## expkV

# an expandable $\langle$ key $\rangle=\langle$ value $\rangle$ implementation 

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2020-04-10 V1.2


#### Abstract

expk ${ }^{v}$ provides a small interface for $\langle k e y\rangle=\langle$ value $\rangle$ parsing. The parsing macro is fully expandable, the 〈code〉 of your keys might be not. expkv is pretty fast, but not the fastest available $\langle k e y\rangle=\langle$ value $\rangle$ solution (keyval is one and a half times as fast, but not expandable and it might strip braces it shouldn't have stripped).


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## 1 Documentation

expkv provides an expandable angle\)key$\rangle=\langle$value$\rangle$parser.The$\langle$key$\rangle=\langle$value$\rangle$pairsshouldbegivenasacommaseparatedlistandtheseparatorbetweena$\langlekey\rangle$andtheassociated<value〉shouldbeanequalsign.Both,thecommasandtheequalsigns,mightbeofcategory12(other)or13(active).Tosupportthisisnecessaryasforexamplebabelturnscharactersactiveforsomelanguages,forinstancetheequalsignisturnedactiveforTurkish.expkvisusableasgenericcodeorasa${}^{\mathrm{EA}}\mathrm{T}_{\mathrm{E}}\mathrm{X}$package.Touseit,justuseoneof:$\backslash$usepackage\{expkv\}\%LaTeX\inputexpkv\%plainTeXThe$\mathrm{IAT}_{\mathrm{E}}\mathrm{X}$packagedoesn'tdomorethanexpkv.tex,exceptcalling$\backslash$ProvidesPackageandsettingthingsupsuchthatexpkv.texwilluse\ProvidesFile.undefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefined

In the expkv family are other packages contained which provide additional functionality. Those packages currently are:
expklidef a key-defining frontend for expkv using a $\langle$ key $\rangle=\langle$ value $\rangle$ syntax
expkvics define expandable $\langle$ key $\rangle=\langle$ value $\rangle$ macros using expkv
Note that while the package names are stylised with a vertical rule, their names are all lower case with a hyphen (e.g., expkv-def).

### 1.1 Setting up Keys

expkv provides a rather simple approach to setting up keys, similar to keyval. However there is an auxiliary package named expklidef which provides a more sophisticated interface, similar to well established packages like pgfkeys or l3keys.

Keys in expkv (as in almost all other $\langle k e y\rangle=\langle v a l u e\rangle$ implementations) belong to a set such that different sets can contain keys of the same name. Unlike many other implementations expk doesn't provide means to set a default value, instead we have keys that take values and keys that don't (the latter are called NoVal keys by expkv), but both can have the same name (on the user level).

The following macros are available to define new keys. Those macros containing "def" in their name can be prefixed by anything allowed to prefix \def (but don't use \outer, keys defined with it won't ever be usable), prefixes allowed for \let can prefix those with "let" in their name, accordingly. Neither $\langle\operatorname{set}\rangle$ nor $\langle k e y\rangle$ are allowed to be empty for new keys and must not contain a \par or tokens that expand to it - they must be legal inside of \csname . . . \endcsname.
\ekvdef $\backslash \operatorname{ekvdef}\{\langle$ set $\rangle\}\{\langle$ key $\rangle\}\{\langle$ code $\rangle\}$
Defines a $\langle k e y\rangle$ taking a value in a $\langle$ set $\rangle$ to expand to $\langle$ code $\rangle$. In $\langle$ code $\rangle$ you can use \#1 to refer to the given value.
\ekvdefNoVal \ekvdefNoVal\{ $\{$ set $\rangle\}\{\langle$ key $\rangle\}\{\langle$ code $\rangle\}$
Defines a no value taking $\langle$ key $\rangle$ in a $\langle$ set $\rangle$ to expand to $\langle$ code $\rangle$.
\ekvlet $\backslash \operatorname{ekvlet}\{\langle$ set $\rangle\}\{\langle$ key $\rangle\}\langle$ cs $\rangle$
Let the value taking $\langle k e y\rangle$ in $\langle s e t\rangle$ to $\langle c s\rangle$, there are no checks on $\langle c s\rangle$ enforced.
\ekvletNoVal \ekvletNoVal\｛〈set $\rangle\}\{\langle k e y\rangle\}\langle c s\rangle$
Let the no value taking $\langle k e y\rangle$ in $\langle$ set $\rangle$ to $\langle c s\rangle$ ，it is not checked whether $\langle c s\rangle$ exists or that it takes no parameter．
\ekvletkv \ekvletkv\｛〈set $\rangle\}\{\langle$ key $\rangle\}\{\langle$ set 2$\rangle\}\{\langle$ key 2$\rangle\}$
Let the $\langle$ key $\rangle$ in $\langle$ set $\rangle$ to $\langle$ key2 $\rangle$ in $\langle$ set2 $\rangle$ ，it is not checked whether that second key exists．
\ekvletkvNoVal\｛〈set $\rangle\}\{\langle$ key $\rangle\}\{\langle$ set2 $\rangle\}\{\langle$ key2 $\rangle\}$
Let the $\langle$ key $\rangle$ in $\langle s e t\rangle$ to $\langle k e y 2\rangle$ in $\langle$ set2 $\rangle$ ，it is not checked whether that second key exists．

## 1．2 Parsing Keys

\ekvset
\ekvset $\{\langle$ set $\rangle\}\{\langle$ key $\rangle=\langle$ value $\rangle, \ldots\}$
Splits $\langle$ key $\rangle=\langle$ value $\rangle$ pairs on commas．From both $\langle$ key $\rangle$ and $\langle$ value $\rangle$ up to one space is stripped from both ends，if then only a braced group remains the braces are stripped as well．So \ekvset $\{f \circ o\}\{b a r=b a z\}$ and $\backslash e k v s e t\{f o o\}\{$ bar\} $=\{b a z\}\}$ will both do $\backslash$ foobarcode $\rangle\{$ baz $\}$ ，so you can hide commas，equal signs and spaces at the ends of either $\langle k e y\rangle$ or 〈value〉 by putting braces around them．If you omit the equal sign the code of the key created with the NoVal variants described in subsection 1.1 will be executed．If $\langle$ key $\rangle=\langle$ value $\rangle$ contains more than a single unhidden equal sign，it will be split at the first one and the others are considered part of the value．\ekvset should be nestable．
\ekvparse $\backslash e k v p a r s e\langle c s 1\rangle\langle c s 2\rangle\{\langle k e y\rangle=\langle$ value $\rangle, \ldots\}$
This macro parses the $\langle k e y\rangle=\langle$ value $\rangle$ pairs and provides those list elements which are only keys as the argument to $\langle\operatorname{cs} 1\rangle$ ，and those which are a $\langle$ key $\rangle=\langle$ value $\rangle$ pair to $\langle\operatorname{cs} 2\rangle$ as two arguments．It is fully expandable as well and returns the parsed list in \unexpanded， which has no effect outside of an \expanded or \edef context ${ }^{1}$ ．If you need control over the necessary steps of expansion you can use \expanded around it．
\ekvbreak，\ekvsneak，and \ekvchangeset and their relatives don＇t work in \ekvparse．It is analogue to expl3＇s \keyval＿parse：NNn，but not with the same parsing rules－\keyval＿parse：NNn throws an error on multiple equal signs per $\langle k e y\rangle=\langle$ value $\rangle$ pair and on empty $\langle k e y\rangle$ names in a $\langle k e y\rangle=\langle v a l u e\rangle$ pair，both of which \ekvparse doesn＇t deal with．

