `\apPIexec` prepares the  $\pi$  constant with `\apFRAC` digits and saves it to the `\apPI` macro. And  $\pi/2$  constant with `\apFRAC` digits is saved to the `\apPIhalf` macro. The `\apPIexec` uses macros `\apPI`: $\langle \text{apFRAC} \rangle$  and `\apPIh`: $\langle \text{apFRAC} \rangle$  where desired values are usually stored. If the values are not prepared here then the macro `\apPIexecA` calculates them.

---

`\apLNr`: 40–41    `\apLNra`: 41    `\apLNrten`: 41    `\apLNtenexec`: 40–41    `\apLNten`: 40–41  
`\apPIvalue`: 41–43    `\apPIDigits`: 41–43    `\apPIexec`: 41–46    `\apPI`: 42–44    `\apPIhalf`: 42–46

```

917: \def\apIexec{%
918:   \expandafter\ifx\csname apPI:\the\apFRAC\endcsname \relax \apIexecA \else
919:     \expandafter\let\expandafter\apPI\csname apPI:\the\apFRAC\endcsname
920:     \expandafter\let\expandafter\apPIhalf\csname apPIh:\the\apFRAC\endcsname
921:   \fi
922: }

```

The macro `\apIexecA` creates the  $\pi$  value with `\apFRAC` digits using the `\apPIvalue`, which is rounded if `\apFRAC < \apPIDigits`. The `\apPIhalf` is calculated from `\apPI`. Finally the macros `\apPI:\langle apFRAC \rangle` and `\apPIh:\langle apFRAC \rangle` are saved globally for saving time when we need such values again.

```

923: \def\apIexecA{%
924:   \ifnum\apPIDigits<\apFRAC \apIexecB \fi
925:   \let\apPI=\apPIvalue
926:   \ifnum\apPIDigits>\apFRAC \apROUND\apPI\apFRAC \fi
927:   \apnumP=\apTOT \apTOT=0 \apDIV\apPI2\let\apPIhalf=\OUT \apTOT=\apnumP
928:   \global\expandafter\let\csname apPI:\the\apFRAC\endcsname=\apPI
929:   \global\expandafter\let\csname apPIh:\the\apFRAC\endcsname=\apPIhalf
930: }

```

If `\apFRAC > \apPIDigits` then new `\apPIvalue` with desired decimal digits is generated using `\apIexecB` macro. The Chudnovsky formula is used:

$$\pi = \frac{53360 \cdot \sqrt{640320}}{S}, \quad S = \sum_{n=0}^{\infty} \frac{(6n)! (13591409 + 545140134n)}{(3n)! (n!)^3 (-262537412640768000)^n}$$

This converges very good with 14 new calculated digits per one step where new  $S_n$  is calculated. Moreover, we use the identity:

$$F_n = \frac{(6n)!}{(3n)! (n!)^3 (-262537412640768000)^n}, \quad F_n = F_{n-1} \cdot \frac{8(6n-1)(6n-3)(6n-5)}{n^3 (-262537412640768000)}$$

and we use auxiliary integer constants  $A_n, B_n, C_n$  with following properties:

$$\begin{aligned}
A_0 &= B_0 = C_0 = 1, \\
A_n &= A_{n-1} \cdot 8(6n-1)(6n-3)(6n-5), \quad B_n = B_{n-1} \cdot n^3, \quad C_n = C_{n-1} \cdot (-262537412640768000), \\
F_n &= \frac{A_n}{B_n C_n}, \\
S_n &= \frac{A_n (13591409 + 545140134n)}{B_n C_n}
\end{aligned}$$

```

931: \def\apIexecB{\apINIT
932:   \localcounts \N \a \c;%
933:   \apTOT=0 \advance\apFRAC by2
934:   \def\apSQRTxo{800.199975006248}% initial value for Newton method for SQRT
935:   \SQRT{640320}%
936:   \let\sqrval=\OUT
937:   \N=0 \def\An{1}\def\Bn{1}\def\Cn{1}\def\S{13591409}%
938:   \loop
939:     \advance\N by 1
940:     \a=\N \multiply\c by6 \advance\c by-1 \c=\a
941:     \advance\c by-2 \multiply\c by\c \a = \a * 8 * (6N-5) *
942:     \advance\c by-2 \multiply\c by8 \a = \a * (6N-3) * (6N-1)
943:     \apMUL\An{\apMUL{\the\c}{\the\c}}\let\An=\OUT
944:     \c=\N \multiply\c by\N \a = \a * n^3
945:     \apMUL\Bn{\apMUL{\the\c}{\the\N}}\let\Bn=\OUT
946:     \apMUL\Cn{-262537412640768000}\let\Cn=\OUT \a = \a * K3
947:     \apDIV\apMUL\An{\apPLUS{13591409}{\apMUL{545140134}{\the\N}}}\apMUL\Bn\Cn}%
948:     \let\S=\OUT \a = \a * (K1 + K2 * N) / (Bn * Cn)