As a small example：
$\backslash e k v p a r s e \backslash h a n d l e k e y \backslash h a n d l e k e y v a l\{f o o=b a r$, key，baz＝\｛zzz\}\}
would expand to
$\backslash$ handlekeyval $\{$ foo $\}\{$ bar $\} \backslash h a n d l e k e y\{k e y\} \backslash h a n d l e k e y v a l\{b a z\}\{z z z\}$
and afterwards \handlekey and \handlekeyval would have to further handle the 〈key〉． There are no macros like these two contained in expkv，you have to set them up yourself if you want to use \ekvparse（of course the names might differ）．If you need the results of \ekvparse as the argument for another macro，you should use \expanded as only then the input stream will contain the output above：
$\backslash$ expandafter $\backslash$ handle $\backslash$ expanded $\{\backslash$ ekvparse $\backslash k \backslash k v\{f o o=b a r$, key，baz＝\｛zzz $\}\}\}$ would expand to
$\backslash$ handle $\backslash k v\{f o o\}\{b a r\} \backslash k\{k e y\} \backslash k v\{b a z\}\{z z z\}$

## 1．3 Miscellaneous

## 1．3．1 Other Macros

expkv provides some other macros which might be of interest．
\ekvVersion
\ekvDate
\ekvifdefined
\ekvifdefinedNoVal
\ekvifdefined\｛〈set $\rangle\}\{\langle$ key $\rangle\}\{\langle$ true $\rangle\}\{\langle$ false $\rangle\}$
$\backslash e k v i f d e f i n e d N o V a l\{\langle$ set $\rangle\}\{\langle$ key $\rangle\}\{\langle$ true $\rangle\}\{\langle$ false $\rangle\}$
These two macros test whether there is a $\langle k e y\rangle$ in $\langle$ set $\rangle$ ．It is false if either a hash table entry doesn＇t exist for that key or its meaning is \relax．

[^0]\ekvbreak
\ekvbreakPreSneak \ekvbreakPostSneak
\ekvbreak\｛〈after〉\}
Gobbles the remainder of the current \ekvset macro and its argument list and rein－ serts 〈after〉．So this can be used to break out of \ekvset．The first variant will also gobble anything that has been sneaked out using \ekvsneak or \ekvsneakPre， while \ekvbreakPreSneak will put 〈after〉 before anything that has been smuggled and \ekvbreakPostSneak will put 〈after〉 after the stuff that has been sneaked out．

## \ekvsneak

\ekvsneakPre
\ekvsneak\｛〈after〉\}
Puts $\langle$ after $\rangle$ after the effects of $\backslash e k v s e t$ ．The first variant will put $\langle a f t e r\rangle$ after any other tokens which might have been sneaked before，while \ekvsneakPre will put 〈after〉 before other smuggled stuff．This reads and reinserts the remainder of the current \ekvset macro and its argument list to do its job．A small usage example is shown in subsubsection 1．4．2．
\ekvchangeset $\{\langle$ new－set $\rangle\}$
Replaces the current set with $\langle$ new－set $\rangle$ ，so for the rest of the current \ekvset call，that call behaves as if it was called with \ekvset $\{\langle$ new－set $\rangle\}$ ．Just like \ekvsneak this reads and reinserts the remainder of the current \ekvset macro to do its job．It is comparable to using 〈key〉／．cd in pgfkeys．
\ekv＠name\｛〈set〉\}\{〈key〉\}
\ekv＠name＠set\｛\｛set〉\}
\ekv＠name＠key $\{\langle k e y\rangle\}$
The names of the macros that correspond to a key in a set are build with these macros．The default definition of \ekv＠name＠set is＂ekv〈set〉（＂and the default of \ekv＠name＠key is＂$\langle k e y\rangle)$＂．The complete name is build using \ekv＠name which is equivalent to \ekv＠name＠set $\{\langle\operatorname{set}\rangle\} \backslash$ ekv＠name＠key\｛〈key $\rangle\}$ ．For NoVal keys an additional N gets ap－ pended irrespective of these macros＇definition，so their name is $\backslash e k v\langle s e t\rangle(\langle k e y\rangle)$ N．You might redefine \ekv＠name＠set and \ekv＠name＠key locally but don’t redefine \ekv＠name！

## 1．3．2 Bugs

Just like keyval，expkv is bug free．But if you find bugshidden features ${ }^{2}$ you can tell me about them either via mail（see the first page）or directly on GitHub if you have an account there：https：／／github．com／Skillmon／tex＿expkv

## 1．3．3 Comparisons

Comparisons of speed are done with a very simple test key and the help of the l3benchmark package．The key and its usage should be equivalent to
$\backslash$ protected $\backslash e k v d e f\{$ test $\}\{h e i g h t\}\{\backslash$ def $\backslash m y h e i g h t\{\# 1\}\}$
\ekvset\｛test\}\{ height = 6 \}

[^1]and only the usage of the key, not its definition, is benchmarked. For the impatient, the essence of these comparisons regarding speed and buggy behaviour is contained in Table 1.

As far as I know expkv is the only fully expandable $\langle k e y\rangle=\langle$ value $\rangle$ parser. I tried to compare expkv to every $\langle$ key $\rangle=\langle$ value $\rangle$ package listed on CTAN, however, one might notice that some of those are missing from this list. That's because I didn't get the others to work due to bugs, or because they just provide wrappers around other packages in this list.

In this subsubsection is no benchmark of \ekvparse and \keyval_parse:NNn contained, as most other packages don't provide equivalent features to my knowledge. \ekvparse is slightly faster than \ekvset, but keep in mind that it does less. The same is true for \keyval_parse: NNn compared to \keys_set:nn of expl3 (where the difference is much bigger).
keyval is about 1.6 times faster and has a comparable feature set just a slightly different way how it handles keys without values. That might be considered a drawback, as it limits the versatility, but also as an advantage, as it might reduce doubled code. Keep in mind that as soon as someone loads xkeyval the performance of keyval gets replaced by xkeyval's.

Also keyval has a bug, which unfortunately can't really be resolved without breaking backwards compatibility for many documents, namely it strips braces from the argument before stripping spaces if the argument isn't surrounded by spaces, also it might strip more than one set of braces. Hence all of the following are equivalent in their outcome, though the last two lines should result in something different than the first two:
\setkeys $\{$ foo $\}\{b a r=b a z\}$
\setkeys $\{$ foo $\}\{$ bar $=\{b a z\}\}$
\setkeys $\{$ foo $\}\{b a r=\{b a z\}\}$
\setkeys $\{$ foo $\}\{b a r=\{\{b a z\}\}\}$
xkeyval is roughly seventeen times slower, but it provides more functionality, e.g., it has choice keys, boolean keys, and so on. It contains the same bug as keyval as it has to be compatible with it by design (it replaces keyval's frontend), but also adds even more cases in which braces are stripped that shouldn't be stripped, worsening the situation.

Itxkeys is over 370 times slower - which is funny, because it aims to be "[...] faster [...] than these earlier packages [referring to keyval and xkeyval]." It needs more time to parse zero keys than four of the packages in this comparison need to parse 100 keys. Since it aims to have a bigger feature set than xkeyval, it most definitely also has a bigger feature set than expkv. Also, it can't parse \long input, so as soon as your values contain a \par, it'll throw errors. Furthermore, Itxkeys doesn't strip outer braces at all by design, which, imho, is a weird design choice. In addition Itxkeys loads catoptions which is known to introduce bugs (e.g., see https://tex.stackexchange.com/questions/461783).

I3keys is around six times slower, but has an, imho, great interface to define keys. It strips all outer spaces, even if somehow multiple spaces ended up on either end. It offers more features, but is pretty much bound to expl3 code. Whether that's a drawback is up to you. Note that this comparison uses the version contained in $\mathrm{T}_{\mathrm{E}} \mathrm{XLive} 2019$ (frozen) which is a bit slower than versions starting with $\mathrm{T}_{\mathrm{E}} \mathrm{XLive} 2020$.
pgfkeys is around 2.7 times slower for one key, but has an enormous feature set. It has the same or a very similar bug keyval has. The brace bug (and also the category fragility) can be fixed by pgfkeyx, but this package was last updated in 2012 and it slows down $\backslash$ pgfkeys by factor 8. Also I don't know whether this might introduce new bugs.
kvsetkeys with kvdefinekeys is about 3.7 times slower, but it works even if commas and equals have category codes different from 12 (just as some other packages in this list). Else the features of the keys are equal to those of keyval, the parser has more features, though.
options is 1.5 times slower for only a single value. It has a much bigger feature set. Unfortunately it also suffers from the premature unbracing bug keyval has.
simplekv is hard to compare because I don't speak French (so I don't understand the documentation) and from what I can see, there is no direct way to define the equivalent test key. Nevertheless, I tested the closest possible equivalent of my test key while siding for simplekv's design not forcing something into it it doesn't seem to be designed for. It is almost seven times slower and has hard to predict behaviour regarding brace and space stripping, similar to keyval. The tested definition was:

```
\usepackage\{simplekv\}
\setKVdefault[simplekv]\{height=\{ abc \}\} \% key setup
\(\backslash\) setKV[simplekv]\{ height \(=6\) \} \% benchmarked
```

yax is over twenty times slower. It has a pretty strange syntax, imho, and again a direct equivalent is hard to define. It has the premature unbracing bug, too. Also somehow loading yax broke options for me. The tested definition was:
kslash\)defactiveparameteryax\{\storevalue$\backslashmyheight~yax:height~\}~\%~key~setup~$\setparameterlist\{yax\}\{height$=6$\}$\quad\%$benchmarkedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefined

### 1.4 Examples

### 1.4.1 Standard Use-Case

Say we have a macro for which we want to create a $\langle$ key $\rangle=\langle$ value $\rangle$ interface. The macro has a parameter, which is stored in the dimension \ourdim having a default value from its initialization. Now we want to be able to change that dimension with the width key to some specified value. For that we'd do

```
\newdimen\ourdim
\ourdim=150pt
\protected \ekvdef{our}{width}{\ourdim=#1\relax}
```

as you can see, we use the set our here. We want the key to behave different if no value is specified. In that case the key should not use its initial value, but be smart and determine the available space from \hsize, so we also define
$\backslash$ protected $\backslash$ ekvdefNoVal\{our $\}\{$ width $\}\{\backslash$ ourdim $=.9 \backslash$ hsize \}
Now we set up our macro to use this $\langle$ key $\rangle=\langle$ value $\rangle$ interface

Table 1: Comparison of $\langle k e y\rangle=\langle$ value $\rangle$ packages. The packages are ordered from fastest to slowest for one $\langle k e y\rangle=\langle$ value $\rangle$ pair. Benchmarking was done using l3benchmark and the scripts in the Benchmarks folder of the git repository. The columns $p_{i}$ are the polynomial coefficients of a linear fit to the run-time, $p_{0}$ can be interpreted as the overhead for initialisation and $p_{1}$ the cost per key. The $T_{0}$ column is the actual mean ops needed for an empty list argument, as the linear fit doesn't match that point well in general. The column "BB" lists whether the parsing is affected by some sort of brace bug, "CF" stands for category code fragile and lists whether the parsing breaks with active commas or equal signs.

| Package | $p_{1}$ | $p_{0}$ | $T_{0}$ | BB | CF | Date |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| keyval | 13.4 | 2.6 | 6.9 | yes | yes | 2014-10-28 |
| expk ${ }^{\text {d }}$ | 19.7 | 7.3 | 9.6 | no | no | 2020-04-07 |
| options | 23.4 | 15.6 | 19.9 | yes | yes | 2015-03-01 |
| pgfkeys | 24.6 | 46.2 | 52.8 | yes | yes | 2020-01-08 |
| kvsetkeys | * | * | 39.8 | no | no | 2019-12-15 |
| l3keys | 108.0 | 56.3 | 50.7 | no | no | 2020-02-25 |
| simplekv | 149.2 | 25.7 | 8.1 | yes | yes | 2017-o8-o8 |
| xkeyval | 248.8 | 234.4 | 161.9 | yes | yes | 2014-12-03 |
| yax | 443.6 | 170.5 | 115.7 | yes | yes | 2010-01-22 |
| Itxkeys | 3516.0 | 5254.8 | 5487.0 | no | no | 2012-11-17 |

*For kvsetkeys the linear model used for the other packages is a poor fit, kvsetkeys seems to have approximately quadratic run-time, the coefficients of the second degree polynomial fit are $p_{2}=7.5, p_{1}=51.3$, and $p_{0}=49.6$. Of course the other packages might not really have linear run-time, but at least from 1 to 20 keys the fits don't seem too bad (the maximum ratio $p_{2} / p_{1}$ for the other packages is $3.5 \times 10^{-3}$ ). If one extrapolates the fits for $100\langle$ key $\rangle=\langle$ value $\rangle$ pairs one finds that most of them match pretty well, the exception being ltxkeys, which behaves quadratic as well with $p_{2}=11.4, p_{1}=3276.7$, and $p_{0}=6132.3$.
$\backslash$ protected $\backslash$ def $\backslash$ ourmacro\#1 $\{\backslash$ begingroup $\backslash$ ekvset $\{o u r\}\{\# 1\} \backslash$ the $\backslash$ ourdim $\backslash$ endgroup $\}$
Finally we can use our macro like in the following

| $\backslash$ ourmacro $\} \backslash$ par | 150.0 pt |
| :--- | :--- |
| $\backslash$ ourmacro $\{$ width $\} \backslash$ par | 192.85382 pt |
| \ourmacro $\{$ width $=5$ pt $\} \backslash$ par | $5 . \mathrm{opt}$ |

The same key using expkVidef Using expkvidef we can set up the equivalent key using a $\langle k e y\rangle=\langle v a l u e\rangle$ interface, after the following we could use \ourmacro in the same way as above. expkVIDef will allocate and initialise \ourdim and define the width key \protected for us, so the result will be exactly the same - with the exception that the default will use \ourdim=. 9 h hsize $\backslash$ relax instead.

```
\input expkv-def % or \usepackage{expkv-def}
\ekvdefinekeys{our}
    {
        dimen width = \ourdim,
        qdefault width = .9\hsize,
        initial width = 150pt
    }
```

1.4.2 An Expandable $\langle$ key $\rangle=\langle$ value $\rangle$ Macro Using \ekvsneak

Let's set up an expandable macro, that uses a $\langle k e y\rangle=\langle$ value $\rangle$ interface. The problems we'll face for this are:

1. ignoring duplicate keys
2. default values for keys which weren't used
3. providing the values as the correct argument to a macro (ordered)

First we need to decide which $\langle k e y\rangle=\langle$ value $\rangle$ parsing macro we want to do this with, \ekvset or \ekvparse. For this example we also want to show the usage of \ekvsneak, hence we'll choose \ekvset. And we'll have to use \ekvset such that it builds a parsable list for our macro internals. To gain back control after \ekvset is done we have to put an internal of our macro at the start of that list, so we use an internal key that uses \ekvsneakPre after any user input.

To ignore duplicates will be easy if the value of the key used last will be put first in the list, so the following will use \ekvsneakPre for the user-level keys. If we wanted some key for which the first usage should be the binding one we would use \ekvsneak instead for that key.

Providing default values can be done in different ways, we'll use a simple approach in which we'll just put the outcome of our keys if they were used with default values before the parsing list terminator.

Ordering the keys can be done simply by searching for a specific token for each argument which acts like a flag, so our sneaked out values will include specific tokens acting as markers.

Now that we have answers for our technical problems, we have to decide what our example macro should do. How about we define a macro that calculates the sine of a number and rounds that to a specified precision? As a small extra this macro should
understand input in radian and degree and the used trigonometric function should be selectable as well. For the hard part of this task (expandably evaluating trigonometric functions) we'll use the xfp package.