```

`\apIexecA: 41–42`    `\apIexecB: 42`

```

949: \apTAYLOR \iftrue \repeat
950: \advance\apFRAC by-2
951: \apDIV{\apMUL{\sqrtval}{53360}}\S
952: \global\let\apPIvalue=\OUT
953: \xdef\apPIdigits{\the\apFRAC}%
954: \apEND
955: }

```

The macros for users `\PI` and `\PIhalf` are implemented as “function-like” macros without parameters.

```

956: \def\PI{\relax \apPIexec \let\OUT=\apPI}
957: \def\PIhalf{\relax \apPIexec \let\OUT=\apPIhalf}

```

apnum.tex

The macros `\SIN` and `\COS` use the Taylor series

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots$$

These series converge good for  $|x| < 1$ . The main problem is to shift the given argument  $x \in \mathbf{R}$  to the range  $[0, 1)$  before the calculation of the series is started. This task is done by `\apSINCOSa` macro, the common code for both, `\SIN` and `\COS` macros.

The macro `\apSINCOSa` does the following steps:

- It advances `\apFRAC` by three and evaluates the argument.
- Note, that the macro `\apSINCOSx` means `\apSINx` or `\apCOSx` depending on the given task.
- The macro `\signK` includes 1. It can be recalculated to -1 later.
- If the argument is zero then the result is set and next computation is skipped. This test is processed by `\apSINCOSo\apCOSx`.
- If the argument is negative then remove minus and save `\sign`. This `\sign` will be applied to the result. The `\sign` is always + when `\COS` is calculated. This follows the identities  $\sin(-x) = -\sin x$  and  $\cos(-x) = \cos x$ .
- The `\apFRAC` is saved and `\apTOT=0`.
- The `\apPIexec` is processed. The `\apPI` and `\apPIhalf` are ready after such processing.
- After `\X div \apPI` (rounded to integer) we have `\K` in `\OUT`, where  $\X = x' + \K \cdot \pi$  and  $x' \in [0, \pi)$ . We set  $\X := x'$  because of the identities  $\sin x = (-1)^k \sin(x + k\pi)$ ,  $\cos x = (-1)^k \cos(x + k\pi)$ . The sign  $(-1)^k$  is saved to `\signK` macro.
- If the  $x'$  is zero then the result is set by `\apSINCOSo\apCOSx` and the rest of calculating is skipped.
- The  $|\X - \pi/2|$  is saved to `\XmPIh` macro.
- If  $\X \in (\pi/4, \pi/2)$  then  $x' = \XmPIh$ . We use identities  $\sin x = \cos(\pi/2 - x)$ ,  $\cos x = \sin(\pi/2 - x)$ . Set  $\X = x'$ . The meaning of `\apSINCOSx` (`\apSINx` or `\apCOSx`) is flipped in such case.
- If the  $x'$  is zero then the result is set by `\apSINCOSo\apSINx` and the rest of calculating is skipped.
- Now  $\X \in (0, \pi/4)$ , i. e.  $|\X| < 1$  and we can use Taylor series. The `\apSINCOSx` (i.e. `\apSINx` or `\apCOSx`) macro initializes the computation of Taylor series mentioned above. The `\XX = \X^2` is prepared. The Taylor series is processed in the loop as usually.
- The the sign of the output is `\sign\signK`.
- If the sign of the result is negative, the “minus” is added to the `\OUT`.

```

959: \def\SIN{\relax \let\apSINCOSx=\apSINx \apSINCOSa}
960: \def\COS{\relax \let\apSINCOSx=\apCOSx \apSINCOSa}
961: \def\apSINCOSa#1{\apINIT
962: \advance\apFRAC by3
963: \evaldef\X{#1}\apEnum\X
964: \def\signK{1}\apSINCOSo\apCOSx
965: \ifnum\apSIGN<0 \apREMfirst\X \def\sign{-}\else\def\sign{+}\fi
966: \ifx\apSINCOSx\apCOSx \def\sign{+}\fi
967: \edef\apFRACsave{\the\apFRAC}%