First we set up our keys according to our earlier considerations and set up the user facing macro \sine. The end marker of the parsing list will be a \sine@stop token, which we don't need to define and we put our defaults right before it.

```
\RequirePackage{xfp }
\ m a k e a t l e t t e r ~
\ekvdef{expex }{f}{\ekvsneakPre{\f{#1}}}
\ekvdef{expex }{round}{\ekvsneakPre{\rnd{#1 }}}
\ekvdefNoVal { expex }{degree }{\ekvsneakPre{\deg{d}}}
\ekvdefNoVal{ expex }{radian }{\ekvsneakPre{\deg { }}}
\ekvdefNoVal{expex}{internal}{\ekvsneakPre{\sine@rnd}}
\newcommand*\sine[2]
    {\ekvset{expex}{#1,internal}\rnd{3}\deg{d}\f{sin}\sine@stop{#2}}
```

For the sake of simplicity we defined the macro \sine with two mandatory arguments, the first being the $\langle k e y\rangle=\langle v a l u e\rangle$ list, the second the argument to the trigonometric function. We could've used xparse's facilities here to define an expandable macro which takes an optional argument instead.

Now we need to define some internal macros to extract the value of each key's last usage (remember that this will be the group after the first special flag-token). For that we use one delimited macro per key.
$\backslash d e f \backslash$ sine@rnd\#1 \rnd\#2\#3\sine@stop $\{\backslash \operatorname{sine@deg\# 1\# 3\backslash sine@stop~\{ \# 2~\} ~\} ~}$
$\backslash d e f \backslash$ sine@deg\#1\deg\#2\#3\sine@stop $\{\backslash \operatorname{sine@f\# 1\# 3\backslash sine@stop~\{ \# 2~\} \} ~}$
$\backslash d e f \backslash \operatorname{sine@f\# 1\backslash f\# 2\# 3\backslash sine@stop\{ \backslash sine@final\{ \# 2\} \} }$
After the macros \sine@rnd, \sine@deg, and \sine@f the macro \sine@final will see $\backslash$ sine@final $\{\langle f\rangle\}\{\langle$ degree/radian $\rangle\}\{\langle$ round $\rangle\}\{\langle$ num $\rangle\}$. Now \sine@final has to expandably deal with those arguments such that the \fpeval macro of xfp gets the correct input. Luckily this is pretty straight forward in this example. In \fpeval the trigonometric functions have names such as $\sin$ or $\cos$ and the degree taking variants sind or cosd. And since the degree key puts a d in \#2 and the radian key leaves \#2 empty all we have to do to get the correct function name is stick the two together.

```
\newcommand*\sine@final[4]{\fpeval{round(#1#2(#4),#3)}}
\ \text { \makeatother}
```

Let's test our macro:

| $\backslash$ sine $\}\{60\} \backslash$ par | 0.866 |
| :---: | :---: |
| $\backslash$ sine $\{$ round $=10\}\{60\} \backslash$ par | 0.8660254038 |
| $\backslash$ sine $\{f=\cos$, radian $\}\{p i\} \backslash p a r$ | -1 |
| $\backslash e d e f \backslash$ myval $\{\backslash$ sine $\{f=$ tan $\}\{1\}\} \backslash$ texttt $\{\backslash$ meaning $\backslash$ myval $\}$ | macro:->0.01 |

The same macro using expkvics Using expkvics we can set up something equivalent with a bit less code. The implementation chosen in expkvics is more efficient than the example above and way easier to code.

[^2]```
    {
        f=sin,
        unit=d,
        round=3,
    }
\ekvcSecondaryKeys\sine
    {
        nmeta degree={unit=d},
        nmeta radian={unit={}},
    }
\newcommand*\sine@[4]{\fpeval{round(#1#2(#4),#3)}}
\ \text { Imakeatother}
```

The resulting macro will behave just like the one previously defined, but will have an additional unit key, since in expkvics every argument must have a value taking key which defines it.

### 1.5 Error Messages

expkv should only send messages in case of errors, there are no warnings and no info messages. In this subsection those errors are listed.

### 1.5.1 Load Time

expkv.tex checks whether $\varepsilon-\mathrm{T}_{\mathrm{E}} \mathrm{X}$ is available. If it isn't, an error will be thrown using \errmessage:
! expkv Error: e-TeX required.

### 1.5.2 Defining Keys

If you get any error from expkv while you're trying to define a key, the definition will be aborted and gobbled.

If you try to define a key with an empty set name you'll get:
! expkv Error: empty set name not allowed.
Similarly, if you try to define a key with an empty key name:
! expkv Error: empty key name not allowed.
Both of these messages are done in a way that doesn't throw additional errors due to \global, \long, etc., not being used correctly if you prefixed one of the defining macros.

### 1.5.3 Using Keys

This subsubsection contains the errors thrown during \ekvset. The errors are thrown in an expandable manner by providing an undefined macro. In the following messages <key> gets replaced with the problematic key's name, and <set> with the corresponding set. If any errors during $\langle k e y\rangle=\langle$ value $\rangle$ handling are encountered, the entry in the comma separated list will be omitted after the error is thrown and the next $\langle k e y\rangle=\langle v a l u e\rangle$ pair will be parsed.

If you're using an undefined key you'll get:

```
! Undefined control sequence.
<argument> \! expkv Error:
unknown key ('<key>', set '<set>').
```

If you're using a key for which only a normal version and no NoVal version is defined, but don't provide a value, you'll get:
! Undefined control sequence. <argument> \! expkv Error:

```
value required ('<key>', set '<set>').
```

If you're using a key for which only a NoVal version and no normal version is defined, but provide a value, you'll get:

```
! Undefined control sequence.
<argument> \! expkv Error:
value forbidden ('<key>', set '<set>').
```

If you're using a set for which you never executed one of the defining macros from subsection 1.1 you'll get a low level $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ error, as that isn't actively tested by the parser (and hence will lead to undefined behaviour and not be gracefully ignored). The error will look like
! Missing \endcsname inserted.
<to be read again>
\! expkv Error: Set '<set>' undefined.

### 1.6 License

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## 2 Implementation

### 2.1 The $\mathrm{It}_{\mathrm{E}} \mathrm{X}$ Package

First we set up the $\mathrm{LA}_{\mathrm{E}} \mathrm{X}$ package. That one doesn't really do much except \inputting the generic code and identifying itself as a package.

```
\def\ekv@tmp
    {%
        \ProvidesFile{expkv.tex}%
            [\ekvDate\space v\ekvVersion\space an expandable key=val implementation]%
    }
\input{expkv.tex}
\ProvidesPackage{expkv}%
    [\ekvDate\space v\ekvVersion\space an expandable key=val implementation]
```


### 2.2 The Generic Code

The rest of this implementation will be the generic code.
We make sure that it's only input once:

```
\expandafter\ifx\csname ekvVersion\endcsname\relax
\else
    \expandafter\endinput
\i
```

    Check whether \(\varepsilon-\mathrm{T}_{\mathrm{E}} \mathrm{X}\) is available - expkv requires \(\varepsilon-\mathrm{T}_{\mathrm{E}} \mathrm{X}\).
    \begingroup $\backslash$ expandafter $\backslash$ expandafter $\backslash$ expandafter $\backslash e n d g r o u p$
\expandafter\ifx\csname numexpr\endcsname\relax
\errmessage\{expkv requires e-TeX\}
\expandafter\endinput
\fi
\ekvVersion
We're on our first input, so lets store the version and date in a macro.
\ekvDate 18 \def \ekvVersion\{1.2\}
19 \def \ekvDate\{2020-04-10\}
(End definition for \ekvVersion and \ekvDate. These functions are documented on page 4.)
If the ${ }^{\mathrm{A}} \mathrm{T}_{\mathrm{E}} \mathrm{X}$ format is loaded we want to be a good file and report back who we are, for this the package will have defined \ekv@tmp to use \ProvidesFile, else this will expand to a \relax and do no harm.
20 \csname ekv@tmp \endcsname
Store the category code of @ to later be able to reset it and change it to 11 for now.
21 \expandafter \chardef \csname ekv@tmp\endcsname=\catcode'\@
22 \catcode'\@=11
\ekv@tmp might later be reused to gobble any prefixes which might be provided to \ekvdef and similar in case the names are invalid, we just temporarily use it here as means to store the current category code of © to restore it at the end of the file, we never care for the actual definition of it.
\@gobble \@firstofone \@firstoftwo
\@secondoftwo \ekv@gobbleto@stop \ekv@fi@secondoftwo
\ekv@gobble@mark \ekv@gobble@from@mark@to@stop
\ekv@ifempty
\ekv@ifempty@ \ekv@ifempty@true \ekv@ifempty@false \ekv@ifempty@true@F
\ekv@ifempty@true@F@gobble \ekv@ifempty@true@F@gobbletwo

Since branching tests are often more versatile than \if . . . \else. . . \fi constructs, we define helpers that are branching pretty fast. Also here are some other utility functions that just grab some tokens. The ones that are also contained in ${ }^{\mathrm{ET}_{\mathrm{E}} \mathrm{X}}$ don't use the ekv prefix.