```

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---

`\PI`: 3, 43    `\PIhalf`: 3, 43    `\SIN`: 3, 43–44    `\COS`: 3, 43–44    `\apSINCOSa`: 43–44

```

968: \apIexec
969: \apFRAC=0 \apDIV\X\apPI          % OUT = X div PI
970: \ifnum\apSIGN=0 \apSIGN=1 \else
971:   \let\K=\OUT
972:   \do\X=\apPLUS\X{-\apMUL\K\apPI};% X := X - K * PI
973:   \apROLL\K{-1}\apROUND\K{0}%
974:   \ifodd 0\XOUT\space \def\signK{-1}\else\def\signK{1}\fi
975: \fi
976: \apSINCOSo\apCOSx
977: \apFRAC=\apFRACsave \relax
978: \do\XmPIh=\apPLUS\X{-\apPIhalf};% XmPIh = | X - PI/2 |
979: \apSINCOSo\apSINx
980: \ifnum\apSIGN<0 \apREMfirst\XmPIh
981: \else % X in (PI/2, PI)
982:   \do\X=\apPLUS\apPI{-\X};%
983:   \ifx\apSINCOSx\apCOSx \apSIGN=-\signK \edef\signK{\the\apSIGN}\fi
984:   \fi % X in (0, PI/2):
985:   \apMINUS\X{.78}% % OUT = X - cca PI/4
986:   \ifnum\apSIGN<0 \else % if X in (PI/4, PI/2) :
987:     \let\X=\XmPIh % X = | X - PI/2 |; SIN <-> COS
988:     \ifx\apSINCOSx\apSINx \let\apSINCOSx=\apCOSx \else \let\apSINCOSx=\apSINx \fi
989:   \fi
990:   \localcounts \N \NN;%
991:   \do\XX=\apPOW\X{2}\ROUND\OUT\apFRAC;%
992:   \apSINCOSx % X in (0, PI/4), initialize Taylor SIN X or COS X
993:   \loop
994:     \advance\N by1 \NN=\N
995:     \advance\N by1 \multiply\NN by\N
996:     \do\Sn=\apDIV{\apMUL\Sn\XX}{-\the\NN};% Sn = - Sn * X^2 / N*(N+1)
997:     \apTAYLOR \iftrue\repeat
998:     \apSIGN=\sign\signK
999:     \ifnum\apTOT=0 \advance\apFRAC by-3 \else \apFRAC=\apTOT \fi
1000:     \ifnum\apFRAC<0 \apFRAC=-\apFRAC \fi
1001:     \apROUND\OUT\apFRAC
1002:     \ifnum\apSIGN<0 \edef\OUT{-\OUT}\fi
1003:     \apEND
1004: }

```

The macros `\apSINx` and `\apCOSx` initialize the calculation of the Taylor series.

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```

1005: \def\apSINx{\let\S=X \N=1 \let\Sn=X}
1006: \def\apCOSx{\def\S{1}\N=0 \let\Sn=S}

```

The `\apSINCOSo` *sequence* macro is used three times in the `\apSINCOSa`. It tests if the current result is zero. If it is true then the `\OUT` is set as zero or it is set to `\signK` (if processed function is equal to the *sequence*).

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```

1007: \def\apSINCOSo#1{\ifnum\apSIGN=0 \ifx#1\SCgo \apSIGN=\signK \let\OUT=\signK \fi \apRETURN\fi}

```

The macro `\TAN` uses the identity  $\tan x = \sin x / \cos x$  and calculates the denominator first. If it is zero then `\apERR` prints “out of range” message else the result is calculated.

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```

1008: \def\TAN#1{\relax \apINIT
1009:   \advance\apFRAC by3
1010:   \evaldef\X{#1}\apEnum\X
1011:   \advance\apFRAC by-3
1012:   \do\denom=\COS\X;%
1013:   \ifnum\apSIGN=0 \apERR{\string\TAN: argument {\X} is out of range}\apRETURN\fi
1014:   \SIN\X\message{\OUT/\denom}%
1015:   \apDIV{\SIN\X}\denom
1016:   \apEND
1017: }

```



The macro `\ATAN` calculates the inverse of tangens using series

$$\arctan \frac{1}{x} = \frac{x}{1+x^2} + \frac{2}{3} \frac{x}{(1+x^2)^2} + \frac{2}{3} \frac{4}{5} \frac{x}{(1+x^2)^3} + \frac{2}{3} \frac{4}{5} \frac{6}{7} \frac{x}{(1+x^2)^4} + \dots$$

This converges relatively good for  $|x| > 1$ . I was inspired by the Claudio Kozický's semestral work from the course "Typography and T<sub>E</sub>X" at ČVUT in Prague.

The macro `\ATAN` takes the argument  $x$  and uses identity  $\arctan(-x) = -\arctan(x)$  when  $x$  is negative. If  $x > 1$  then the identity

$$\arctan(x) = \frac{\pi}{2} - \arctan \frac{1}{x}$$

is used and  $\arctan(1/x)$  is calculated by `\apATANox` macro using the series above. Else the argument is re-calculated  $x := 1/x$  and the `\apATANox` is used. When  $x = 1$  then the `\apPIhalf/2` is returned directly.