```
\long\def\@gobble#1{}
\long\def\@firstofone#1{#1}
\long\def\@firstoftwo#1#2{#1}
\long\def\@secondoftwo#1#2{#2}
\long\def\ekv@fi@secondoftwo\fi\@firstoftwo#1#2{\fi#2}
\long\def\ekv@gobbleto@stop#1\ekv@stop{}
\def\ekv@gobble@mark\ekv@mark{}
\long\def\ekv@gobble@from@mark@to@stop\ekv@mark#1\ekv@stop{}
(End definition for \@gobble and others.)
```

As you can see \ekv@gobbleto@stop uses a special marker \ekv@stop. The package will use three such markers, the one you've seen already, \ekv@mark and \ekv@nil. Contrarily to how for instance expl3 does things, we don't define them, as we don't need them to have an actual meaning. This has the advantage that if they somehow get expanded - which should never happen if things work out - they'll throw an error directly.

We can test for a lot of things building on an if-empty test, so lets define a really fast one. Since some tests might have reversed logic (true if something is not empty) we also set up macros for the reversed branches.

```
\long\def\ekv@ifempty#1%
    {%
        \ekv@ifempty@\ekv@ifempty@A#1\ekv@ifempty@B\ekv@ifempty@true
            \ekv@ifempty@A\ekv@ifempty@B\@secondoftwo
    }
\long\def\ekv@ifempty@#1\ekv@ifempty@A\ekv@ifempty@B{}
\long\def\ekv@ifempty@true\ekv@ifempty@A\ekv@ifempty@B\@secondoftwo#1#2{#1}
\long\def\ekv@ifempty@false\ekv@ifempty@A\ekv@ifempty@B\@firstoftwo#1#2{#2}
\long\def\ekv@ifempty@true@F\ekv@ifempty@A\ekv@ifempty@B\@firstofone#1{}
\long\def\ekv@ifempty@true@F@gobble\ekv@ifempty@A\ekv@ifempty@B\@firstofone#1#2%
    {}
\long\def\ekv@ifempty@true@F@gobbletwo
            \ekv@ifempty@A\ekv@ifempty@B\@firstofone#1#2#3%
        {}
```

(End definition for \ekv@ifempty and others.)
\ekv@ifblank@ The obvious test that can be based on an if-empty is if-blank, meaning a test checking whether the argument is empty or consists only of spaces. Our version here will be tweaked a bit, as we want to check this, but with one leading \ekv@mark token that is to be ignored.
45 \long\def\ekv@ifblank@\ekv@mark\#1\{\ekv@ifempty@\ekv@ifempty@A\}
(End definition for \ekv@ifblank@.)
\ekv@ifdefined
We'll need to check whether something is defined quite frequently, so why not define a macro that does this. The following test is expandable, slower than the typical expandable test for undefined control sequences, but faster for defined ones. Since we want to be as fast as possible for correct input, this is to be preferred.

```
\def\ekv@ifdefined#1%
    {%
        \expandafter
        \ifx\csname\ifcsname #1\endcsname #1\else relax\fi\endcsname\relax
            \ekv@fi@secondoftwo
        \i
        \@firstoftwo
    }
```

(End definition for \ekv@ifdefined.)

```
\ekv@name The keys will all follow the same naming scheme, so we define it here.
\ekv@name@set
54 \def\ekv@name\#1\#2\{\ekv@name@set\{\#1\}\ekv@name@key\{\#2\}\}
\ekv@name@key 55 \def\ekv@name@set\#1\{ekv\#1(\}
\({ }_{56}\) \def \(\backslash\) ekv@name@key\#1\{\#1)\}
(End definition for \ekv@name, \ekv@name@set, and \ekv@name@key. These functions are documented on page 5.)
```

We can misuse the macro name we use to expandably store the set-name in a single token - since this increases performance drastically, especially for long set-names to throw a more meaningful error message in case a set isn't defined. The name of \ekv@undefined@set is a little bit misleading, as it is called in either case inside of \csname, but the result will be a control sequence with meaning \relax if the set is undefined, hence will break the \csname building the key-macro which will throw the error message.

```
57 \def\ekv@undefined@set#1{! expkv Error: Set '#1' undefined.}
(End definition for \ekv@undefined@set.)
```

\ekv@checkvalid We place some restrictions on the allowed names, though, namely sets and keys are not allowed to be empty - blanks are fine (meaning set- or key-names consisting of spaces). The \def \ekv@tmp gobbles any $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ prefixes which would otherwise throw errors. This will, however, break the package if an \outer has been gobbled this way. I consider that good, because keys shouldn't be defined \outer anyways.

```
\protected\def\ekv@checkvalid#1#2%
    {%
        \ekv@ifempty{#1}%
            {%
            \def \ekv@tmp{}%
            \errmessage{expkv Error: empty set name not allowed}%
            }%
            {%
            \ekv@ifempty{#2}%
                        {%
                        \def\ekv@tmp{}%
                            \errmessage{expkv Error: empty key name not allowed}%
                    }%
                    \@secondoftwo
            }%
        \@gobble
    }
(End definition for \ekv@checkvalid.)
```

\ekvifdefined And provide user-level macros to test whether a key is defined.
\ekvifdefinedNoVal 75 \def\ekvifdefined\#1\#2\{\ekv@ifdefined\{\ekv@name\{\#1\}\{\#2\}\}\}
${ }_{76}$ \def \ekvifdefinedNoVal\#1\#2\{\ekv@ifdefined\{\ekv@name\{\#1\}\{\#2\}N\}\}
(End definition for \ekvifdefined and \ekvifdefinedNoVal. These functions are documented on page 4.)
\ekvdef Set up the key defining macros \ekvdef etc.
\ekvdefNoVal
\protected $\backslash$ long $\backslash$ def $\backslash$ ekvdef\#1\#2\#3\%
\{\%
\ekvletNoVal 79 \ekv@checkvalid\{\#1\}\{\#2\}\%
\ekvletkv
\ekvletkvNoVal
\{\%
\expandafter\def\csname\ekv@name\{\#1\}\{\#2\}\endcsname\#\#1\{\#3\}\%
\expandafter\ekv@defset\csname\ekv@undefined@set\{\#1\}\endcsname\{\#1\}\%
3\%
\}
\protected $\backslash$ long \def $\backslash$ ekvdefNoVal\#1\#2\#3\% \{\%
\ekv@checkvalid\{\#1\}\{\#2\}\%
\{\%
\expandafter $\backslash$ def $\backslash c$ sname $\backslash e k v @ n a m e\{\# 1\}\{\# 2\} N \backslash e n d c s n a m e\{\# 3\} \%$
\expandafter\ekv@defset\csname\ekv@undefined@set\{\#1\}\endcsname\{\#1\}\% \}\%
\}
\protected\def\ekvlet\#1\#2\#3\%
\{\%
\ekv@checkvalid\{\#1\}\{\#2\}\%
\{\%
\expandafter\let\csname\ekv@name\{\#1\}\{\#2\}\endcsname\#3\%
\expandafter\ekv@defset\csname\ekv@undefined@set\{\#1\}\endcsname\{\#1\}\%
\}\%
\}
\protected\def \ekvletNoVal\#1\#2\#3\%
\{\%
\ekv@checkvalid\{\#1\}\{\#2\}\%
\{\%
\expandafter\let\csname\ekv@name\{\#1\}\{\#2\}N\endcsname\#3\%
\expandafter\ekv@defset\csname\ekv@undefined@set\{\#1\}\endcsname\{\#1\}\%
\}\%
\}
\protected\def \ekvletkv\#1\#2\#3\#4\%
\{\%
\ekv@checkvalid\{\#1\}\{\#2\}\%
\{\%
\expandafter\let\csname\ekv@name\{\#1\}\{\#2\}\expandafter\endcsname
\csname\ekv@name\{\#3\}\{\#4\}\endcsname
\expandafter\ekv@defset\csname\ekv@undefined@set\{\#1\}\endcsname\{\#1\}\%
3\%
\}
\protected\def \ekvletkvNoVal\#1\#2\#3\#4\%
\{\%
\ekv@checkvalid\{\#1\}\{\#2\}\%
\{\%
\expandafter\let\csname\ekv@name\{\#1\}\{\#2\}N\expandafter\endcsname
\csname\ekv@name\{\#3\}\{\#4\}N\endcsname

```
124 \expandafter\ekv@defset\csname\ekv@undefined@set{#1}\endcsname{#1}%
}%
}
(End definition for \ekvdef and others. These functions are documented on page 2.)
```

In order to enhance the speed the set name given to \ekvset will be turned into a control sequence pretty early, so we have to define that control sequence.

```
\protected\def\ekv@defset#1#2%
{%
            \ifx#1\relax
                \edef#1##1{\ekv@name@set{#2}\ekv@name@key{##1}}%
            \i
}
(End definition for \ekv@defset.)
```

\ekvset Set up \ekvset, which should not be affected by active commas and equal signs. The equal signs are a bit harder to cope with and we'll do that later, but replacing the active commas with commas of category other can be done beforehand. That's why we define \ekvset here with a temporary meaning just to set up the things with two different category codes. \#1 will be a , ${ }_{13}$ and \#2 will be a $=_{13}$.

```
\def\ekvset#1#2{%
\endgroup
\long\def\ekvset##1##2%
    {%
            \expandafter\ekv@set\csname\ekv@undefined@set{##1}\endcsname
                \ekv@mark##2#1\ekv@stop#1{}%
    }
```

(End definition for \ekvset. This function is documented on page 3.)
 there were unprotected other commas and resplit there.
140 \long\def \ekv@set\#\#1\#\#2\#1\%
\{\%
Test whether we're at the end, if so invoke \ekv@endset,
142 \ekv@gobble@from@mark@to@stop\#\#2\ekv@endset\ekv@stop
else go on with other commas,
143 \ekv@set@other\#\#1\#\#2, \ekv@stop,\%
and get the next active comma delimited $\langle$ key $\rangle=\langle$ value $\rangle$ pair.

```
444 \ekv@set##1\ekv@mark
145 }
```

(End definition for \ekv@set.)
\ekv@endset \ekv@endset is a hungry little macro. It will eat everything that remains of \ekv@set and unbrace the sneaked stuff.

146 \long\def \ekv@endset
\ekv@stop\ekv@set@other\#\#1, \ekv@stop, \ekv@set\#\#2\ekv@mark
\#\#3\%
149 \{\#\#3\}
\ekv@eq@other \ekv@eq@active
\ekv@set@other
\ekv@set@eq@other@a \ekv@set@eq@other@b

Splitting at equal signs will be done in a way that checks whether there is an equal sign and splits at the same time. This gets quite messy and the code might look complicated, but this is pretty fast (faster than first checking for an equal sign and splitting if one is found). The splitting code will be adapted for \ekvset and \ekvparse to get the most speed, but some of these macros don't require such adaptions. \ekv@eq@other and \ekv@eq@active will split the argument at the first equal sign and insert the macro which comes after the first following \ekv@mark. This allows for fast branching based on $\mathrm{T}_{\mathrm{E} X}$ 's argument grabbing rules and we don't have to split after the branching if the equal sign was there.

```
\long\def\ekv@eq@other##1=##2\ekv@mark##3##4\ekv@stop
    {%
        ##3##1\ekv@stop\ekv@mark##2%
        }
\long\def\ekv@eq@active##1#2##2\ekv@mark##3##4\ekv@stop
    {%
        ##3##1\ekv@stop\ekv@mark##2%
    }
(End definition for \ekv@eq@other and \ekv@eq@active.)
```

The macro \ekv@set@other is guaranteed to get only single $\langle k e y\rangle=\langle$ value $\rangle$ pairs.

```
\long\def\ekv@set@other##1##2,%
    {%
```

First we test whether we're done.
160 \ekv@gobble@from@mark@to@stop\#\#2\ekv@endset@other\ekv@stop

If not we split at the equal sign of category other.

```
161 \ekv@eq@other##2\ekv@nil\ekv@mark\ekv@set@eq@other@a
162 =\ekv@mark\ekv@set@eq@active\ekv@stop
```

And insert the set name and the next recursion step of \ekv@set@other.

```
463 ##1%
    \ekv@set@other##1\ekv@mark
    }
```

(End definition for \ekv@set@other.)

The first of these two macros runs the split-test for equal signs of category active. It will only be inserted if the $\langle k e y\rangle=\langle$ value $\rangle$ pair contained at least one equal sign of category other and \#\#1 will contain everything up to that equal sign.

```
\long\def\ekv@set@eq@other@a##1\ekv@stop
    {%
        \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@set@eq@other@active@a
            #2\ekv@mark\ekv@set@eq@other@b\ekv@stop
    }
```

The second macro will have been called by \ekv@eq@active if no active equal sign was found. All it does is remove the excess tokens of that test and forward the $\langle k e y\rangle=\langle$ value $\rangle$ pair to \ekv@set@pair.

```
171 \long\def\ekv@set@eq@other@b
    ##1\ekv@nil\ekv@mark\ekv@set@eq@other@active@a\ekv@stop\ekv@mark
172 {% %
```

```
74 \ekv@strip{##1}\ekv@set@pair
    }
```

(End definition for \ekv@set@eq@other@a and \ekv@set@eq@other@b.)
\ekv@set@eq@other@active@a \ekv@set@eq@other@active@b
\ekv@set@eq@other@active@a will be called if the $\langle k e y\rangle=\langle v a l u e\rangle$ pair was wrongly split on an equal sign of category other but has an earlier equal sign of category active. \#\#1 will be the contents up to the active equal sign and \#\#2 everything that remains until the first found other equal sign. It has to reinsert the equal sign and passes things on to \ekv@set@eq@other@active@b which calls \ekv@set@pair on the then correctly split $\langle k e y\rangle=\langle v a l u e\rangle$ pair.
${ }_{176}$ \long\def $\backslash e k v @ s e t @ e q @ o t h e r @ a c t i v e @ a \# \# 1 \backslash e k v @ s t o p \# \# 2 \backslash e k v @ n i l \backslash e k v @ m a r k ~$

```
    {%
```

            \ekv@set@eq@other@active@b\{\#\#1\}\#\#2=\%
        \}
    \long\def \ekv@set@eq@other@active@b\#\#1\%
\{\%
\ekv@strip\{\#\#1\}\ekv@set@pair
\}
(End definition for \ekv@set@eq@other@active@a and \ekv@set@eq@other@active@b.)
\ekv@set@eq@active \ekv@set@eq@active@
\ekv@set@eq@active will be called when there was no equal sign of category other in the $\langle$ key $\rangle=\langle$ value $\rangle$ pair. It removes the excess tokens of the prior test and split-checks for an active equal sign.