```

1018: \def\ATAN#1{\relax \apINIT
1019:   \advance\apFRAC by3
1020:   \evaldef\X{#1}\apEnum
1021:   \ifnum\apSIGN=0 \def\OUT{0}\apRETURN\fi
1022:   \ifnum\apSIGN<0 \def\sign{-}\apREMfirst\X \else\def\sign{}\fi
1023:   \let\tmp=\X \apDIG\tmp\relax
1024:   \ifnum\apnumD>0 % if X > 1:
1025:     \apIexec % OUT = apPIhalf - apATANox
1026:     \def\tmp{1}\ifx\tmp\X \apDIV\apPIhalf2\else \apATANox \apPLUS\apPIhalf{-\OUT}\fi
1027:   \else % else
1028:     \do\X=\apDIV{1}\X;% X := 1/X
1029:     \apATANox % OUT = apATANox
1030:   \fi
1031:   \ifnum\apTOT=0 \advance\apFRAC by-3 \else \apFRAC=\apTOT \fi
1032:   \ifnum\apFRAC<0 \apFRAC=-\apFRAC \fi
1033:   \apROUND\OUT\apFRAC
1034:   \ifx\sign\empty\apSIGN=1 \else \edef\OUT{-\OUT}\apSIGN=-1 \fi
1035:   \apEND
1036: }

```

The macro `\apATANox` calculates  $\arctan(1/x)$  using series mentioned above.

```

1037: \def\apATANox{%
1038:   \localcounts \N;%
1039:   \do\XX=\apPLUS{1}{\apPOW\X{2}}\apROUND\OUT\apFRAC;% XX = 1 + X^2
1040:   \do\Sn=\apDIV\X\XX \apROUND\OUT\apFRAC;% % Sn = X / (1+X^2)
1041:   \N=1 \let\S=\Sn
1042:   \loop
1043:     \advance\N by1
1044:     \do\Sn=\apMUL{\the\N}\Sn;%
1045:     \advance\N by1
1046:     \do\Sn=\apDIV\Sn{\apMUL{\the\N}\XX};% Sn = Sn * N / ((N+1) * (1+X^2))
1047:     \apTAYLOR \iftrue \repeat
1048: }

```

The macros `\ASIN` and `\ACOS` for functions  $\arcsin(x)$  and  $\arccos(x)$  are implemented using following identities:

$$\arcsin(x) = \arctan \frac{x}{\sqrt{1-x^2}}, \quad \arccos(x) = \frac{\pi}{2} - \arcsin(x)$$

```

1049: \def\ASIN#1{\relax \apINIT
1050:   \evaldef\X{#1}\apEnum\X \edef\sign{\the\apSIGN}%
1051:   \apPLUS 1{-\apPOW\X2}% OUT = 1 - X^2
1052:   \ifnum\apSIGN<0 \apERR{\string\ASIN: argument {\X} is out of range}\apRETURN\fi
1053:   \do\sqr=\SQRT\OUT;% sqrt = SRQT {1 - X^1}
1054:   \ifnum\apSIGN=0 \apIexec

```

```

1055:      \ifnum\sign<0 \edef\OUT{-\apPIhalf}\apSIGN=-1 % ASIN(-1) = -PI/2
1056:      \else \let\OUT=\apPIhalf \apSIGN=1 \fi % ASIN(1) = PI/2
1057:      \apRETURN \fi
1058:      \ATAN{X/\sqrt{}}% OUT = arctan ( X / SQRT {1 - X^2} )
1059:      \apEND
1060: }
1061: \def\ACOS#1{\relax \apPiexec \apPLUS\apPIhalf{-\ASIN{#1}}}

```

## 2.11 Conclusion

This code is here only for backward compatibility with old versions of `apnum.tex`. Don't use these sequences if you are implementing an internal feature because users can re-define these sequences.

```

1066: \let\PLUS=\apPLUS \let\MINUS=\apMINUS \let\MUL=\apMUL \let\DIV=\apDIV \let\POW=\apPOW
1067: \let\SIGN=\apSIGN \let\ROUND=\apROUND \let\NORM=\apNORM \let\ROLL=\apROLL

```

`apnum.tex`

Here is my little joke. Of course, this macro file works in LaTeX without problems because only TeX primitives (from classical TeX) and the `\newcount` macro are used here. But I wish to print my opinion about LaTeX. I hope that this doesn't matter and LaTeX users can use my macro because a typical LaTeX user doesn't read a terminal nor `.log` file.

```

1069: \ifx\documentclass\undefined \else % please, don't remove this message
1070: \message{SORRY, you are using LaTeX. I don't recommend this. Petr Olsak}\fi
1071: \catcode'\@=\apnumZ
1072: \endinput

```

`apnum.tex`

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