```
\long\def\ekv@set@eq@active
            ##1\ekv@nil\ekv@mark\ekv@set@eq@other@a\ekv@stop\ekv@mark
        {%
            \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@set@eq@active@
                #2\ekv@@mark\ekv@set@noeq\ekv@stop
    }
```

If an active equal sign was found in \ekv@set@eq@active we'll have to pass the now split $\langle k e y\rangle=\langle v a l u e\rangle$ pair on to \ekv@set@pair.
90 \long\def \ekv@set@eq@active@\#\#1 \ekv@stop

## \{\%

\ekv@strip\{\#\#1\}\ekv@set@pair \}
(End definition for \ekv@set@eq@active and \ekv@set@eq@active@.)
\ekv@set@noeq
If no active equal sign was found by \ekv@set@eq@active there is no equal sign contained in the parsed list entry. In that case we have to check whether the entry is blank in order to ignore it (in which case we'll have to gobble the set-name which was put after these tests by \ekv@set@other). Else this is a NoVal key and the entry is passed on to \ekv@set@key.

```
94 \long\def\ekv@set@noeq##1\ekv@nil\ekv@mark\ekv@set@eq@active@\ekv@stop\ekv@mark
195 {%
196 \ekv@ifblank@##1\ekv@nil\ekv@ifempty@B\ekv@ifempty@true@F@gobble
97 \ekv@ifempty@A\ekv@ifempty@B\@firstofone
98 {\ekv@strip{##1}\ekv@set@key}%
199 }
(End definition for \ekv@set@noeq.)
```

\ekv@endset@other
All that's left for \ekv@set@other is the macro which breaks the recursion loop at the end. This is done by gobbling all the remaining tokens.
200 \long\def \ekv@endset@other
201 \ekv@stop
202 \ekv@eq@other\#\#1\ekv@nil\ekv@mark\ekv@set@eq@other@a
203 = \ekv@mark\ekv@set@eq@active\ekv@stop
204 \#\#2\%
205 \ekv@set@other\#\#3\ekv@mark
206 \{\}
(End definition for \ekv@endset@other.)
\ekvbreak Provide macros that can completely stop the parsing of \ekvset, who knows what it'll \ekvbreakPreSneak \ekvbreakPostSneak be useful for.
207 \long\def \ekvbreak\#\#1\#\#2\ekv@stop\#1\#\#3\{\#\#1\}
\long\def \ekvbreakPreSneak \#\#1\#\#2\ekv@stop\#1\#\#3\{\#\#1\#\#3\}
209 \long\def \ekvbreakPostSneak\#\#1\#\#2\ekv@stop\#1\#\#3\{\#\#3\#\#1\}
(End definition for \ekvbreak, \ekvbreakPreSneak, and \ekvbreakPostSneak. These functions are documented on page 5.)
\ekvsneak One last thing we want to do for \ekvset is to provide macros that just smuggle stuff

## \ekvsneakPre

 after \ekvset's effects.```
210 \long\def\ekvsneak##1##2\ekv@stop#1##3%
        {%
            ##2\ekv@stop#1{##3##1}%
        }
\long\def\ekvsneakPre##1##2\ekv@stop#1##3%
        {%
            ##2\ekv@stop#1{##1##3}%
        }
```

(End definition for \ekvsneak and \ekvsneakPre. These functions are documented on page 5.)
\ekvparse Additionally to the \ekvset macro we also want to provide an \ekvparse macro, that has the same scope as \keyval_parse: NNn from expl3. This is pretty analogue to the \ekvset implementation, we just put an \unexpanded here and there instead of other macros to stop the \expanded on our output.

```
218 \long\def\ekvparse##1##2##3%
```

219 \{\%
220 \ekv@parse\#\#1\#\#2\ekv@mark\#\#3\#1 \ekv@stop\#1\%
221 \}
(End definition for \ekvparse. This function is documented on page 4.)

## \ekv@parse

```
28 \long\def\ekv@endparse
        \ekv@stop\ekv@parse@other##1,\ekv@stop,\ekv@parse##2\ekv@mark
        {}
(End definition for \ekv@endparse.)
```

```
31 \long\def\ekv@parse@other##1##2##3,%
```

31 \long\def\ekv@parse@other\#\#1\#\#2\#\#3,%
{%
{%
\ekv@gobble@from@mark@to@stop\#\#3\ekv@endparse@other\ekv@stop
\ekv@gobble@from@mark@to@stop\#\#3\ekv@endparse@other\ekv@stop
\ekv@eq@other\#\#3\ekv@nil\ekv@mark\ekv@parse@eq@other@a
\ekv@eq@other\#\#3\ekv@nil\ekv@mark\ekv@parse@eq@other@a
=\ekv@mark\ekv@parse@eq@active\ekv@stop
=\ekv@mark\ekv@parse@eq@active\ekv@stop
\#\#1\#\#2%
\#\#1\#\#2%
\ekv@parse@other\#\#1\#\#2\ekv@mark
\ekv@parse@other\#\#1\#\#2\ekv@mark
}
}
(End definition for \ekv@parse@other.)

```
(End definition for \ekv@parse@other.)
```

    \ekv@parse@other
    \ekv@parse@eq@other@a
    \ekv@parse@eq@other@b
    ```
\long\def\ekv@parse@eq@other@a##1\ekv@stop
    {%
            \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@parse@eq@other@active@a
            #2\ekv@mark\ekv@parse@eq@other@b\ekv@stop
        }
\long\def\ekv@parse@eq@other@b
        ##1\ekv@nil\ekv@mark\ekv@parse@eq@other@active@a\ekv@stop\ekv@mark
        {%
            \ekv@strip{##1}\ekv@parse@pair
    }
```

(End definition for \ekv@parse@eq@other@a and \ekv@parse@eq@other@b.)

```
\long\def\ekv@parse@eq@other@active@a##1\ekv@stop##2\ekv@nil\ekv@mark
        {%
            \ekv@parse@eq@other@active@b{##1}##2=%
        }
\long\def\ekv@parse@eq@other@active@b##1%
    {%
            \ekv@strip{##1}\ekv@parse@pair
        }
(End definition for \ekv@parse@eq@other@active@a and \ekv@parse@eq@other@active@b.)
```

    \ekv@parse@eq@active
    \ekv@parse@eq@active@
    257
\long\def\ekv@parse@eq@active
\#\#1 \ekv@nil\ekv@mark\ekv@parse@eq@other@a\ekv@stop\ekv@mark
\{\%
\ekv@eq@active\#\#1 \ekv@nil\ekv@mark\ekv@parse@eq@active@
\#2\ekv@mark\ekv@parse@noeq\ekv@stop
\}
\long\def\ekv@parse@eq@active@\#\#1 \ekv@stop
\{\%
\ekv@strip\{\#\#1\}\ekv@parse@pair
\}

```
(End definition for \ekv@parse@eq@active and \ekv@parse@eq@active@.)
```

\ekv@parse@noeq

```
\long\def\ekv@parse@noeq
            ##1\ekv@nil\ekv@mark\ekv@parse@eq@active@\ekv@stop\ekv@mark
        {%
            \ekv@ifblank@##1\ekv@nil\ekv@ifempty@B\ekv@ifempty@true@F@gobbletwo
                \ekv@ifempty@A\ekv@ifempty@B\@firstofone
                {\ekv@strip{##1}\ekv@parse@key}%
        }
(End definition for \ekv@parse@noeq.)
```

\ekv@endparse@other

```
\long\def\ekv@endparse@other
        \ekv@stop
        \ekv@eq@other##1\ekv@nil\ekv@mark\ekv@parse@eq@other@a
        =\ekv@mark\ekv@parse@eq@active\ekv@stop
        ##2%
        \ekv@parse@other##3\ekv@mark
        {}
```

(End definition for \ekv@endparse@other.)
\ekv@parse@pair
\ekv@parse@pair@

```
\long\def\ekv@parse@pair##1##2\ekv@nil
        {%
            \ekv@strip{##2}\ekv@parse@pair@{##1}%
        }
\long\def\ekv@parse@pair@##1##2##3##4%
        {%
            \unexpanded{##4{##2}{##1}}%
        }
```

(End definition for \ekv@parse@pair and \ekv@parse@pair@.)
\ekv@parse@key
\long\def \ekv@parse@key\#\#1\#\#2\#\#3\%
\{\%
\unexpanded\{\#\#2\{\#\#1\}\}\%
\}
(End definition for \ekv@parse@key.)
Finally really setting things up with \ekvset's temporary meaning:
$293\}$
294 \begingroup
295 \catcode' $\backslash,=13$
296 ไcatcode' $\backslash==13$
297 \ekvset,=
\ekvchangeset Provide a macro that is able to switch out the current $\langle\operatorname{set}\rangle$ in \ekvset. This operation is slow (by comparison, it should be slightly faster than \ekvsneak), but allows for something similar to pgfkeys's $\langle k e y\rangle /$. cd mechanism. However this operation is more expensive than /.cd as we can't just redefine some token to reflect this, but have to switch out the set expandably, so this works similar to the \ekvsneak macros reading and reinserting the remainder of the $\langle$ key $\rangle=\langle$ value $\rangle$ list.

```
29 \def\ekvchangeset#1%
299 {%
300 \expandafter\ekv@changeset\csname\ekv@undefined@set{#1}\endcsname\ekv@mark
}
```

(End definition for $\backslash e k v c h a n g e s e t$. This function is documented on page 5.)
This macro does the real change-out of \ekvchangeset. We introduced an \ekv@mark to not accidentally remove some braces which we have to remove again.
302 \long\def \ekv@changeset\#1\#2\ekv@set@other\#3\#4\ekv@set\#5\%
303 \{\%
〔\%
\ekv@gobble@mark\#2\ekv@set@other\#1\#4\ekv@set\#1\%
\}
(End definition for \ekv@changeset.)
\ekv@set@pair
\ekv@set@pair gets invoked with the space and brace stripped key-name as its first argument, the value as the second argument, and the set name as the third argument. It builds the key-macro name and provides everything to be able to throw meaningful error messages if it isn't defined. \ekv@set@pair@ will space and brace strip the value if the macro is defined and call the key-macro. Else it'll branch into the error messages provided by \ekv@set@pair.

```
\long\def\ekv@set@pair#1#2\ekv@nil#3%
    {%
        \expandafter\ekv@set@pair@
            \csname
                    \ifcsname #3{#1}\endcsname
                #3{#1}%
            \else
                relax%
            \fi
            \endcsname
            {#2}%
            {%
                    \ekv@ifdefined{#3{#1}N}%
                    \ekv@err@noarg
                    \ekv@err@unknown
                    #3{#1}%
            }%
    }
\long\def\ekv@set@pair@#1#2%
    {%
        \ifx#1\relax
            \ekv@fi@secondoftwo
        \i
        \@firstoftwo
        {\ekv@strip{#2}#1}%
    }
```

```
(End definition for \ekv@set@pair.)
```

\ekv@set@key
Analogous to \ekv@set@pair, \ekv@set@key builds the NoVal key-macro and provides an error-branch. \ekv@set@key@ will test whether the key-macro is defined and if so call it, else the errors are thrown.

```
```

\long\def\ekv@set@key\#1\#2%

```
```

\long\def\ekv@set@key\#1\#2%
{%
{%
\expandafter\ekv@set@key@
\expandafter\ekv@set@key@
\csname
\csname
\ifcsname \#2{\#1}N\endcsname
\ifcsname \#2{\#1}N\endcsname
\#2{\#1}N%
\#2{\#1}N%
\else
\else
relax%
relax%
\i
\i
\endcsname
\endcsname
{%
{%
\ekv@ifdefined{\#2{\#1}}%
\ekv@ifdefined{\#2{\#1}}%
\ekv@err@reqval
\ekv@err@reqval
\ekv@err@unknown
\ekv@err@unknown
\#2{\#1}%
\#2{\#1}%
}%
}%
}
}
\def\ekv@set@key@\#1%
\def\ekv@set@key@\#1%
{%
{%
\ifx\#1\relax
\ifx\#1\relax
\ekv@fi@secondoftwo
\ekv@fi@secondoftwo
\fi
\fi
\@firstoftwo\#1%
\@firstoftwo\#1%
}
}
(End definition for \ekv@set@key.)

```
```

(End definition for \ekv@set@key.)

```
```

\ekv@err lekv@err@

Since \ekvset is fully expandable as long as the code of the keys is (which is unlikely) we want to somehow throw expandable errors, in our case via undefined control sequences.
\ekv@err@common \ekv@err@common@ \ekv@err@unknown
\ekv@err@noarg \ekv@err@reqval

```
\begingroup
```

\begingroup
\edef\ekv@err
\edef\ekv@err
{%
{%
\endgroup
\endgroup
\unexpanded{\long\def\ekv@err}\#\#1%
\unexpanded{\long\def\ekv@err}\#\#1%
{%
{%
\unexpanded{\expandafter\ekv@err@\@firstofone}%
\unexpanded{\expandafter\ekv@err@\@firstofone}%
{\unexpanded\expandafter{\csname ! expkv Error:\endcsname}\#\#1.}%
{\unexpanded\expandafter{\csname ! expkv Error:\endcsname}\#\#1.}%
\unexpanded{\ekv@stop}%
\unexpanded{\ekv@stop}%
}%
}%
}
}
\ekv@err
\ekv@err
\def\ekv@err@{\expandafter\ekv@gobbleto@stop}
\def\ekv@err@{\expandafter\ekv@gobbleto@stop}
(End definition for \ekv@err and \ekv@err@.)

```
(End definition for \ekv@err and \ekv@err@.)
```

Now we can use \ekv@err to set up some error messages so that we can later use those instead of the full strings.
369 \long\def \ekv@err@common \#1\#2\{\expandafter\ekv@err@common@\string\#2\{\#1\}\}
${ }_{370} \backslash l o n g \backslash d e f \backslash e k v @ e r r @ c o m m o n @ \# 1 ' \# 2 ' ~ \# 3 . \# 4 \# 5\{\backslash e k v @ e r r\{\# 4 ~(' \# 5 ', ~ s e t ~ ' \# 2 ')\}\}$
\long\def\ekv@err@noarg \#1\#2\{\ekv@err@common\{value forbidden\}\#1\{\#2\}\}
73 \long\def\ekv@err@reqval \#1\#2\{\ekv@err@common\{value required\}\#1\{\#2\}\}
(End definition for \ekv@err@common and others.)
\ekv@strip \ekv@strip@a \ekv@strip@b \ekv@strip@c

Finally we borrow some ideas of expl3's l3tl to strip spaces from keys and values. This \ekv@strip also strips one level of outer braces after stripping spaces, so an input of \{abc\} becomes abc after stripping. It should be used with \#1 prefixed by \ekv@mark. Also this implementation at most strips one space from both sides.

```
\def\ekv@strip#1%
    {%
            \long\def\ekv@strip##1%
                {%
                    \ekv@strip@a
                    ##1%
                    \ekv@nil
                    \ekv@mark#1%
                    #1\ekv@nil{}%
                    \ekv@stop
                }%
            \long\def\ekv@strip@a##1\ekv@mark#1##2\ekv@nil##3%
                {%
                    \ekv@strip@b##3##1##2\ekv@nil
                }%
            \long\def\ekv@strip@b##1#1\ekv@nil
                {%
                    \ekv@strip@c##1\ekv@nil
                }%
            \long\def\ekv@strip@c\ekv@mark##1\ekv@nil##2\ekv@stop##3%
                {%
                    ##3{##1}%
                }%
    }
\ekv@strip{ }
(End definition for \ekv@strip and others.)
```

    Now everything that's left is to reset the category code of @.
    399 \catcode'\@=\ekv@tmp

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[^0]:    ${ }^{1}$ This is a change in behaviour，previously（vo． 3 and before）\ekvparse would expand in exactly two steps． This isn＇t always necessary，but makes the parsing considerably slower．If this is necessary for your application you can put an \expanded around it and will still be faster since you need only a single \expandafter this way．

[^1]:    ${ }^{2}$ Thanks，David！

[^2]:    $\backslash$ makeatletter
    $\backslash e k v c S p l i t A n d F o r w a r d \backslash$ sine $\backslash$ sine@

